

Gödel

Hashing

`matt.might.net`

`@mattnight`

Disclaimer

“simple, fun idea”

“simple, fun idea”

“works well in practice,”

“simple, fun idea”

“works well in practice,”

“but theory says it will not.”

An old problem

An older solution

A big impact

An old problem

“CFA is slow!”

An older solution



Gödel hashing



functional
monotonic
compact
dynamic
incremental
perfect

Inspired by a true theorem.

Word-level parallelism!

Great cache behavior!

A big impact

Minutes of work

2x

2x5x

2x 5x 8x

^{2x} 5x 8x

1000x

Motivation

(f x)

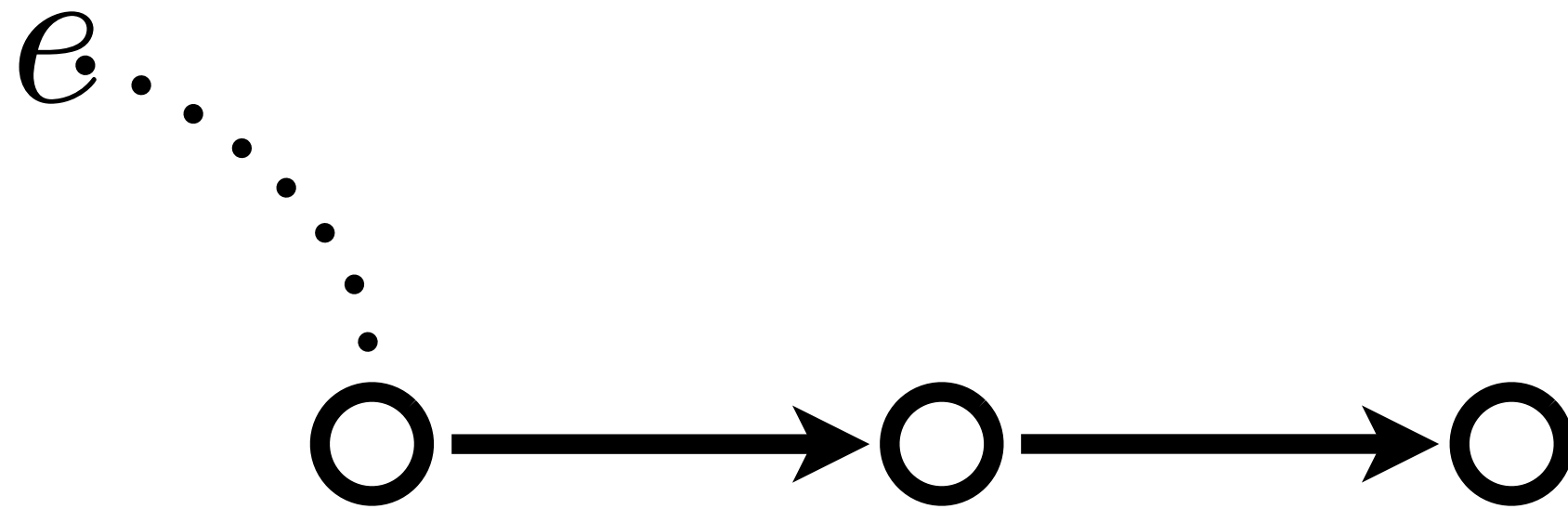
$f(\mathbf{x})$

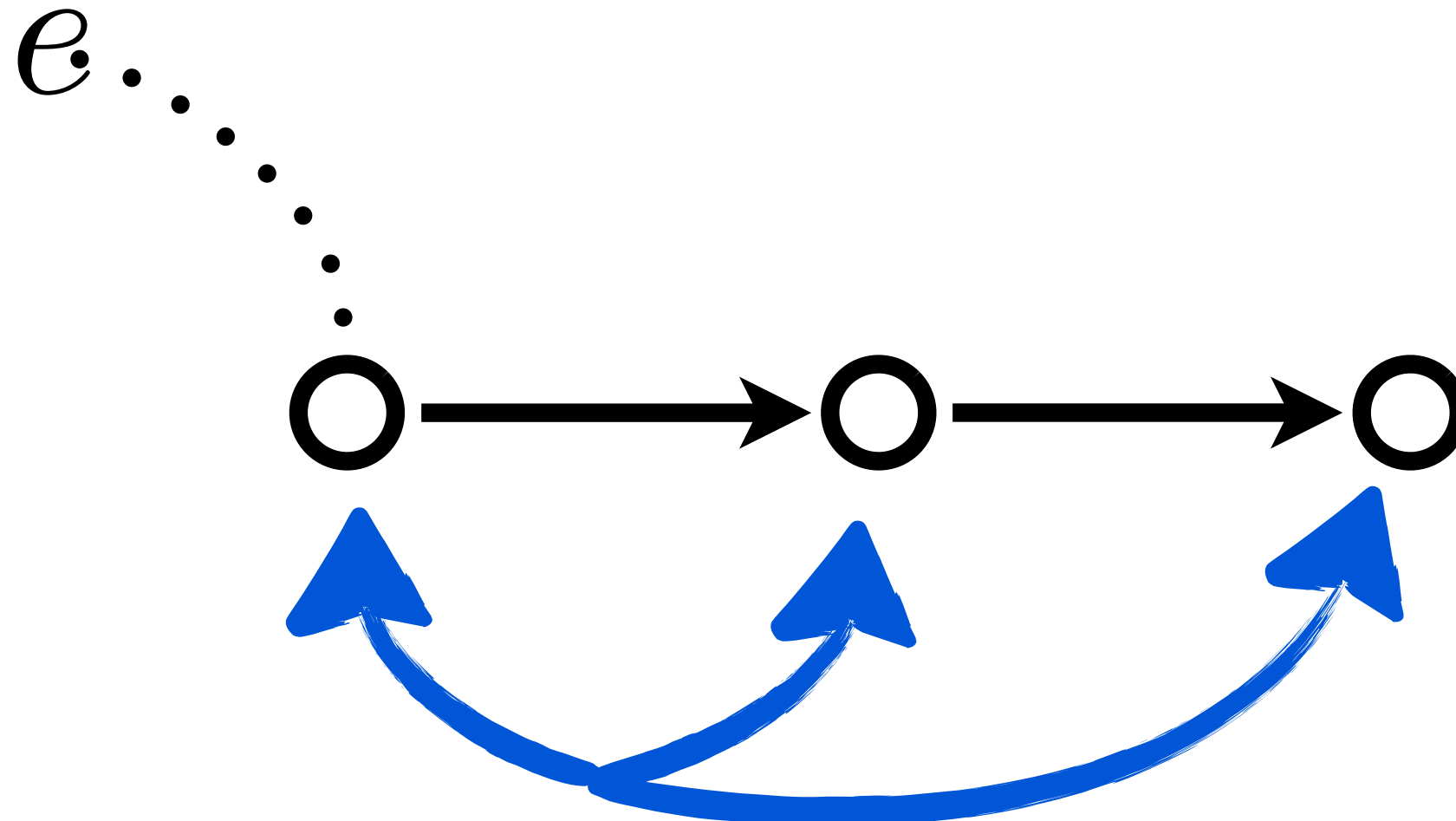
What is f ?

Why not run the program?

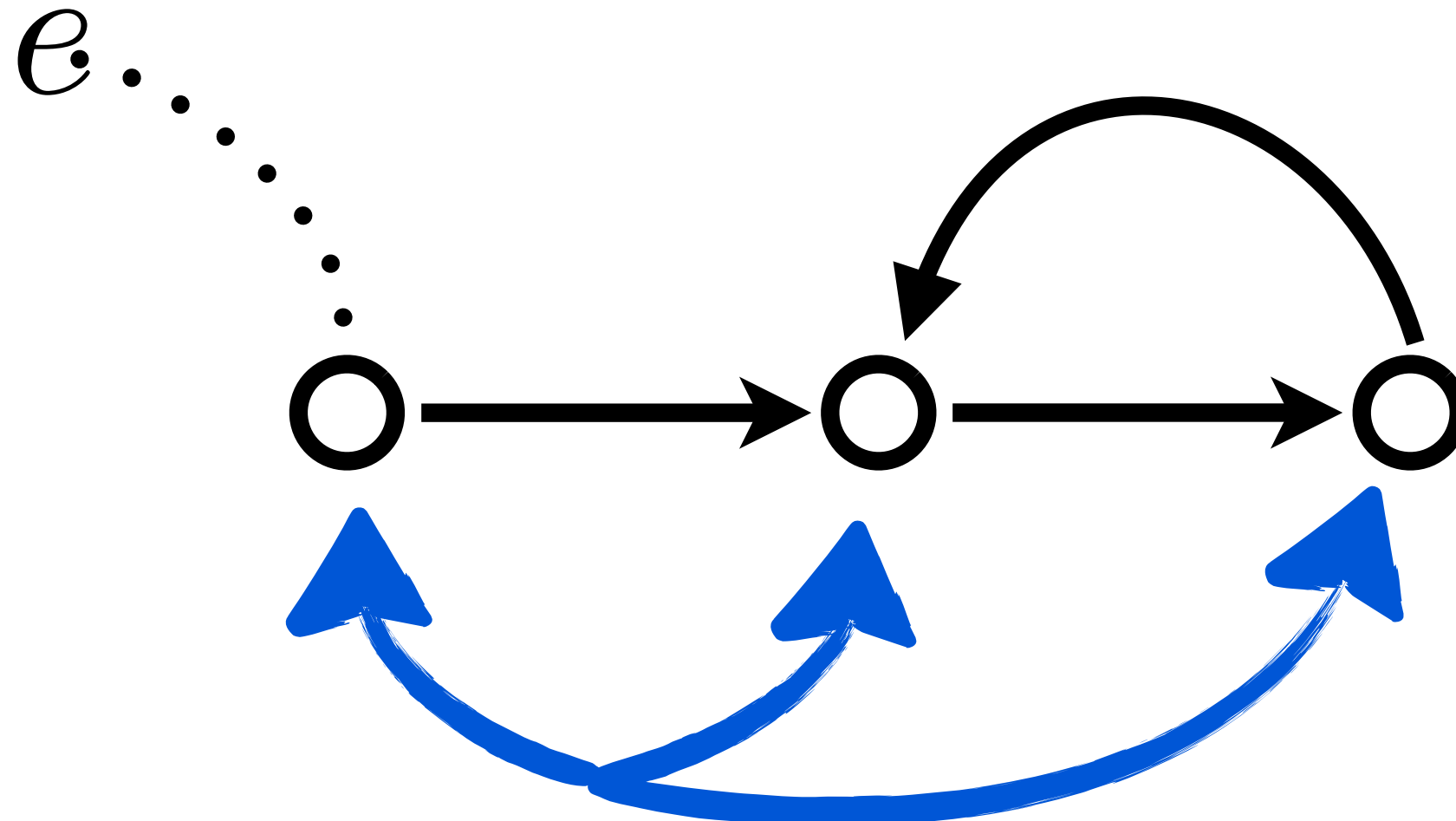
e

e

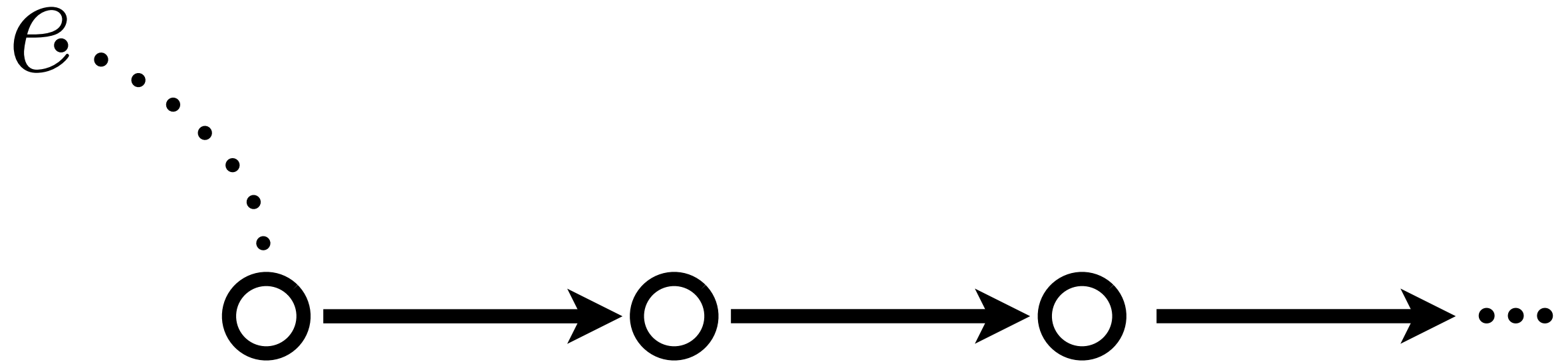




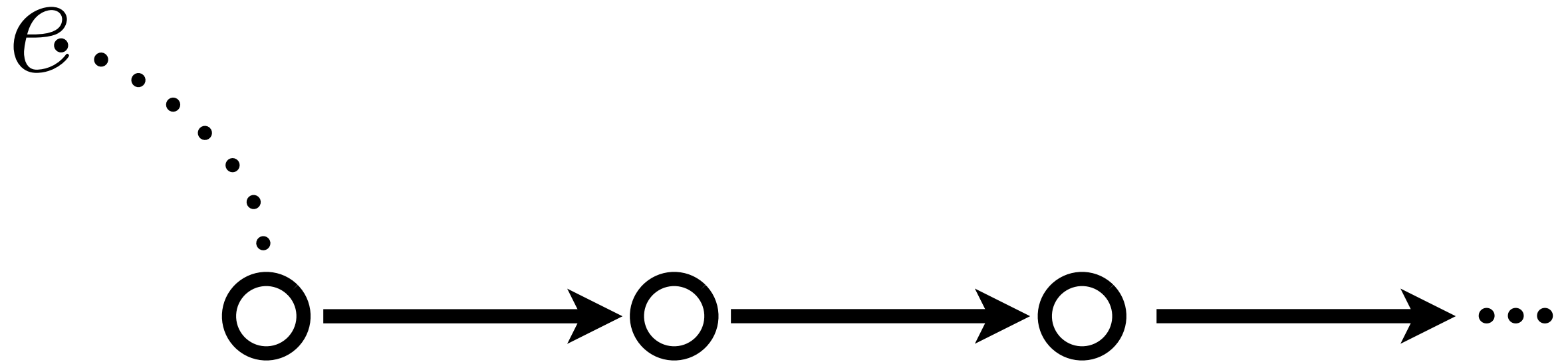
What is f , here?

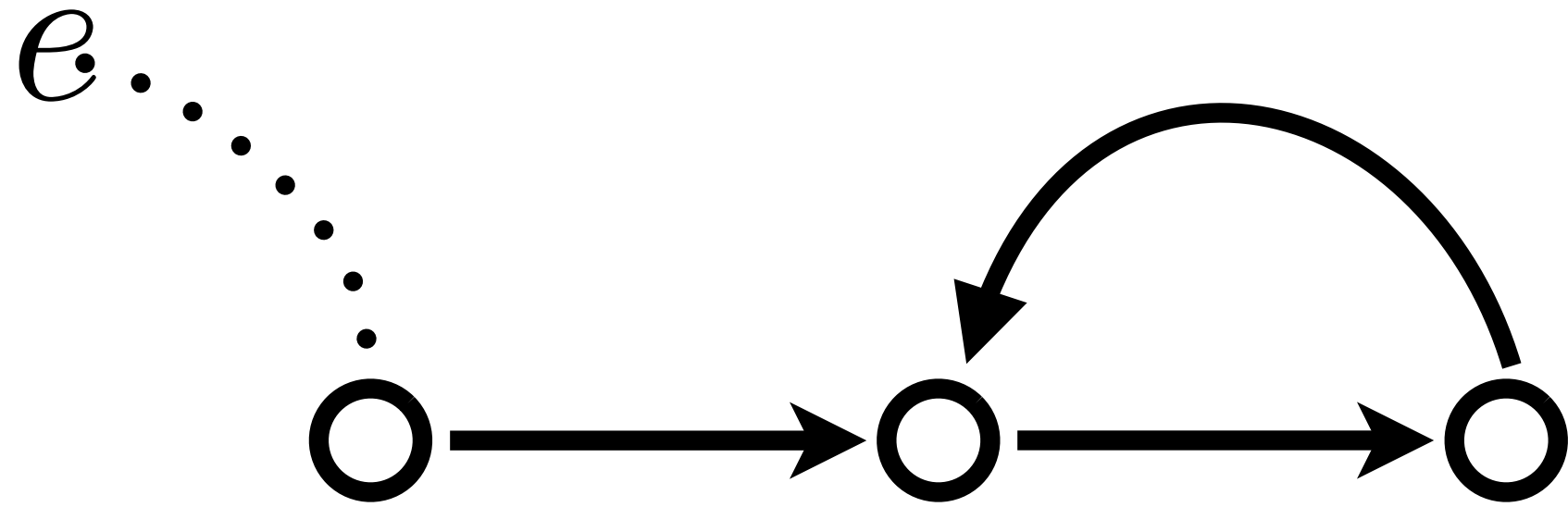


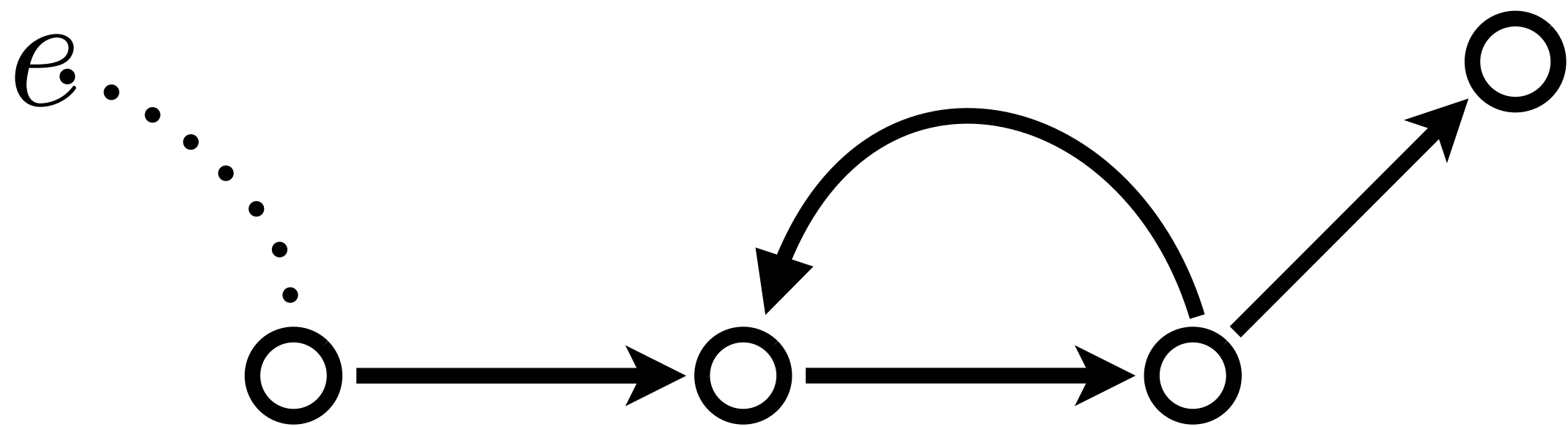
What is f , here?



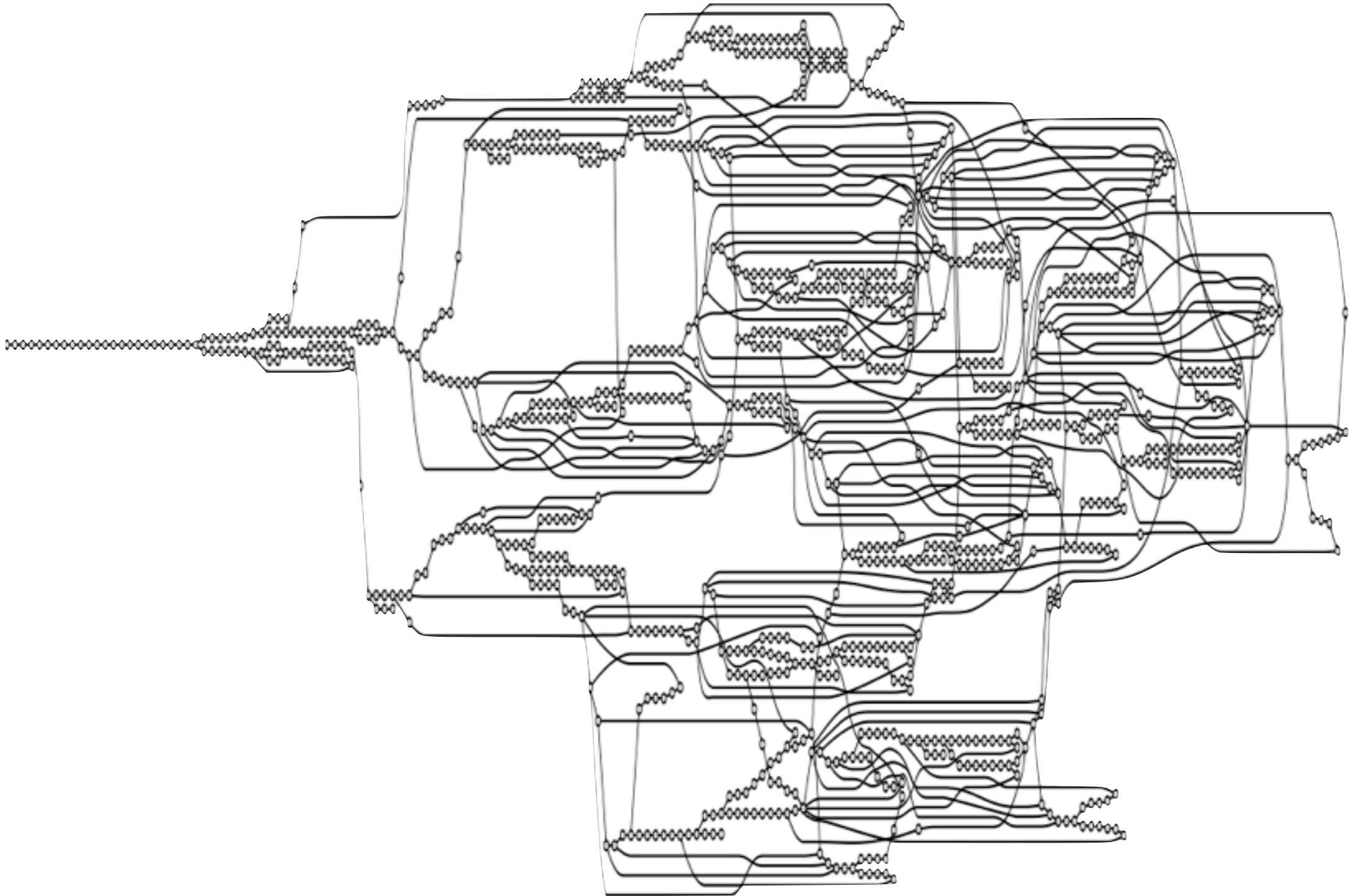
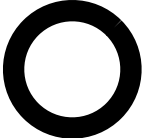


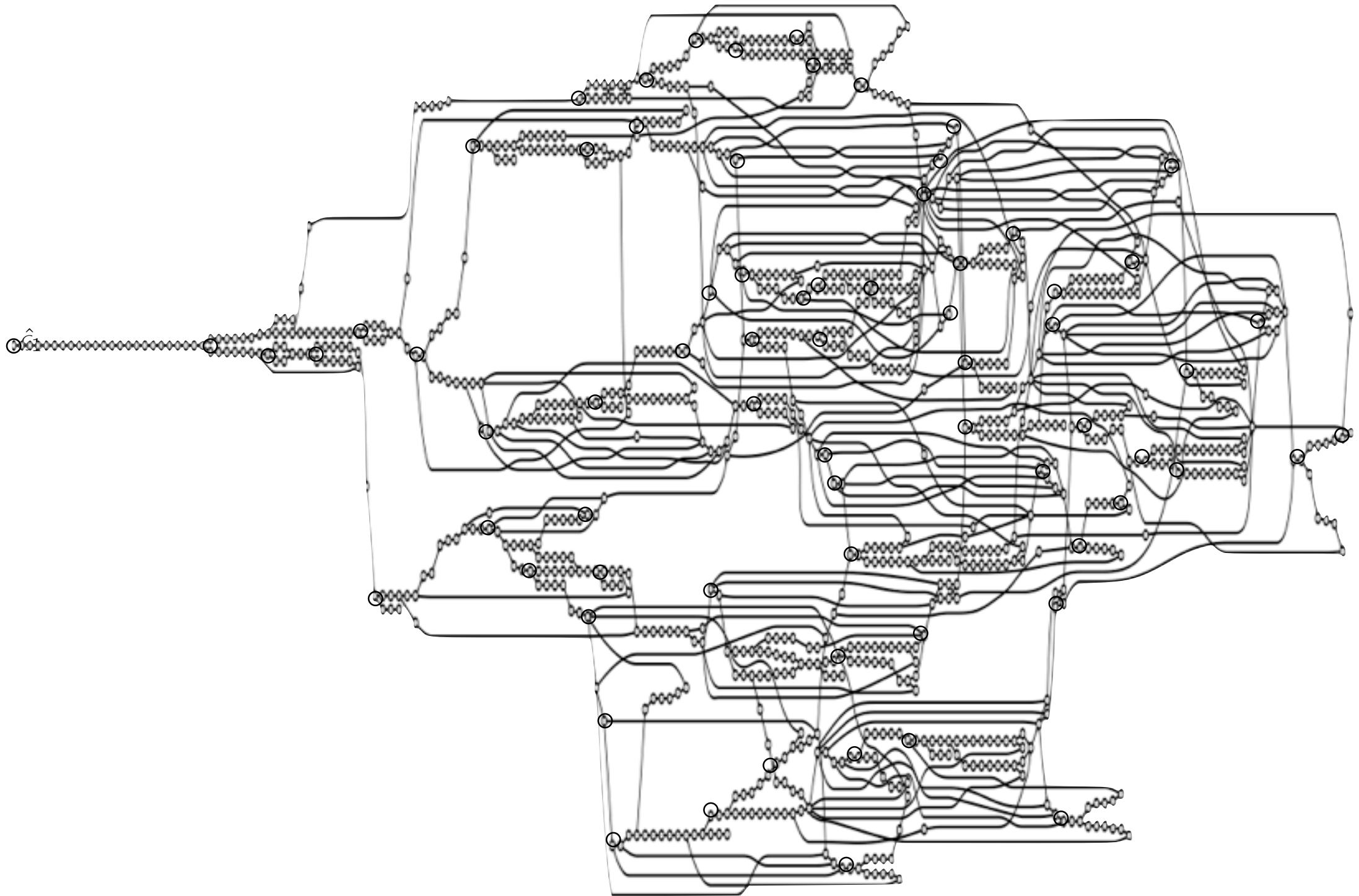
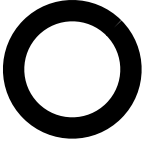


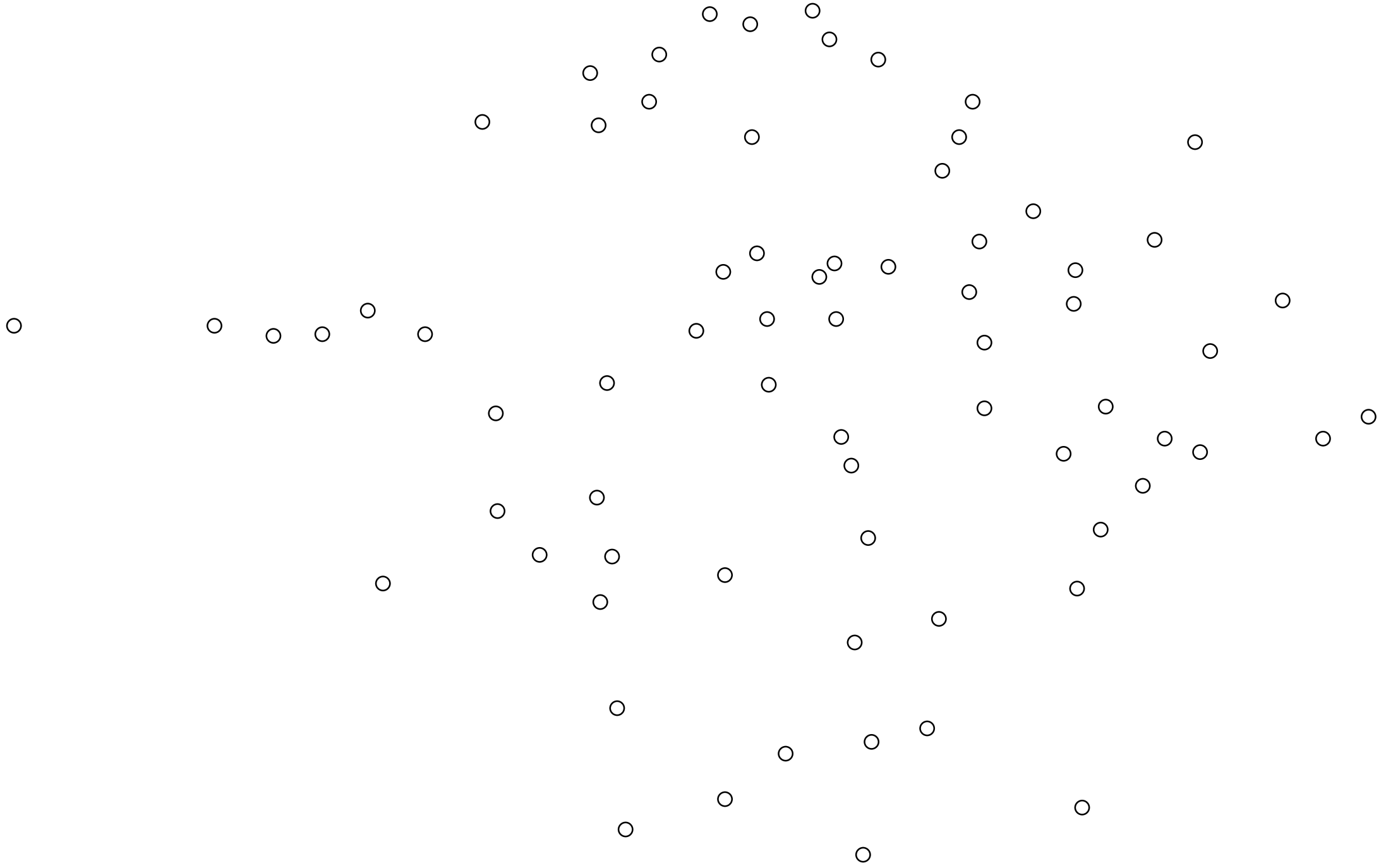
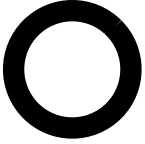


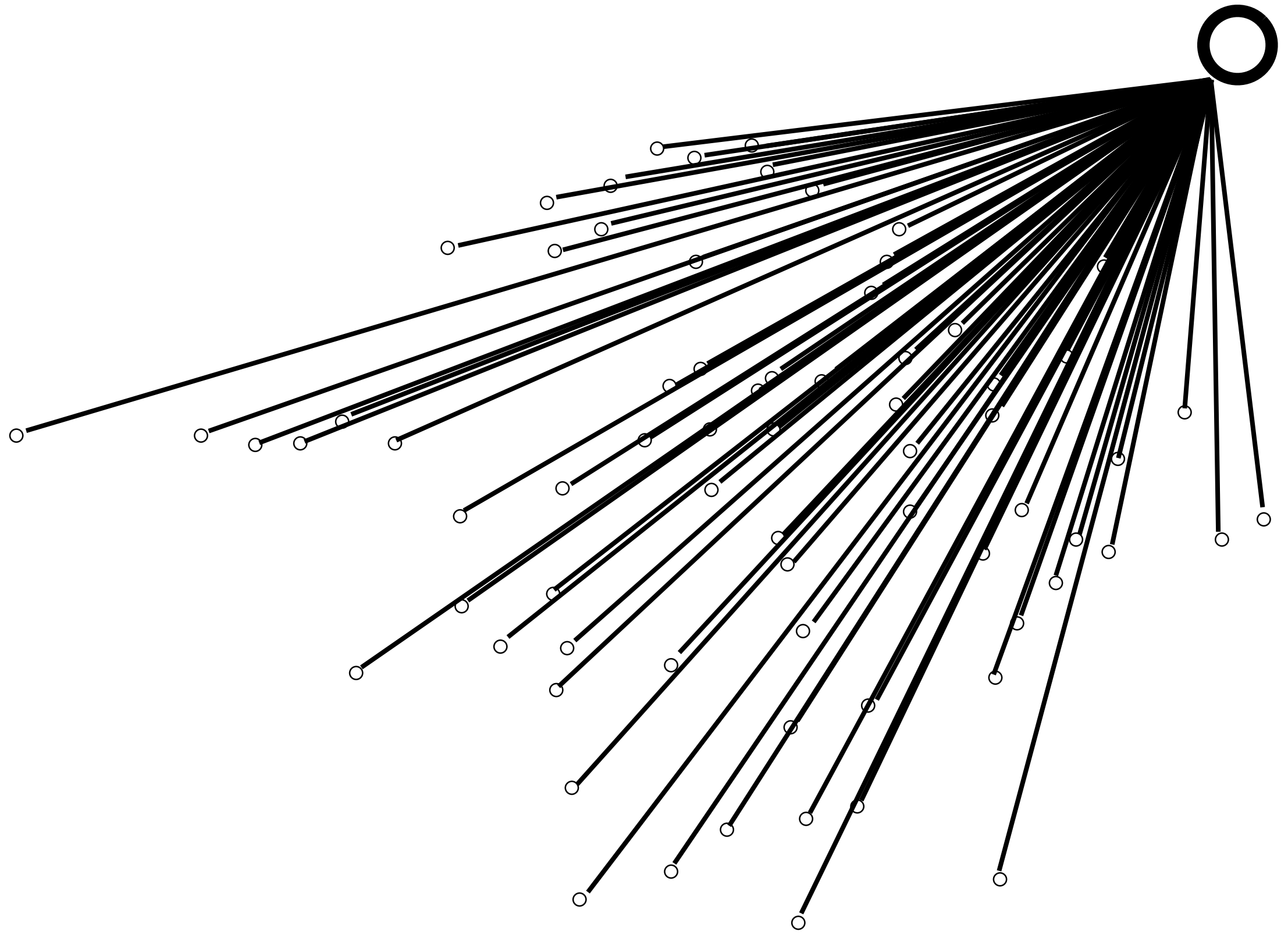


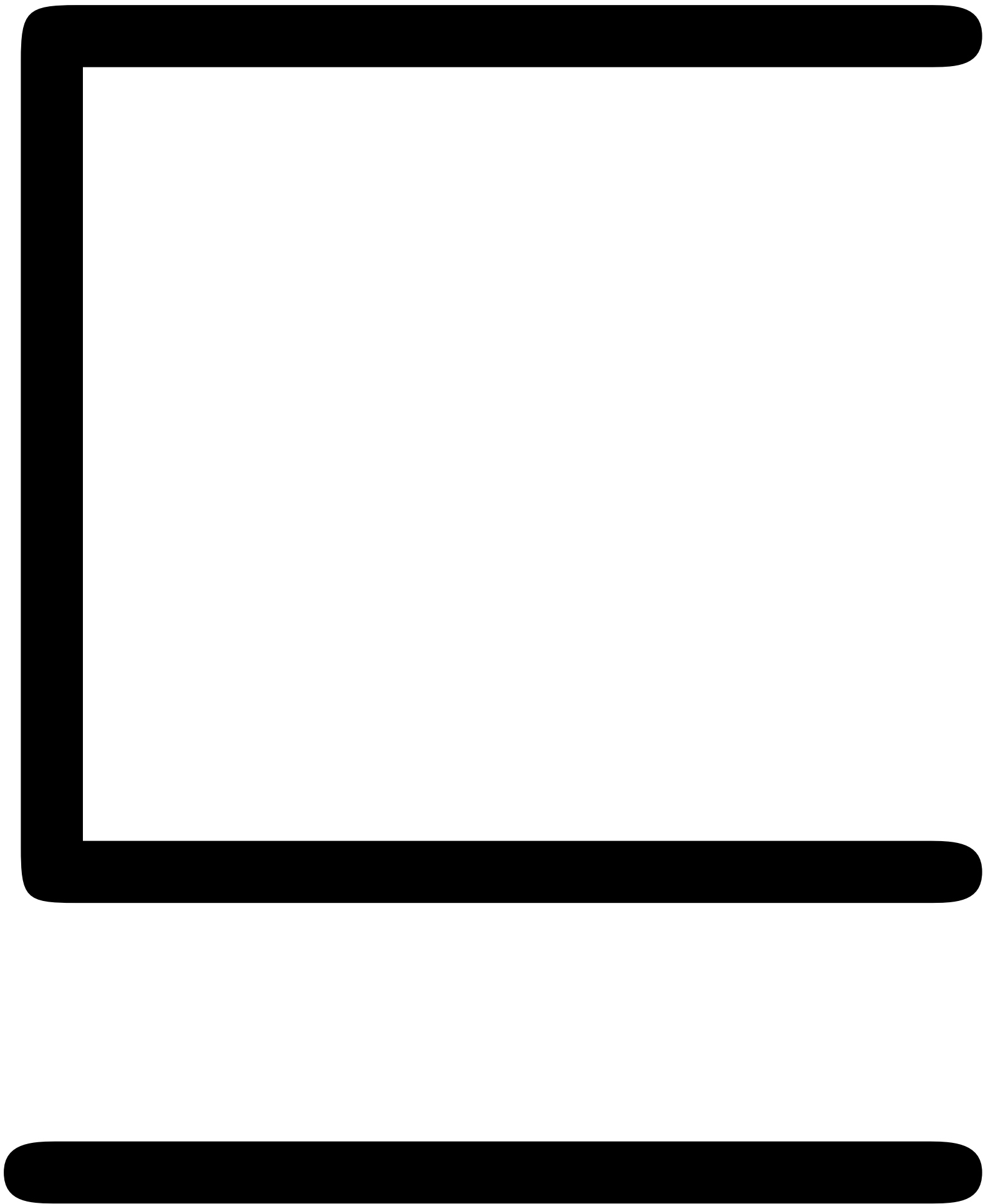
Problem ^o

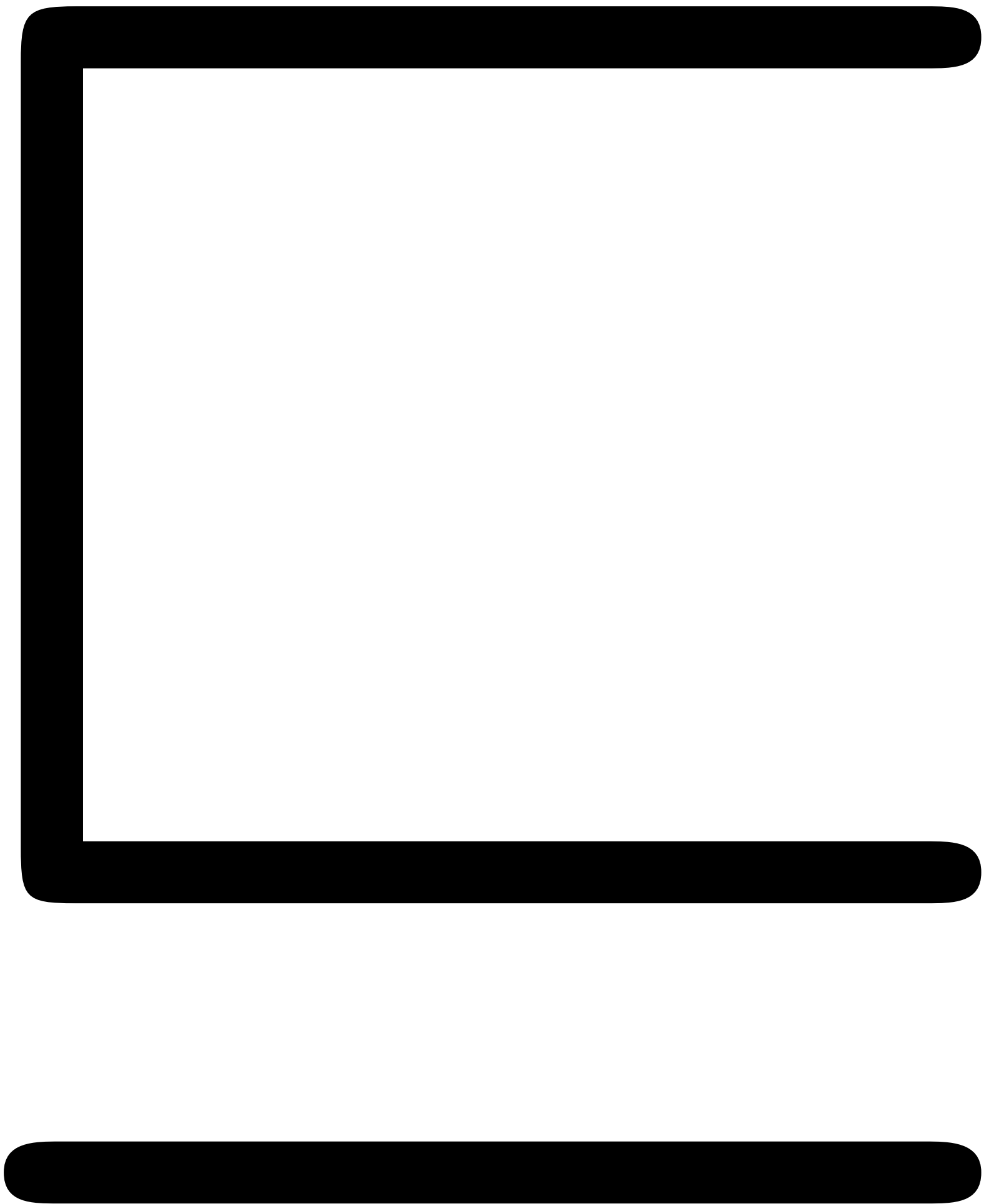




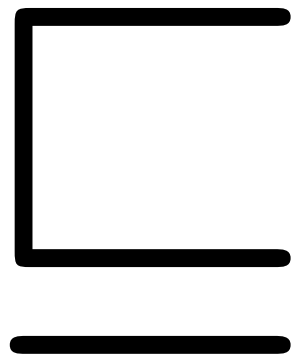








\hat{S}_1



\hat{S}_2

$$\hat{S}_1 = (e, \hat{\rho}, \hat{\sigma}, \hat{k})$$

expression

store

$$\hat{S}_1 = (e, \hat{\rho}, \hat{\sigma}, \hat{k})$$

environment

stack

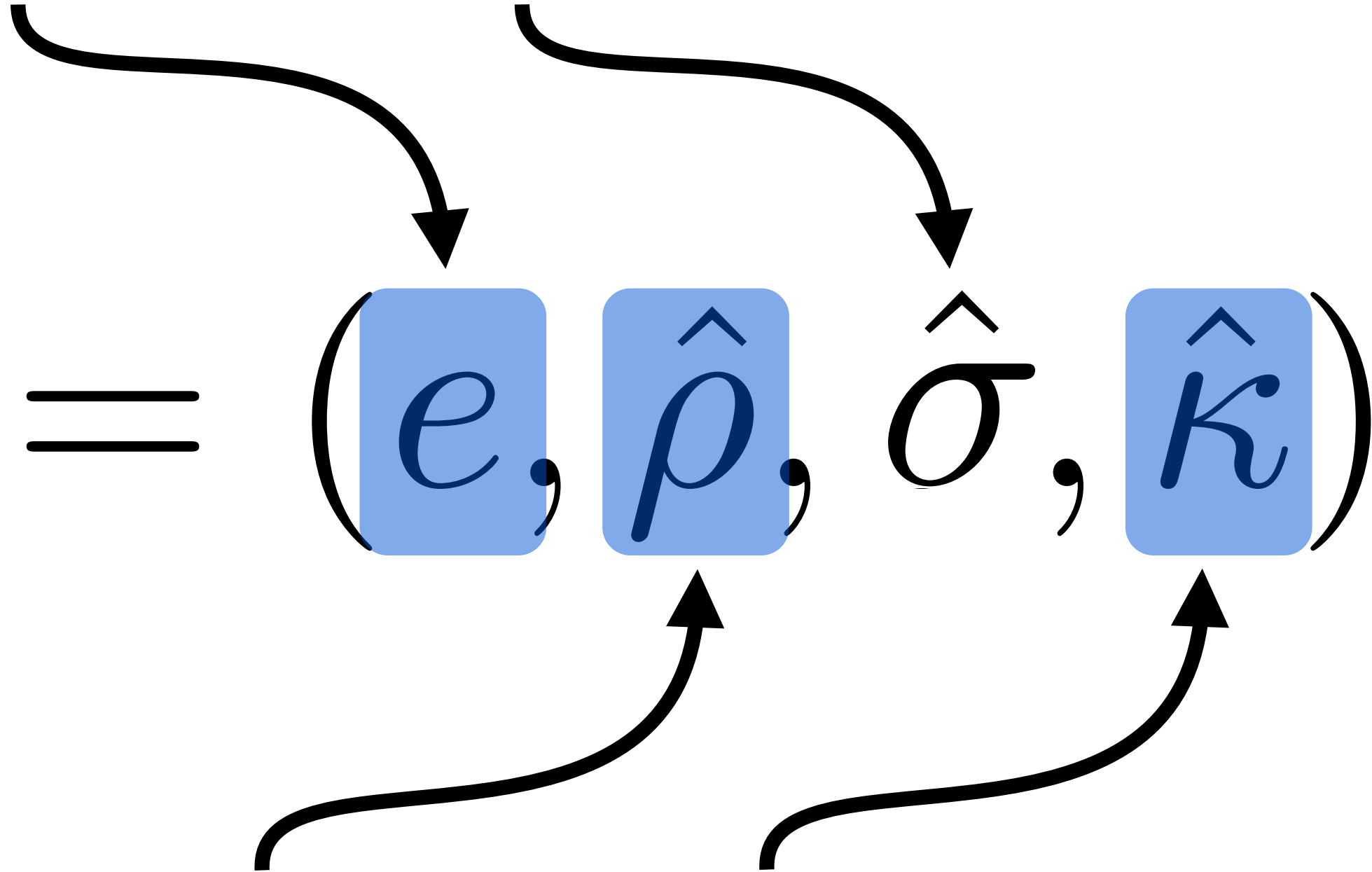
expression

store

$$\hat{S}_1 = (e, \hat{\rho}, \hat{\sigma}, \hat{K})$$

environment

stack



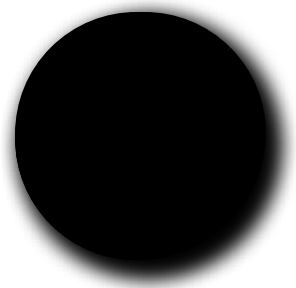
$$\hat{\sigma} : \widehat{Addr} \rightarrow \mathcal{P}(\widehat{Value})$$

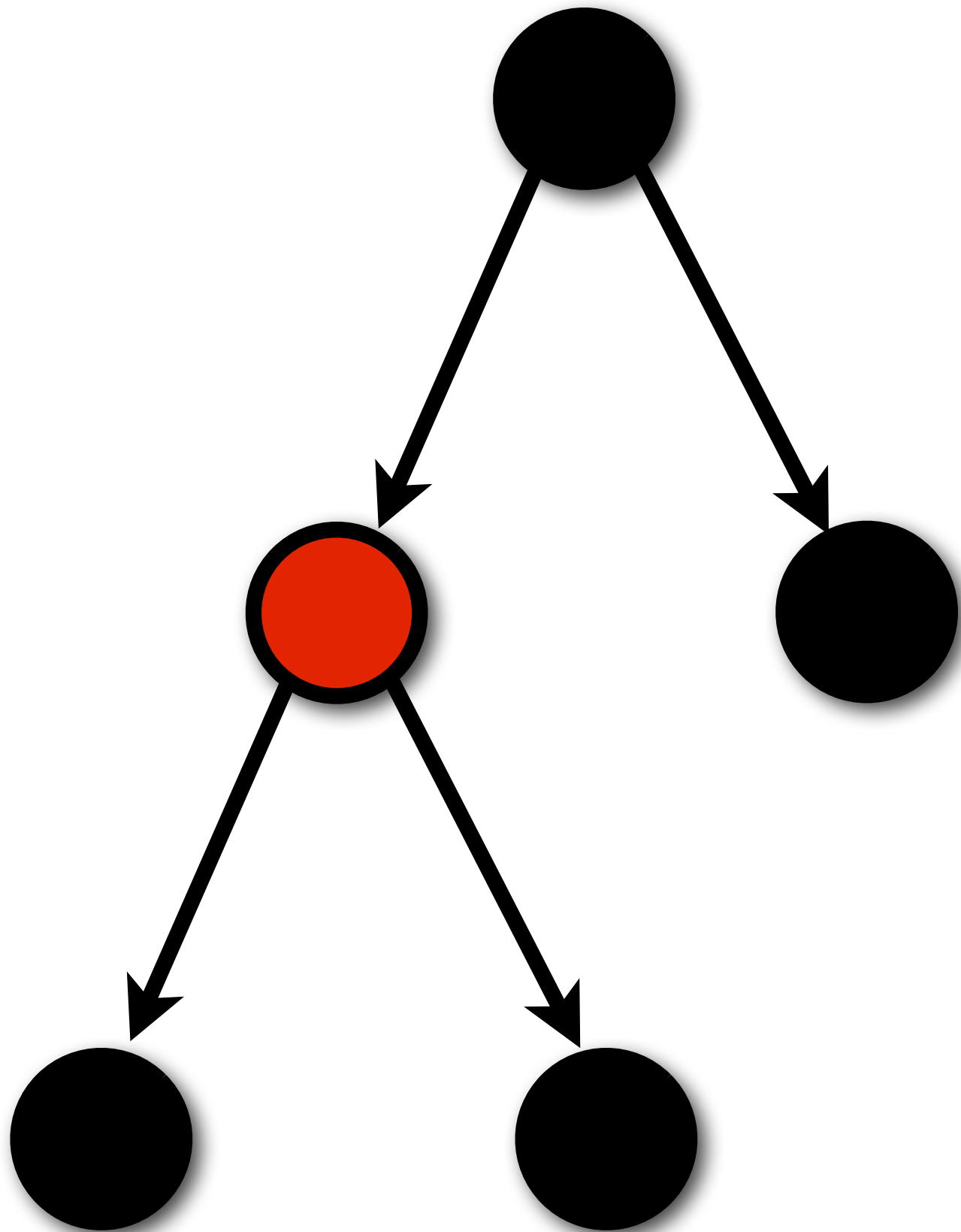
Addr

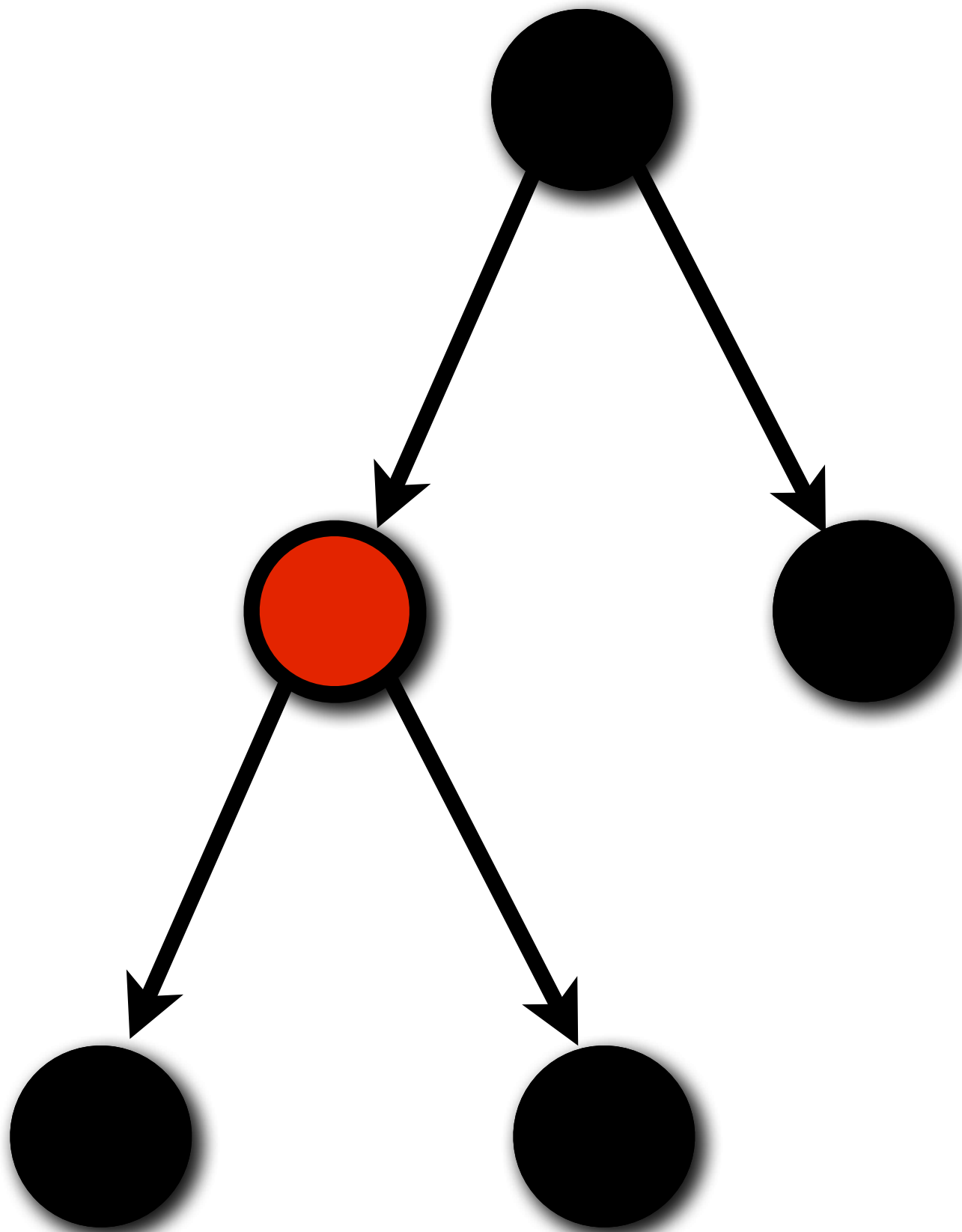
0	1	0	0	0	1	0	1
0	0	0	0	1	1	0	0
0	0	0	1	0	1	0	0
0	1	0	0	0	0	1	0
1	0	0	0	0	0	0	0
0	0	0	1	0	0	1	0
0	0	0	0	0	1	1	1
0	0	1	0	0	0	1	0

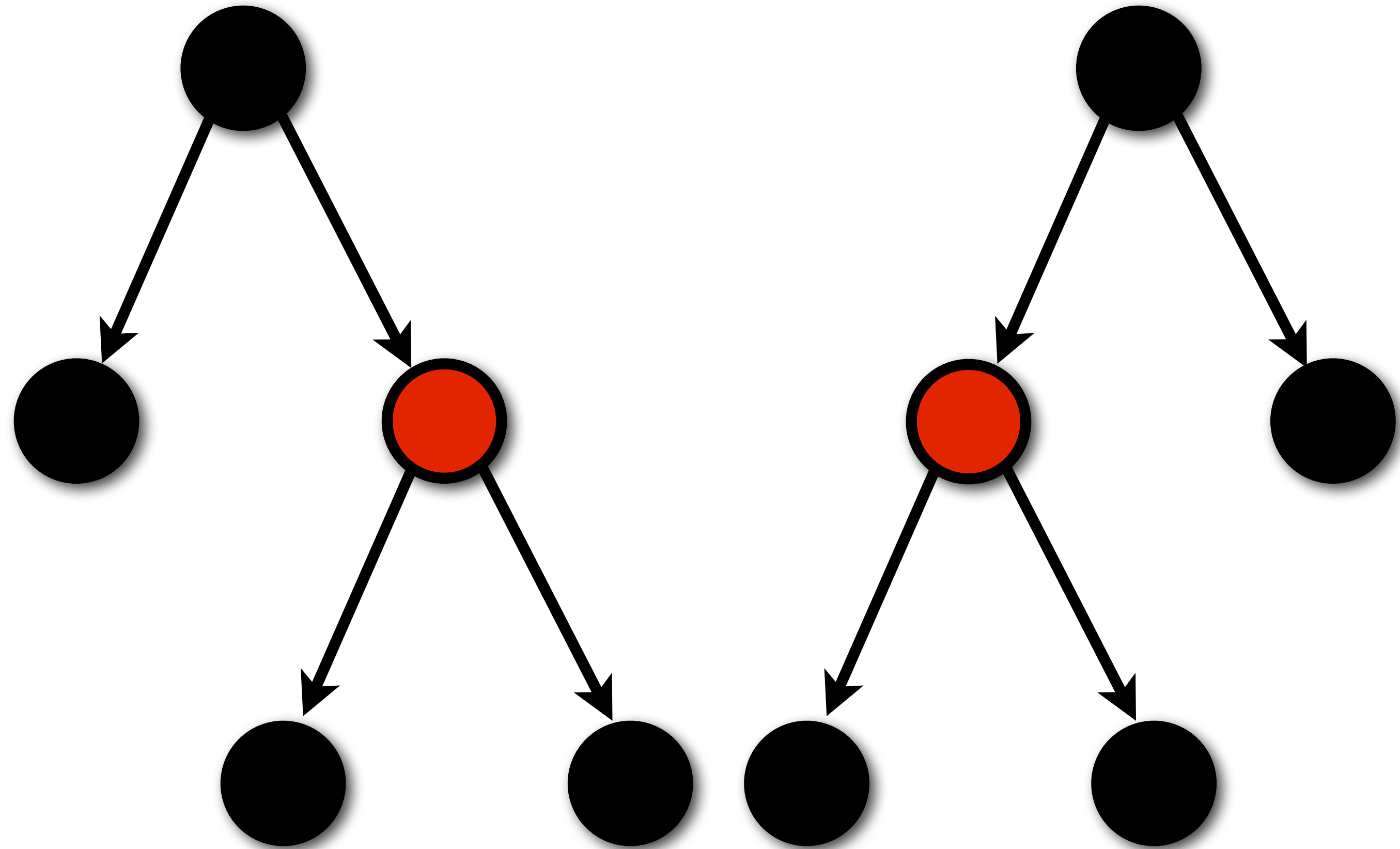
Value

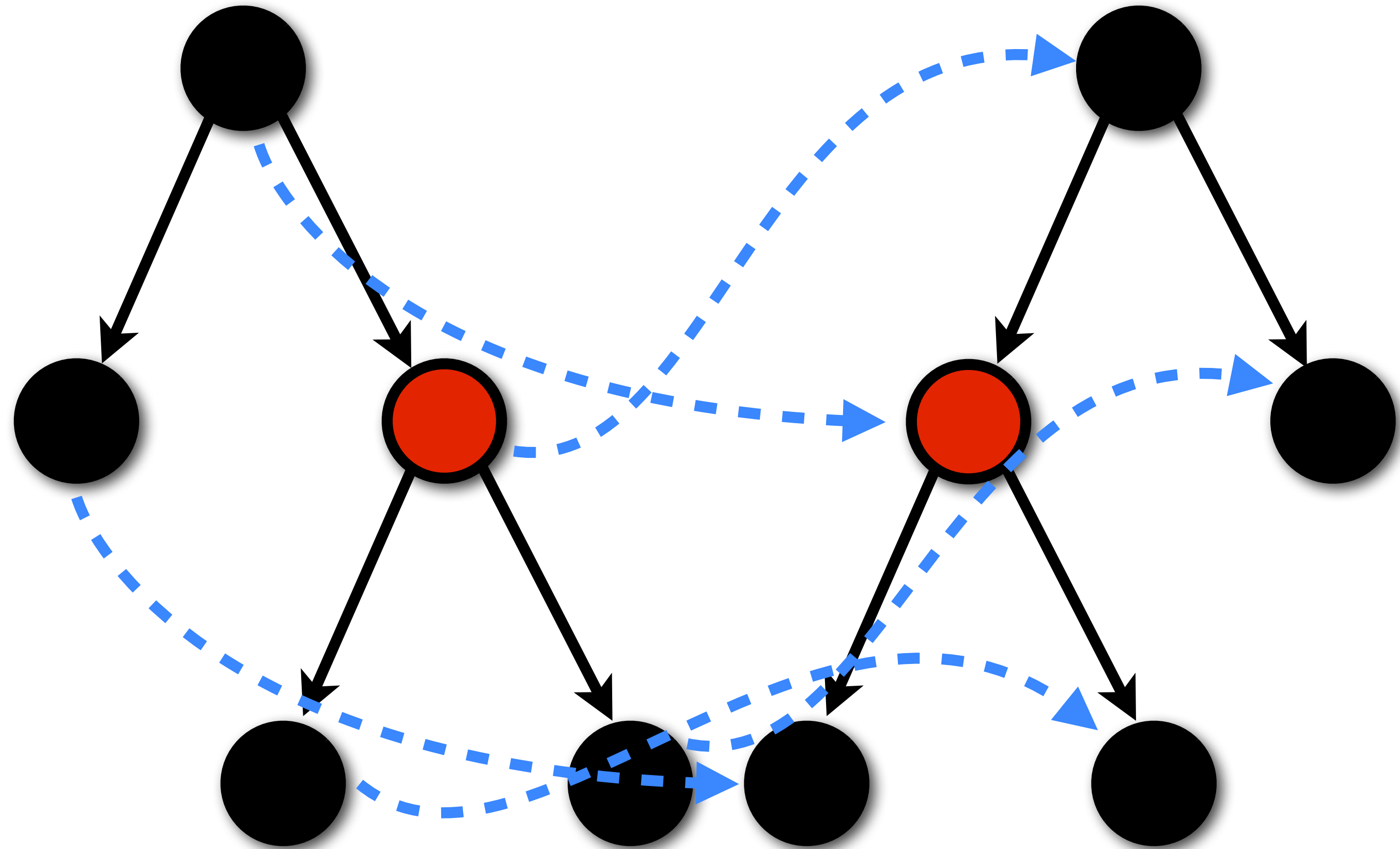
0	0	0	0	0	0	0	1	0	0	0	0	0	0	0	1	0	0	0	0	0	0	0	1	0	0	0	0	0	0	0	1	0	0			
0	0	1	0	0	1	0	0	0	0	1	0	0	1	0	0	0	0	1	0	0	1	0	0	0	0	1	0	0	1	0	0	0	0			
0	0	0	0	1	1	1	0	0	0	0	0	1	1	1	0	0	0	0	0	1	1	1	0	0	0	0	0	1	1	1	0	0	0			
0	1	0	0	0	1	0	0	0	1	0	0	0	1	0	0	0	1	0	0	0	1	0	0	0	1	0	0	0	1	0	0	0	1			
1	0	0	0	1	0	1	0	1	0	0	0	1	0	1	0	1	0	0	0	1	0	1	0	1	0	0	0	1	0	1	0	1	0			
0	0	0	1	1	0	0	0	0	0	0	1	1	0	0	0	0	0	0	1	1	0	0	0	0	0	0	0	1	1	0	0	0	0	0		
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0	0	0	0	0	0	0	1	0	0	0	0	0	0	0	1	0	0	0	0	0	0	0	1	0	0	0	0	0	0	0	0	0	1	0	0	
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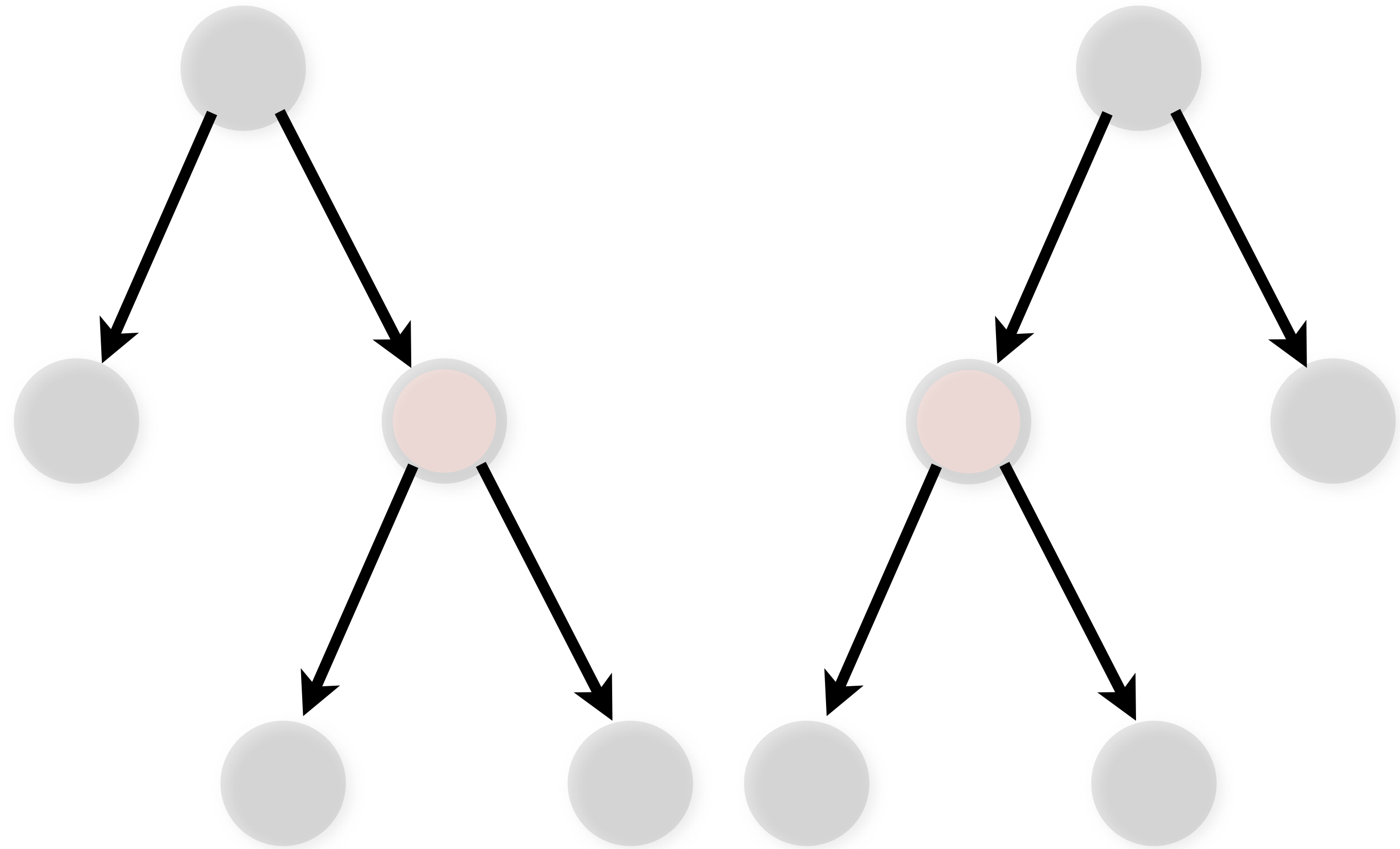


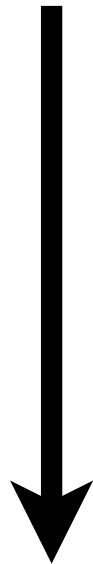
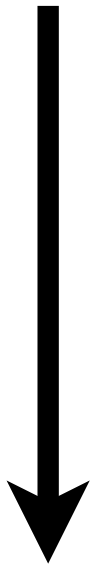
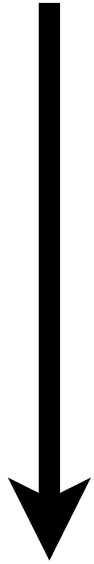
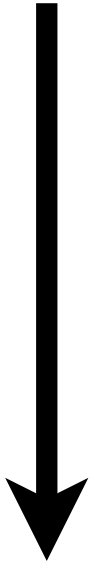
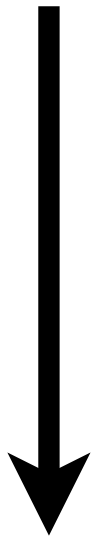


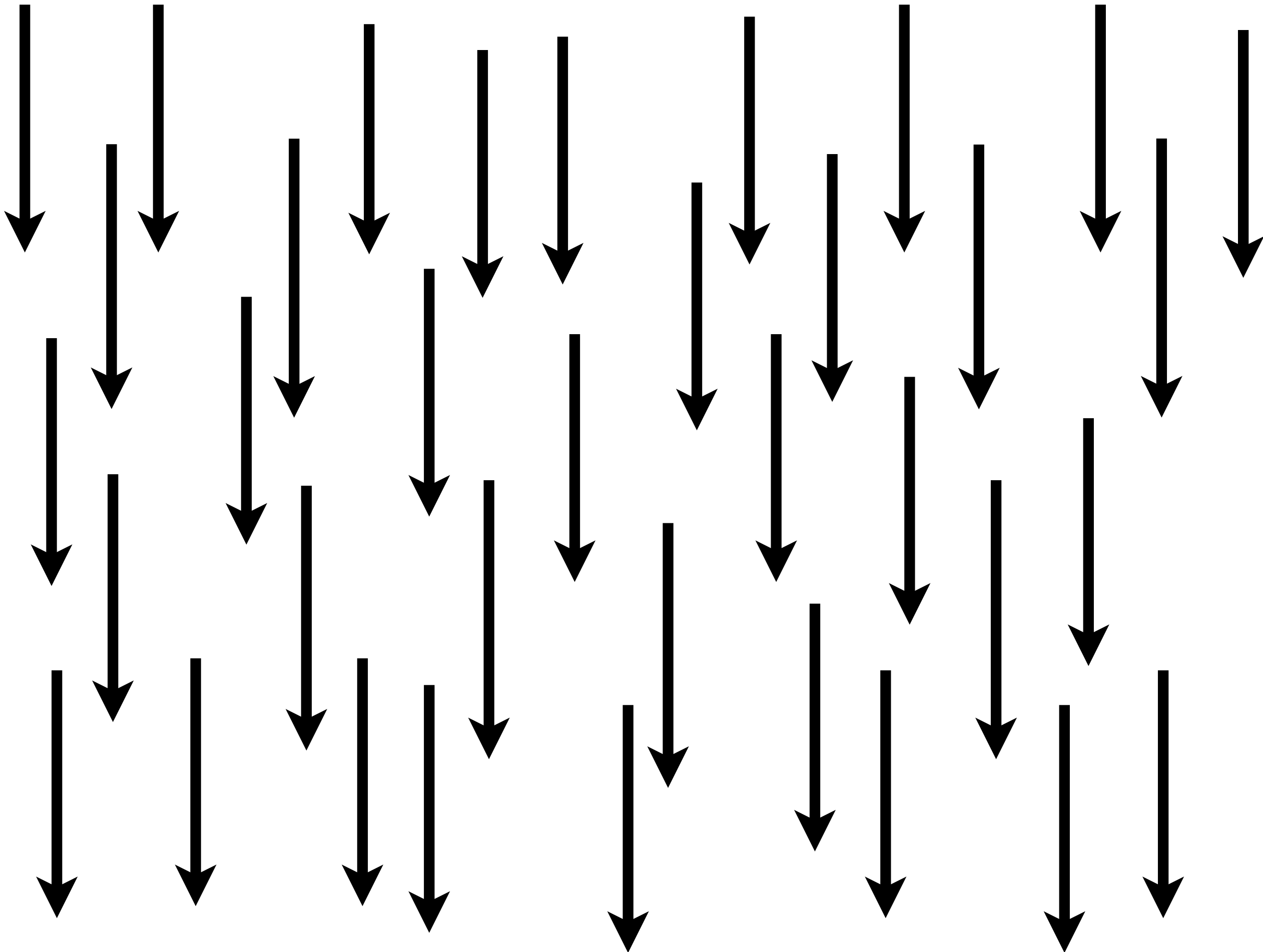


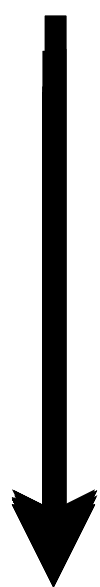








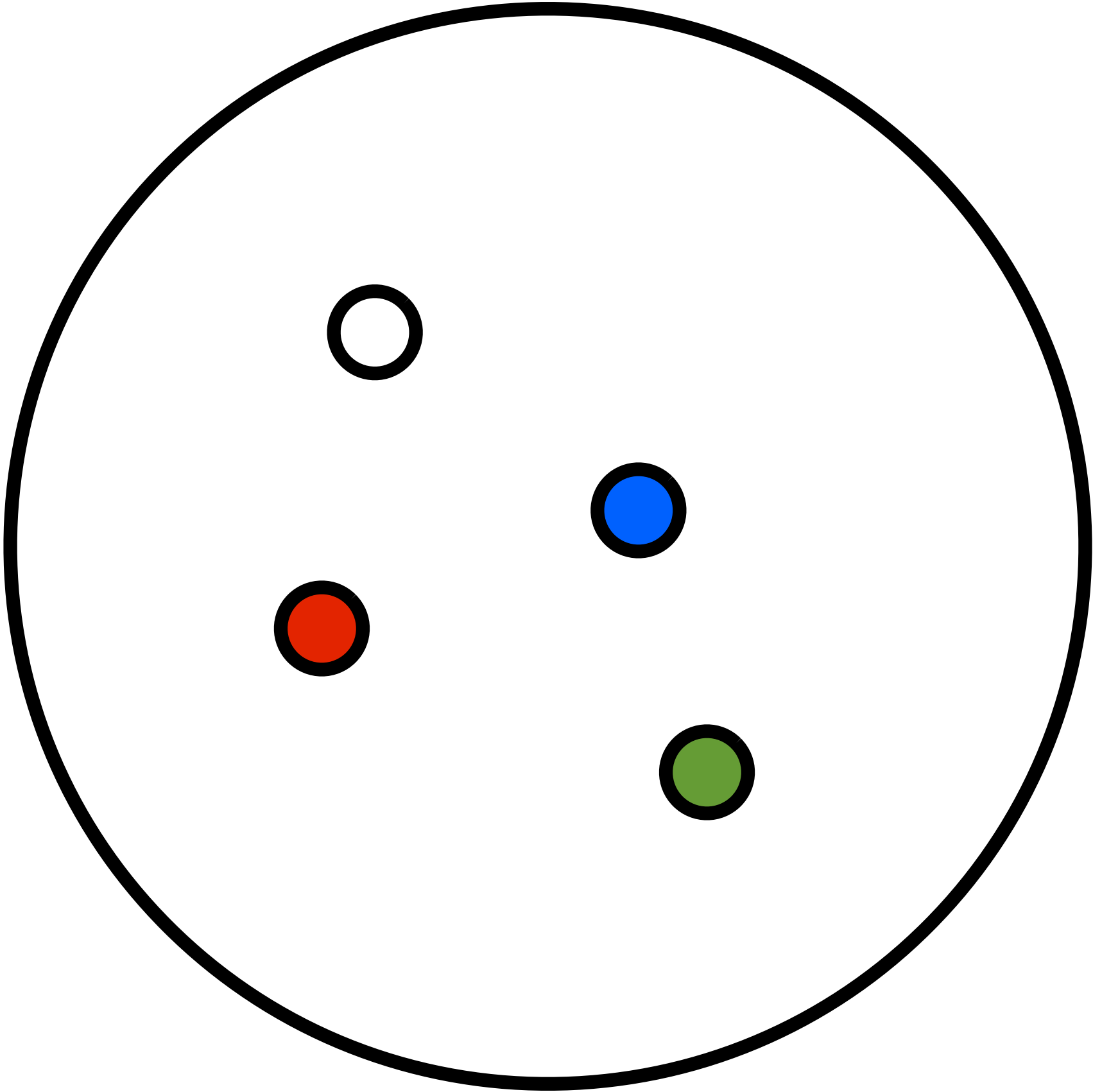


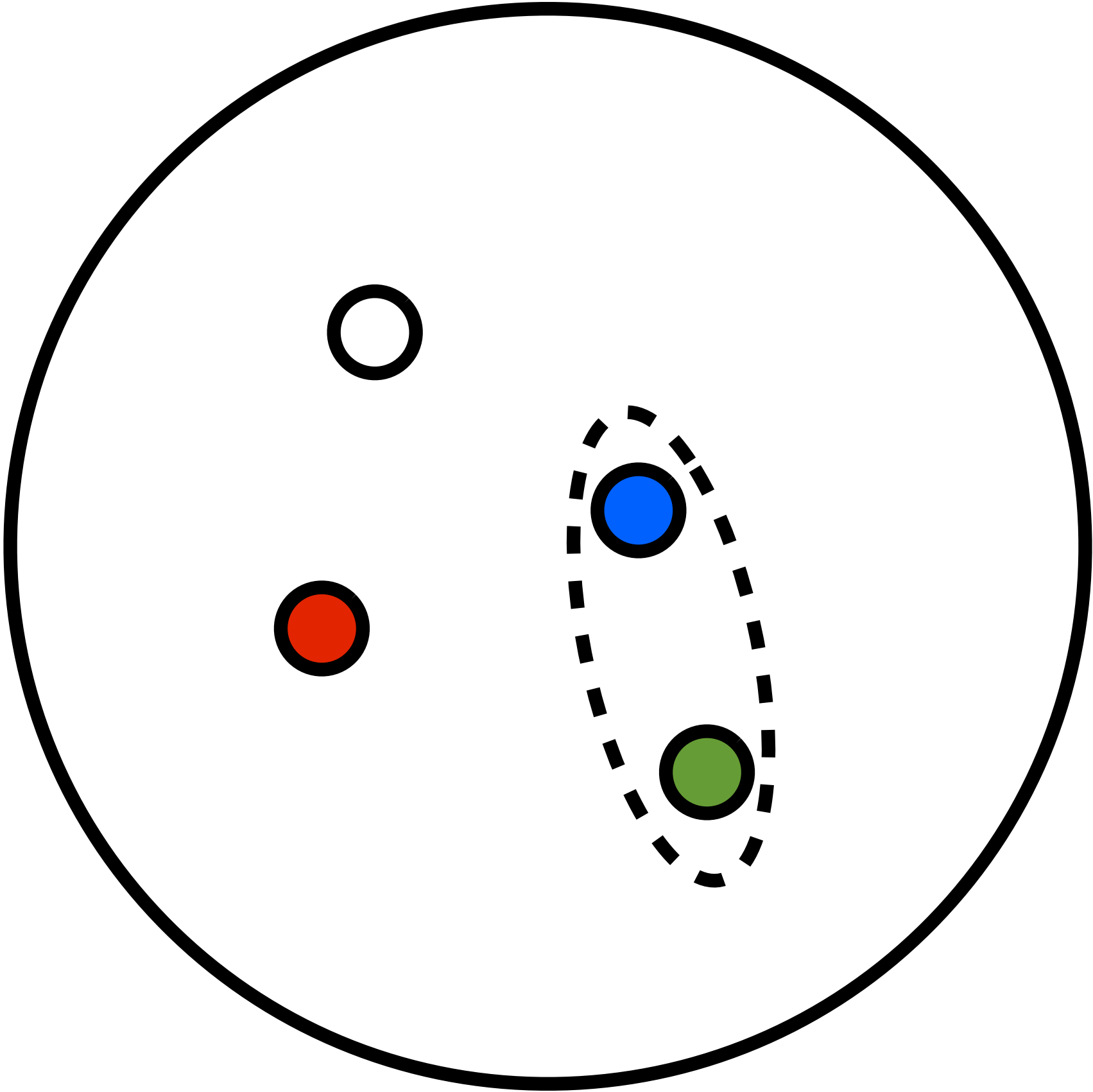




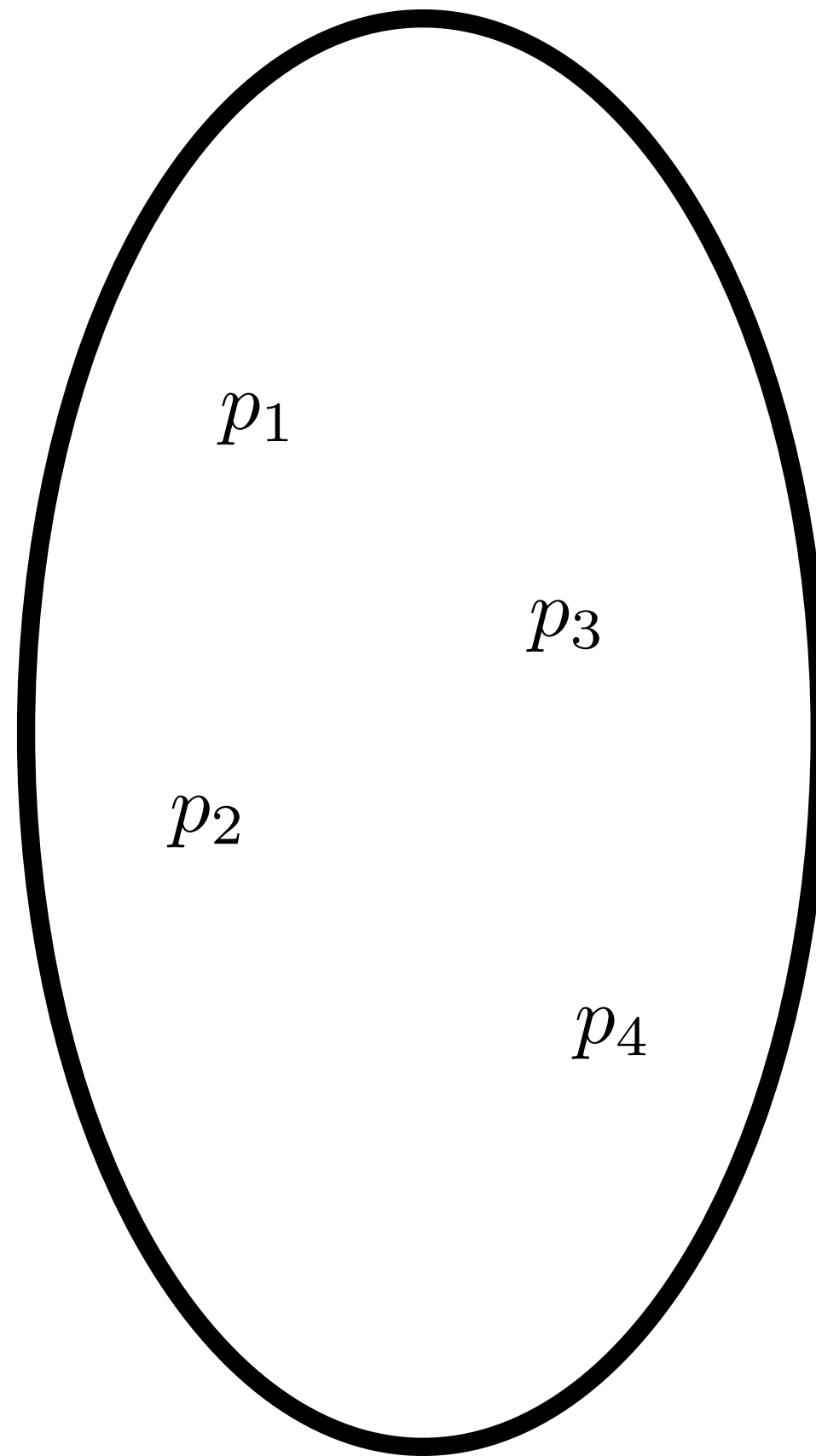
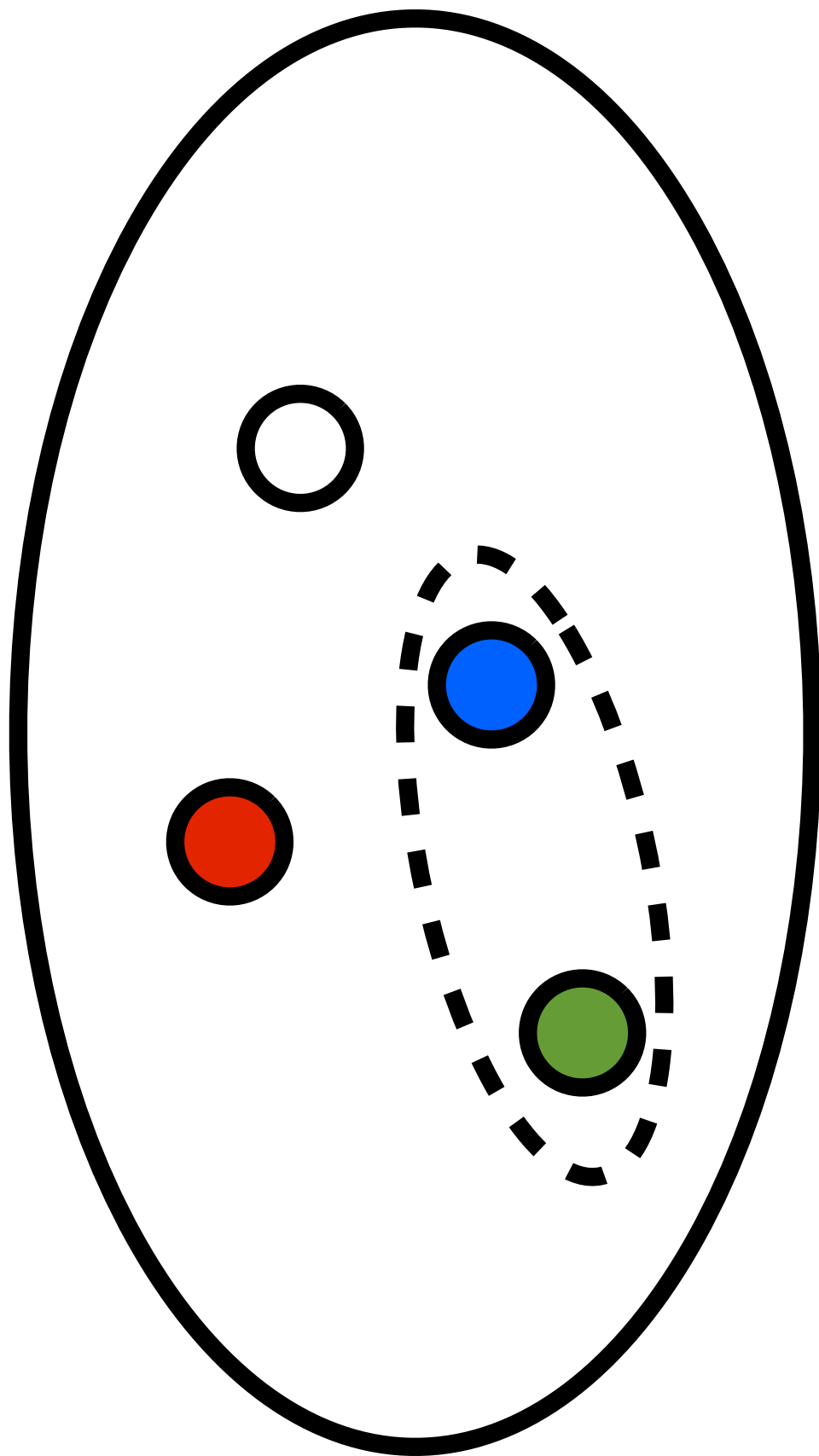
First: Hash sets

Prime decomposition

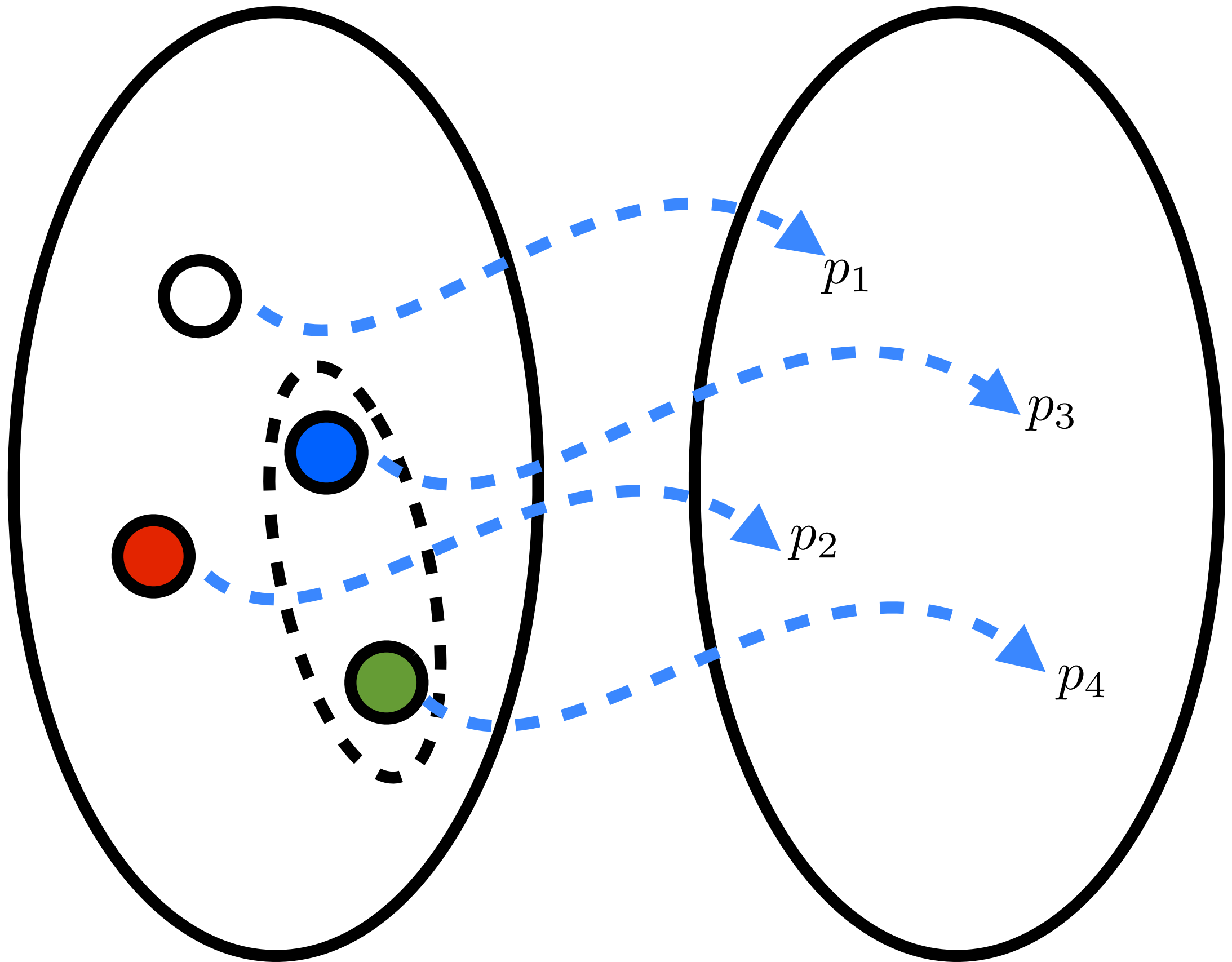


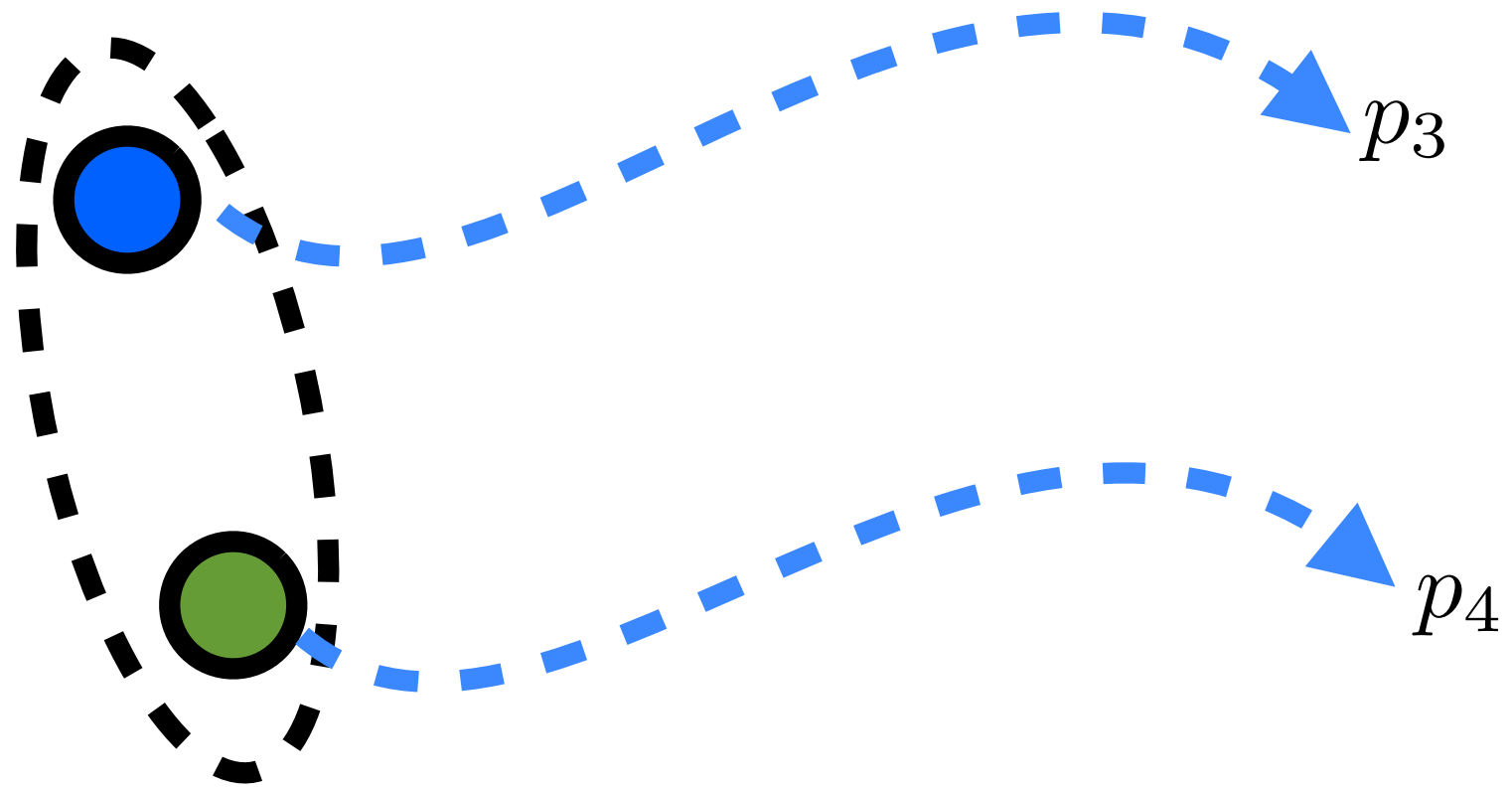


Primes



Primes



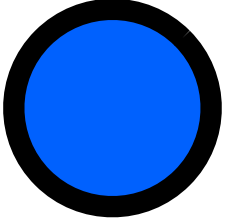
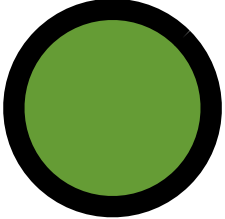


p_3

p_4

$$p_3 \times p_4$$

{  ,  }

{  ,  }

$$A \subseteq B$$

$$[B] \bmod [A] = 0$$

A \cap *B*

$\text{gcd}(\llbracket A \rrbracket, \llbracket B \rrbracket)$

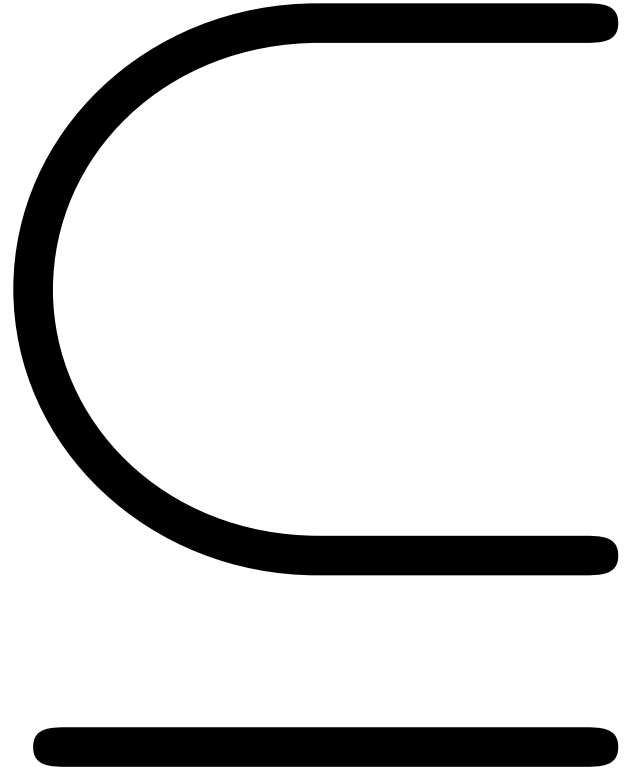
$$\text{lcm}([A], [B])$$

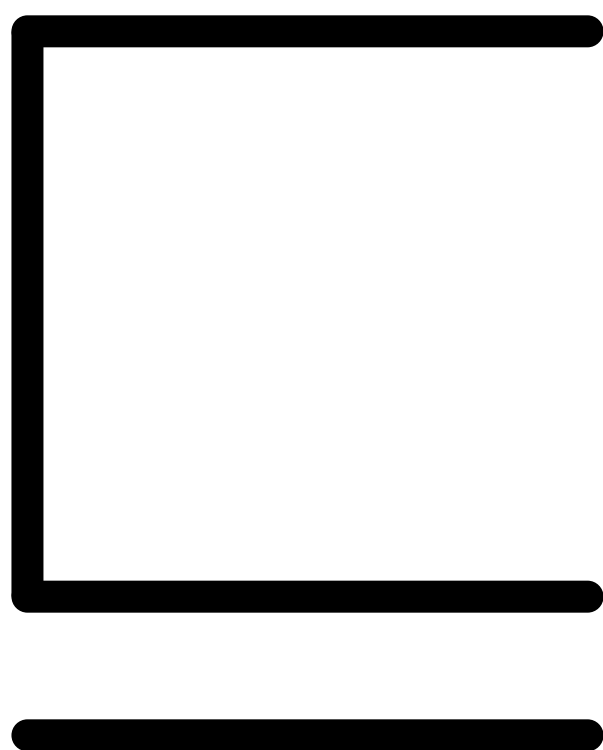
A \cup *B*

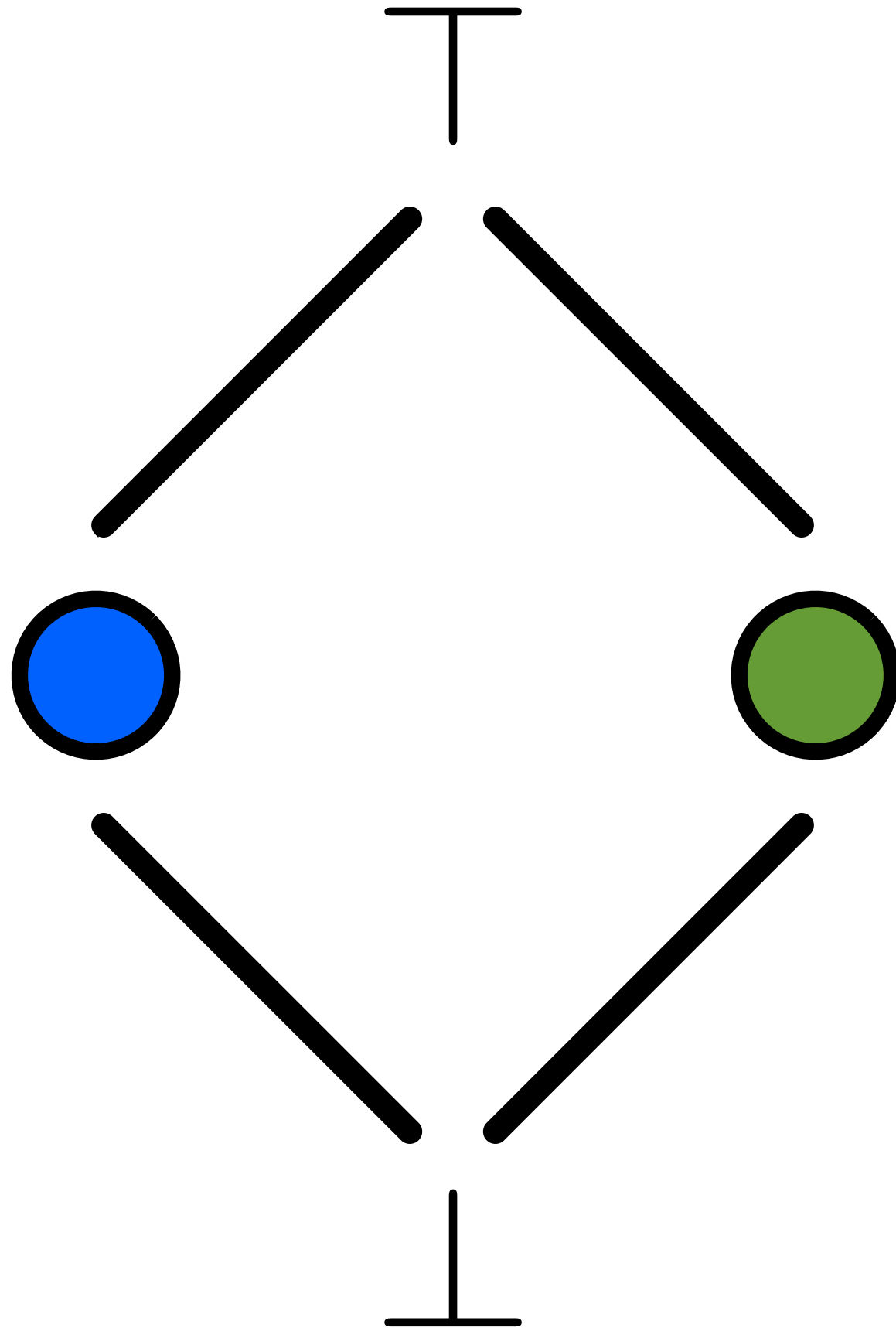
A \cup *B*

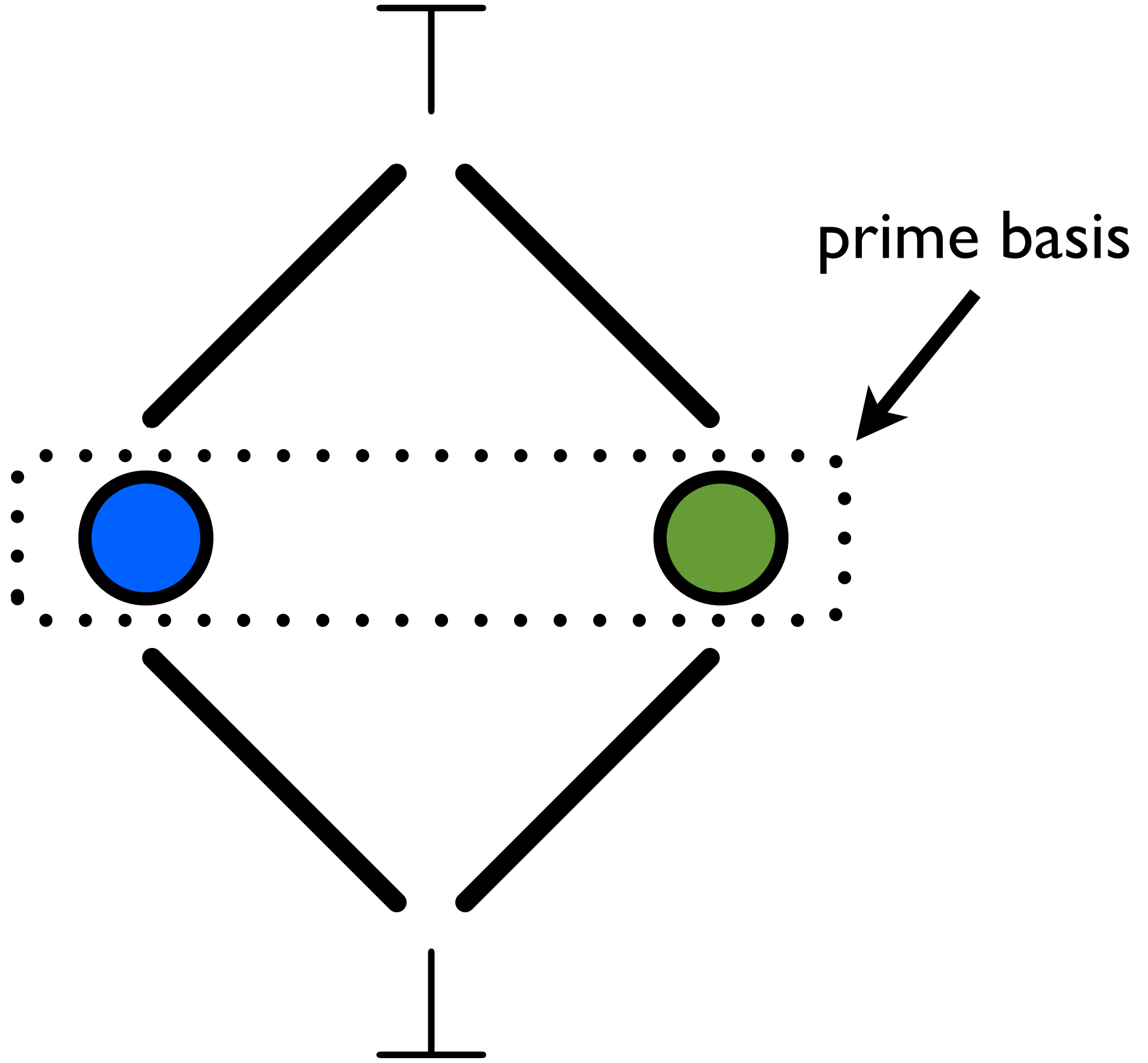
A — *B*

$$\llbracket A \rrbracket / \gcd(\llbracket A \rrbracket, \llbracket B \rrbracket)$$



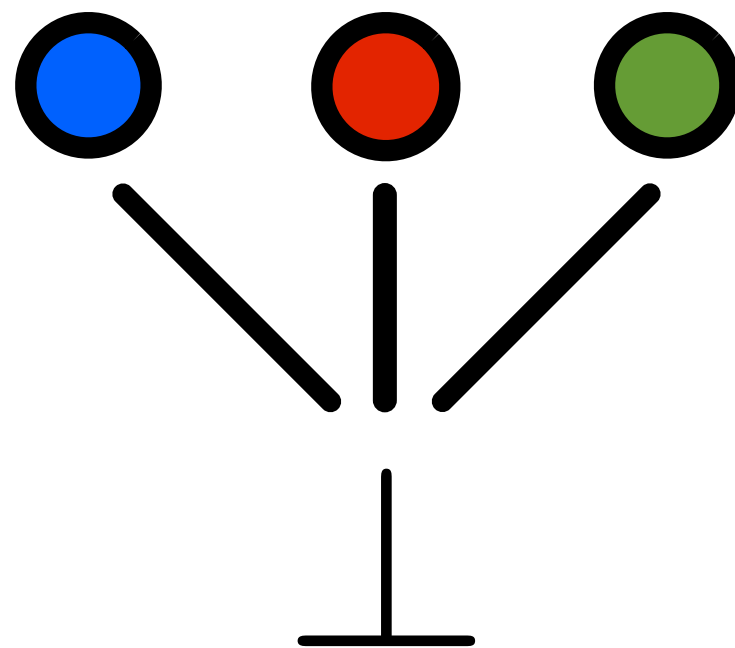


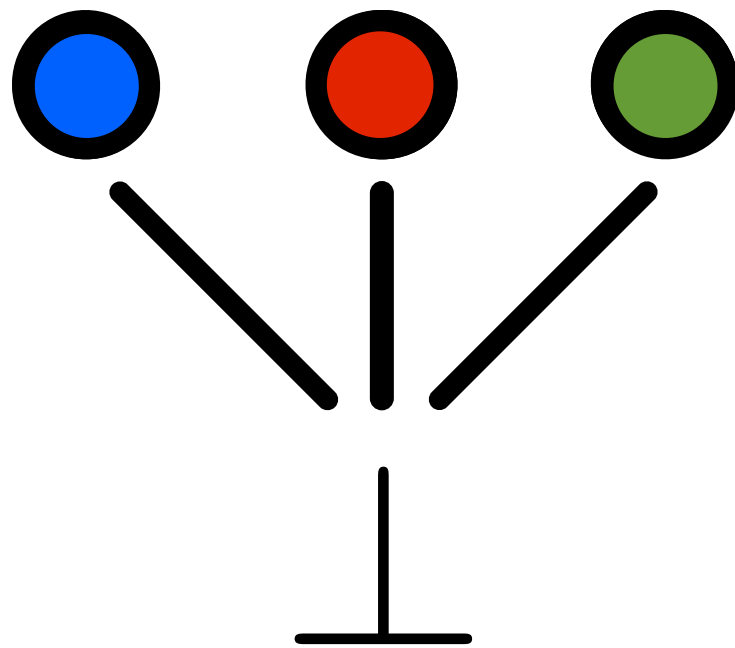


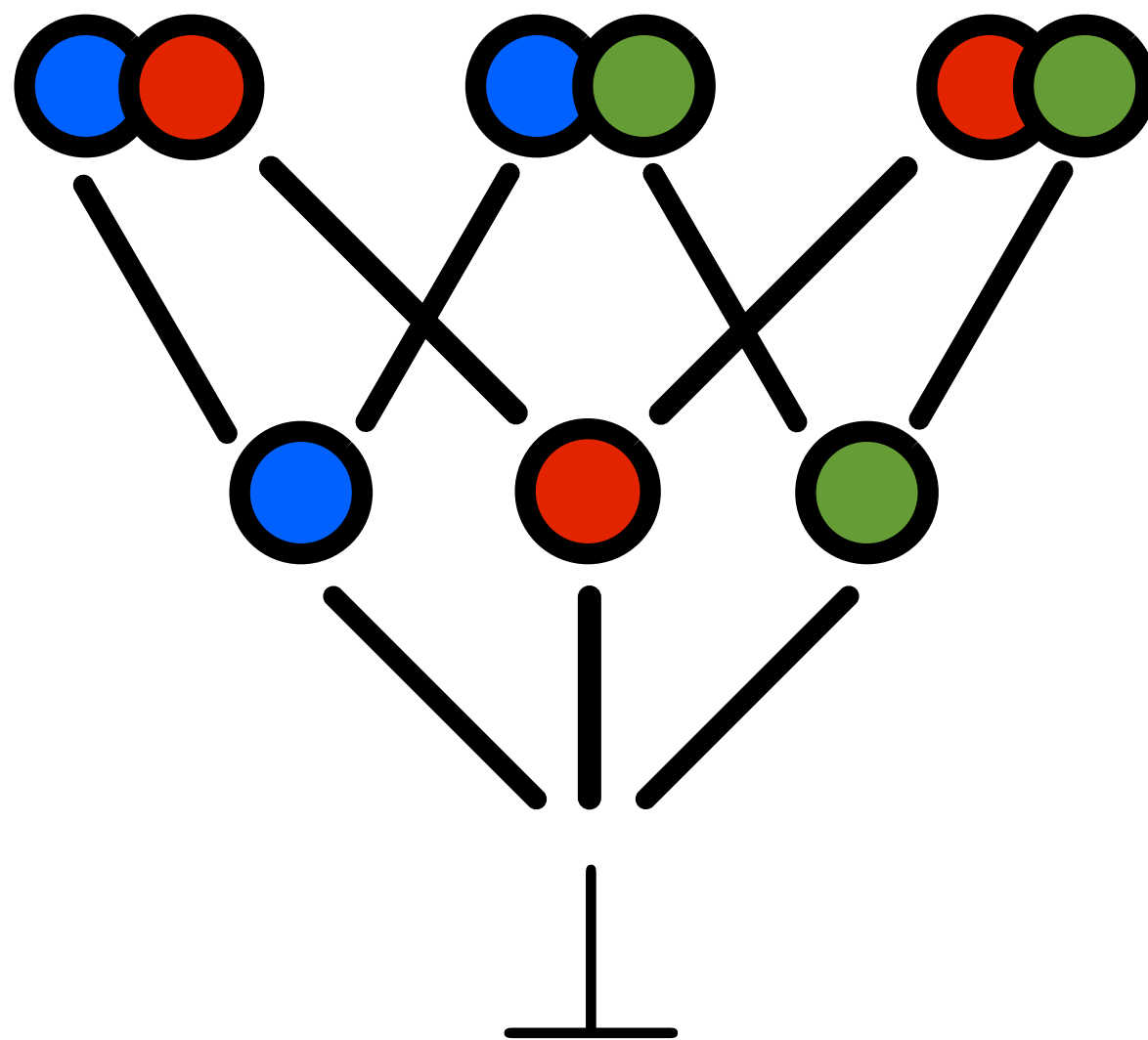


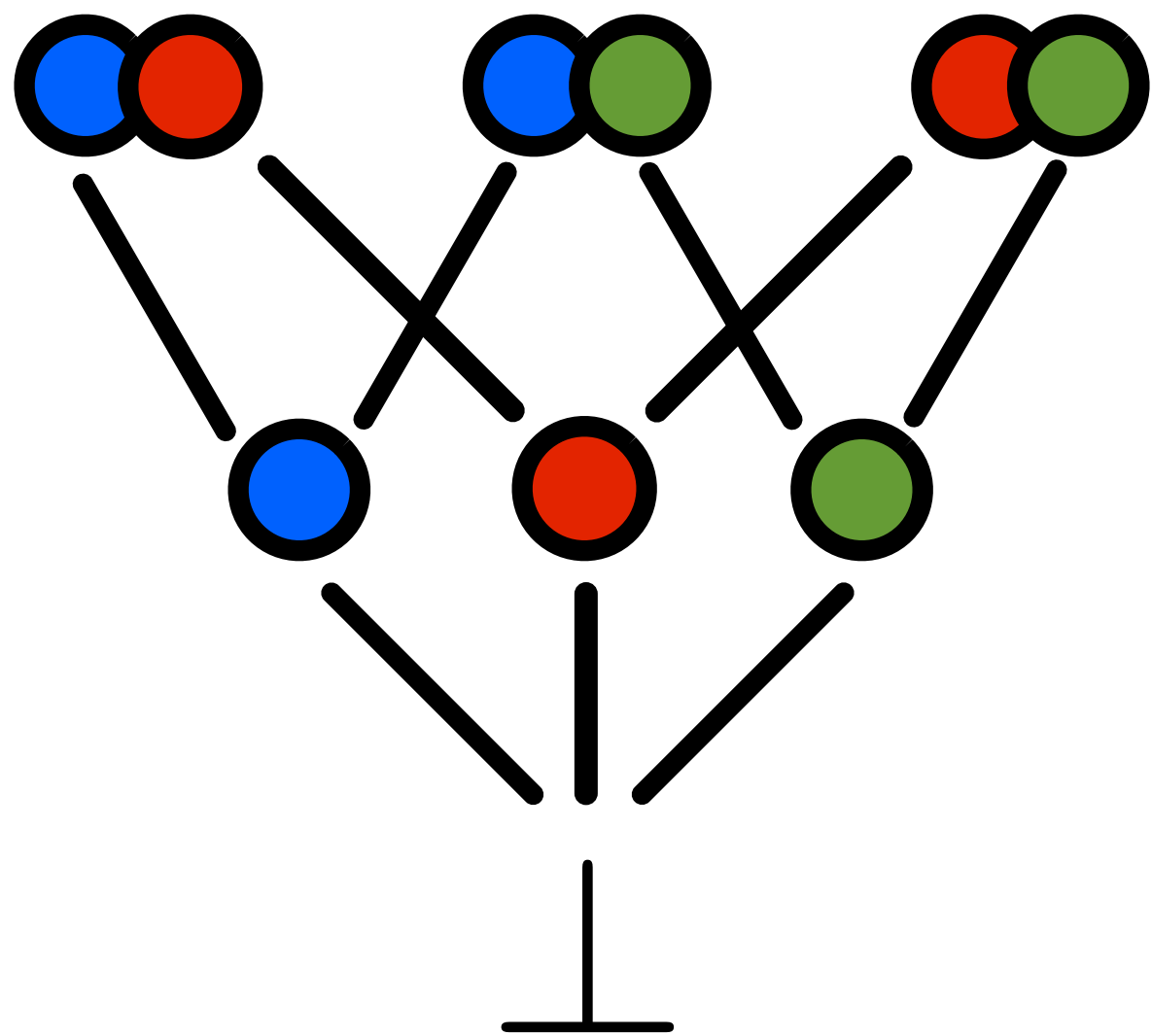
$$n = p_1^{m_1} p_2^{m_2} p_3^{m_3} \dots$$

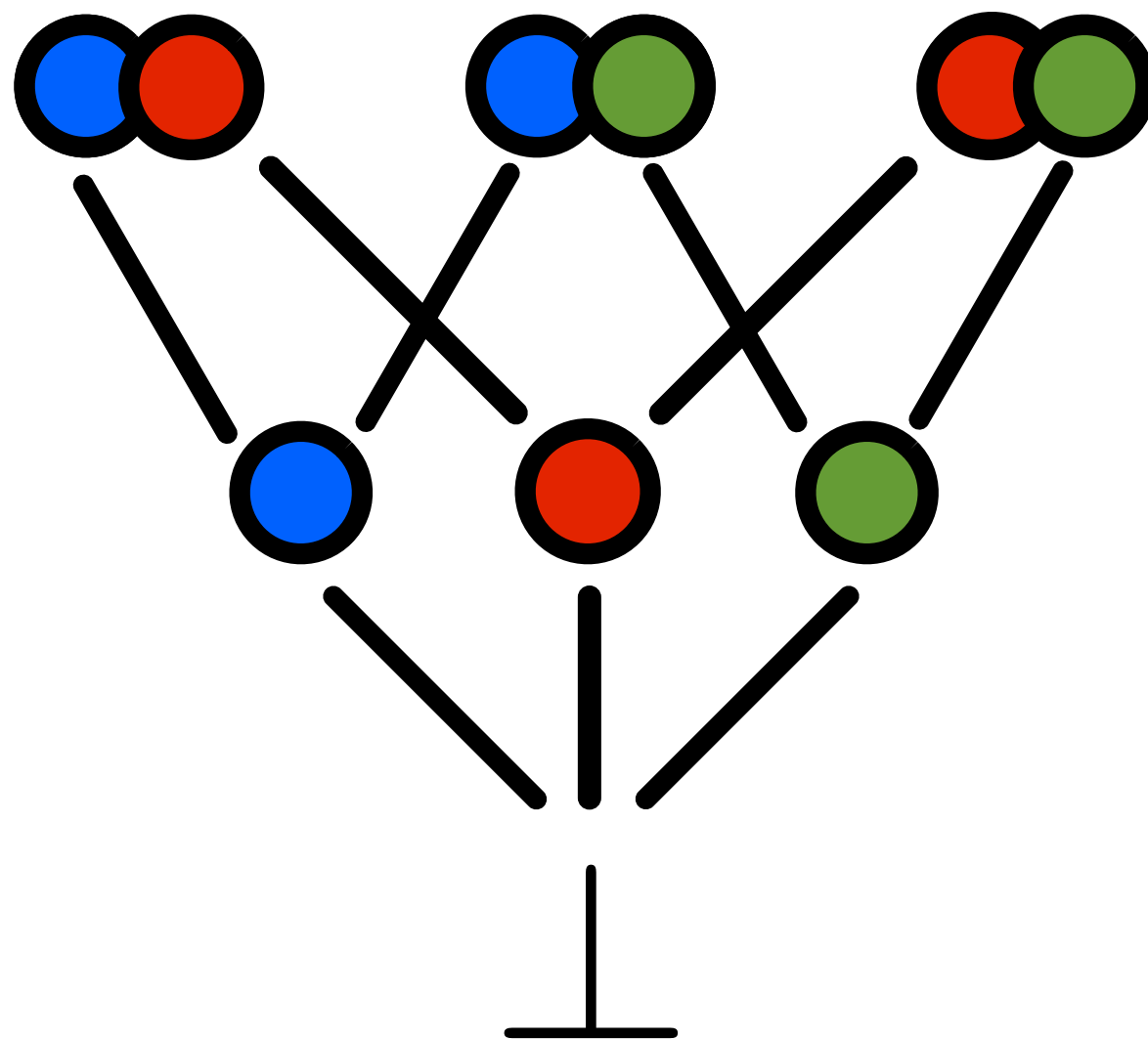
$$n = \sqcup \{ \text{green circle}, \text{blue circle}, \text{red circle} \}$$

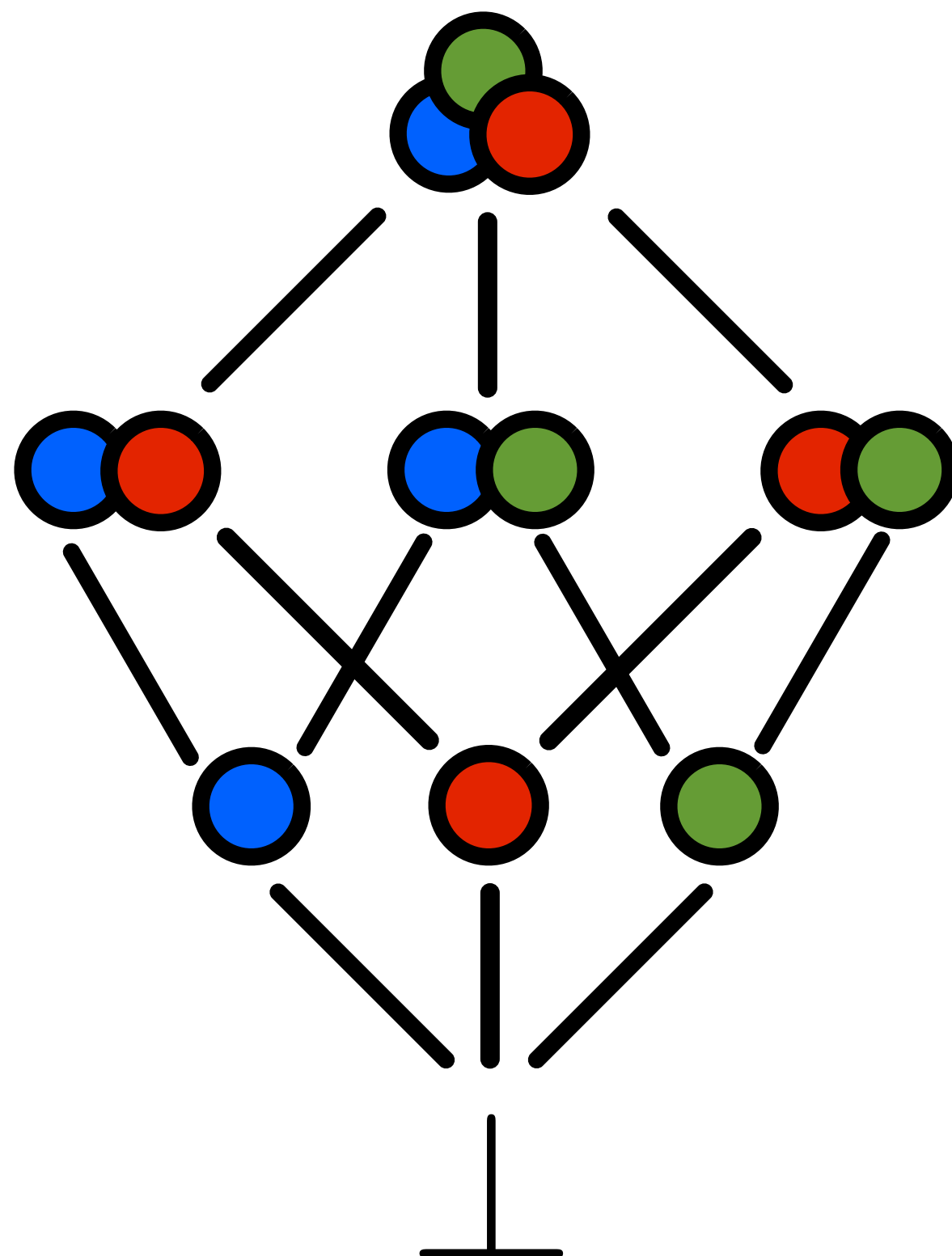


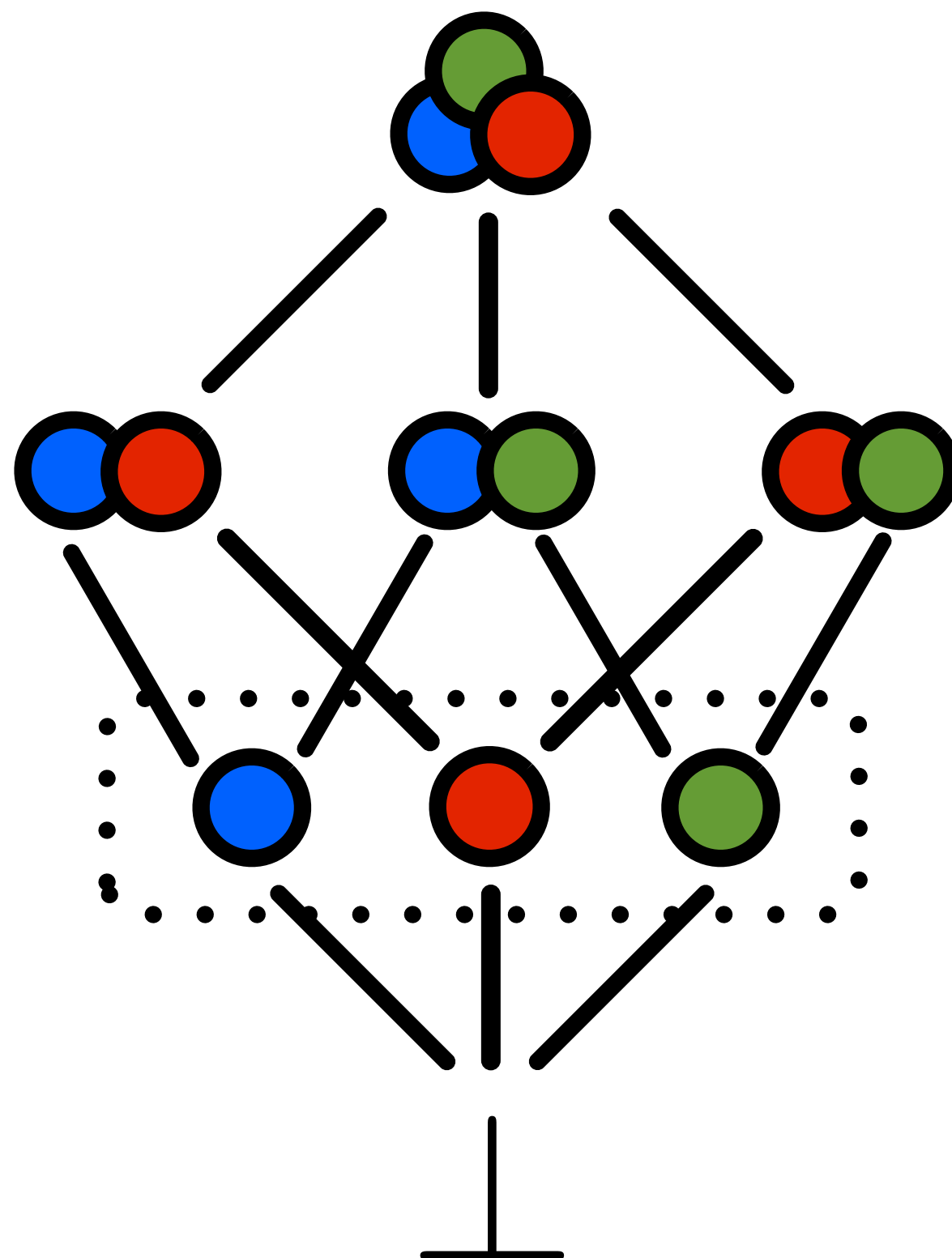


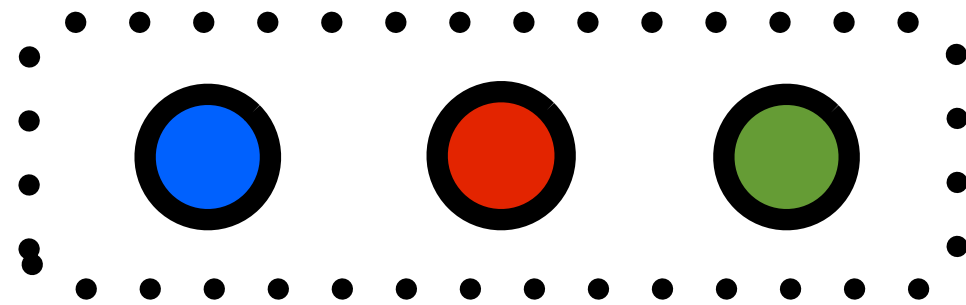


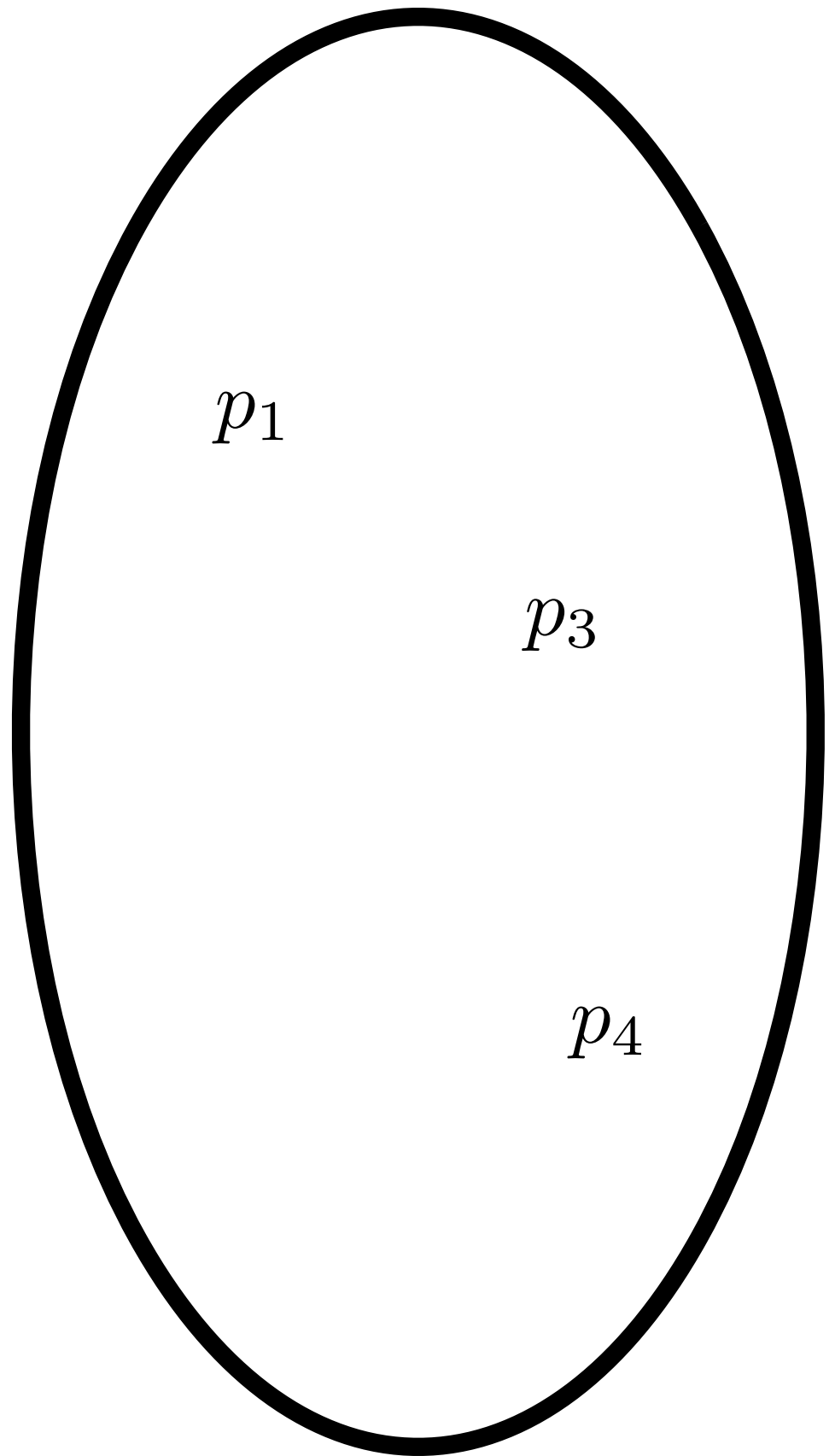
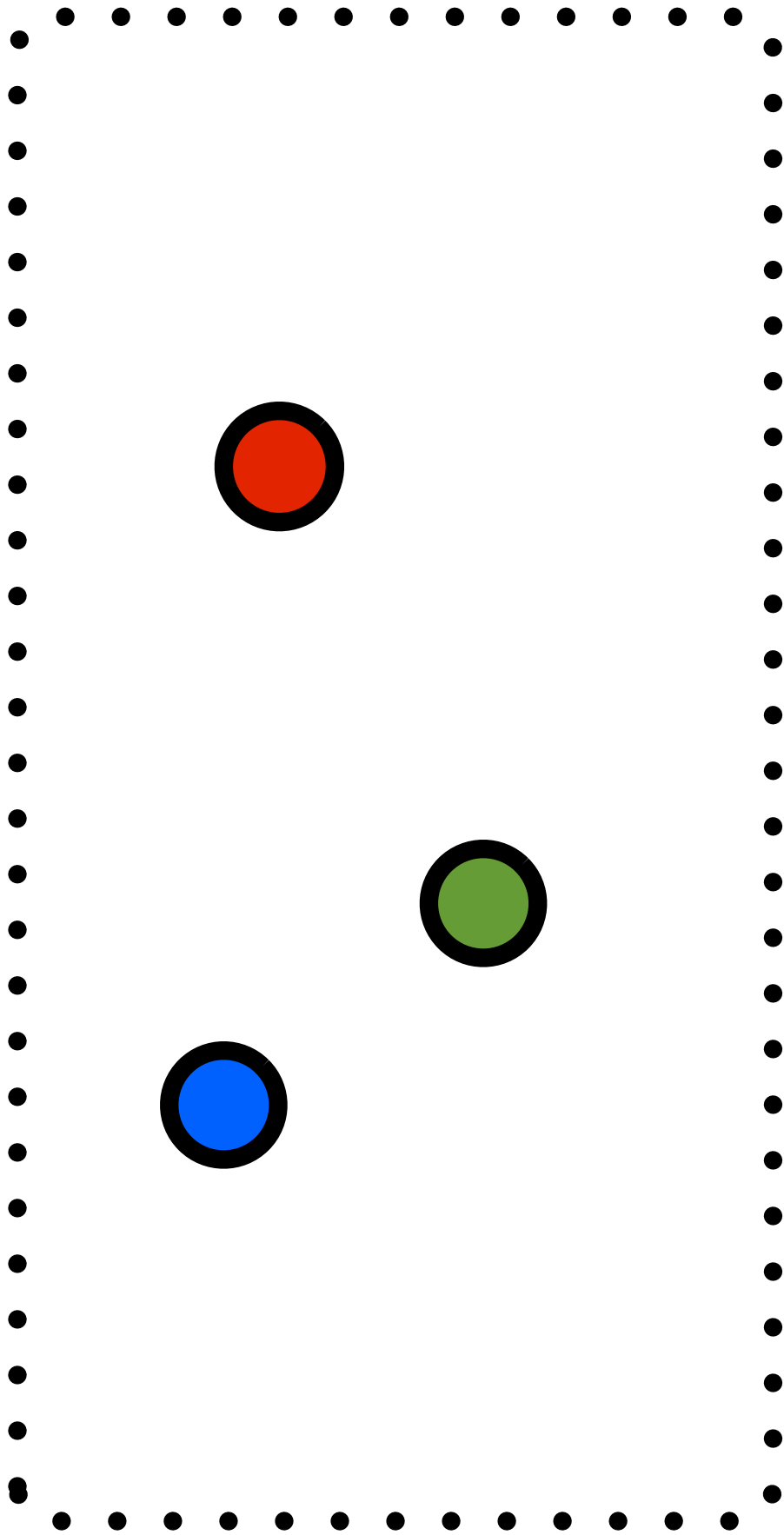


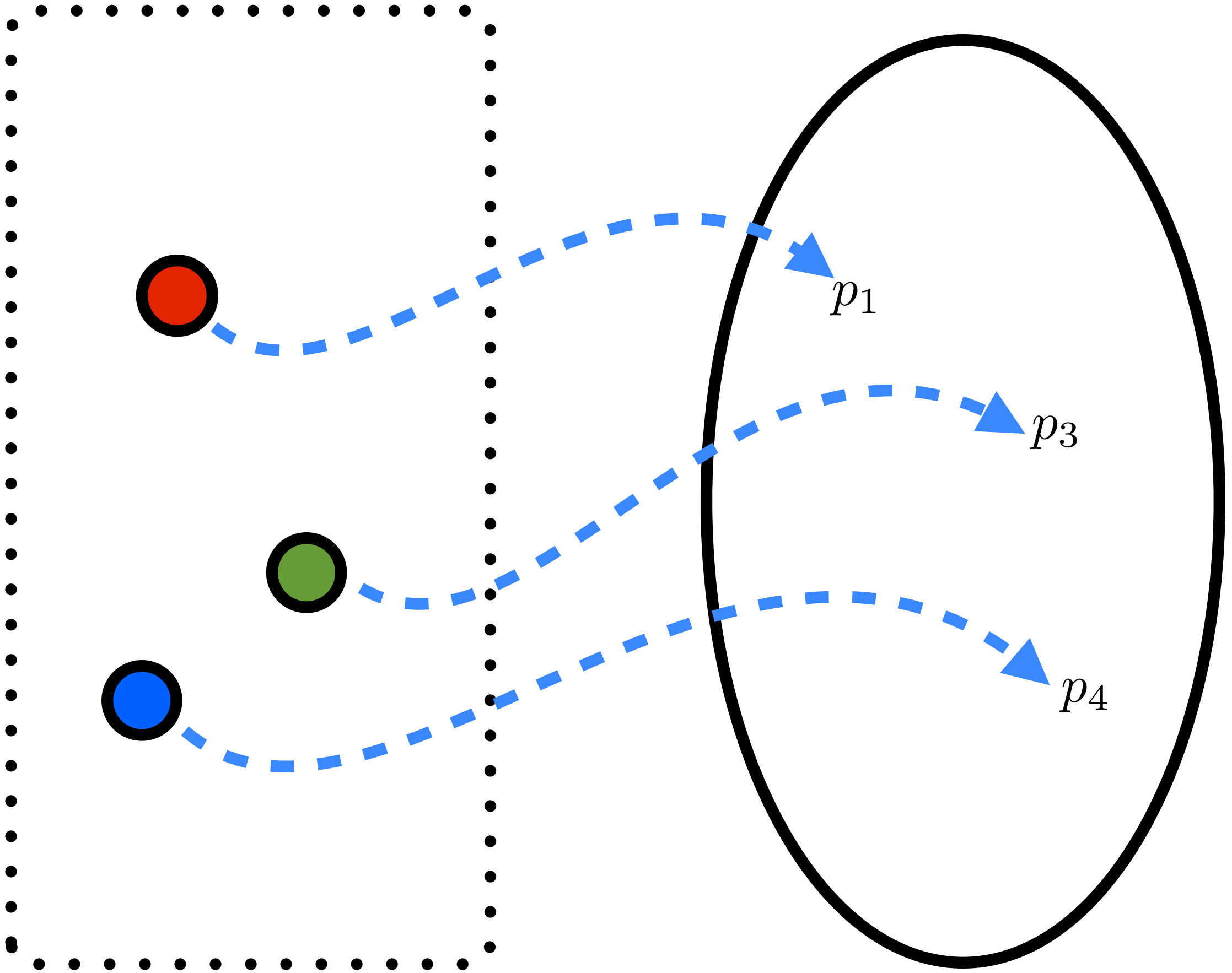












$$x \sqcup y$$

$$\text{lcm}(\llbracket x \rrbracket, \llbracket y \rrbracket)$$

$x \sqsubseteq y$

$$\llbracket y \rrbracket \bmod \llbracket x \rrbracket = 0$$

x \sqsupseteq *y*

$\text{gcd}(\llbracket x \rrbracket, \llbracket y \rrbracket)$

But, does it work for CFA?

$$\hat{\sigma} : \widehat{Addr} \rightarrow \mathcal{P}(\widehat{Value})$$

\widehat{Addr}

\widehat{Value}

$$\{ \hat{a}_1, \hat{a}_2 \}$$

$$\{ \hat{v}_1, \hat{v}_2 \}$$

$$[\hat{a}_1 \mapsto \{\hat{v}_2\}]$$

$$[\hat{a}_2 \mapsto \{\hat{v}_1\}]$$

$$[\hat{a}_2 \mapsto \{\hat{v}_2\}]$$

$$[\hat{a}_1 \mapsto \{\hat{v}_1\}]$$

L_1 has a prime basis.

L_2 has a prime basis.

$L_1 \times L_2$ has a prime basis.

$L_1 + L_2$ has a prime basis.

$X \rightarrow L_2$ has a prime basis.

What else?

$$\llbracket \{ a^n, b^m \} \rrbracket$$

$$[a]^n [b]^m$$

$$A \subseteq B$$

$$[B] \bmod [A] = 0$$

A \cup *B*

$$\llbracket A \rrbracket \times \llbracket B \rrbracket$$

$\llbracket \langle a, b, c \rangle \rrbracket$

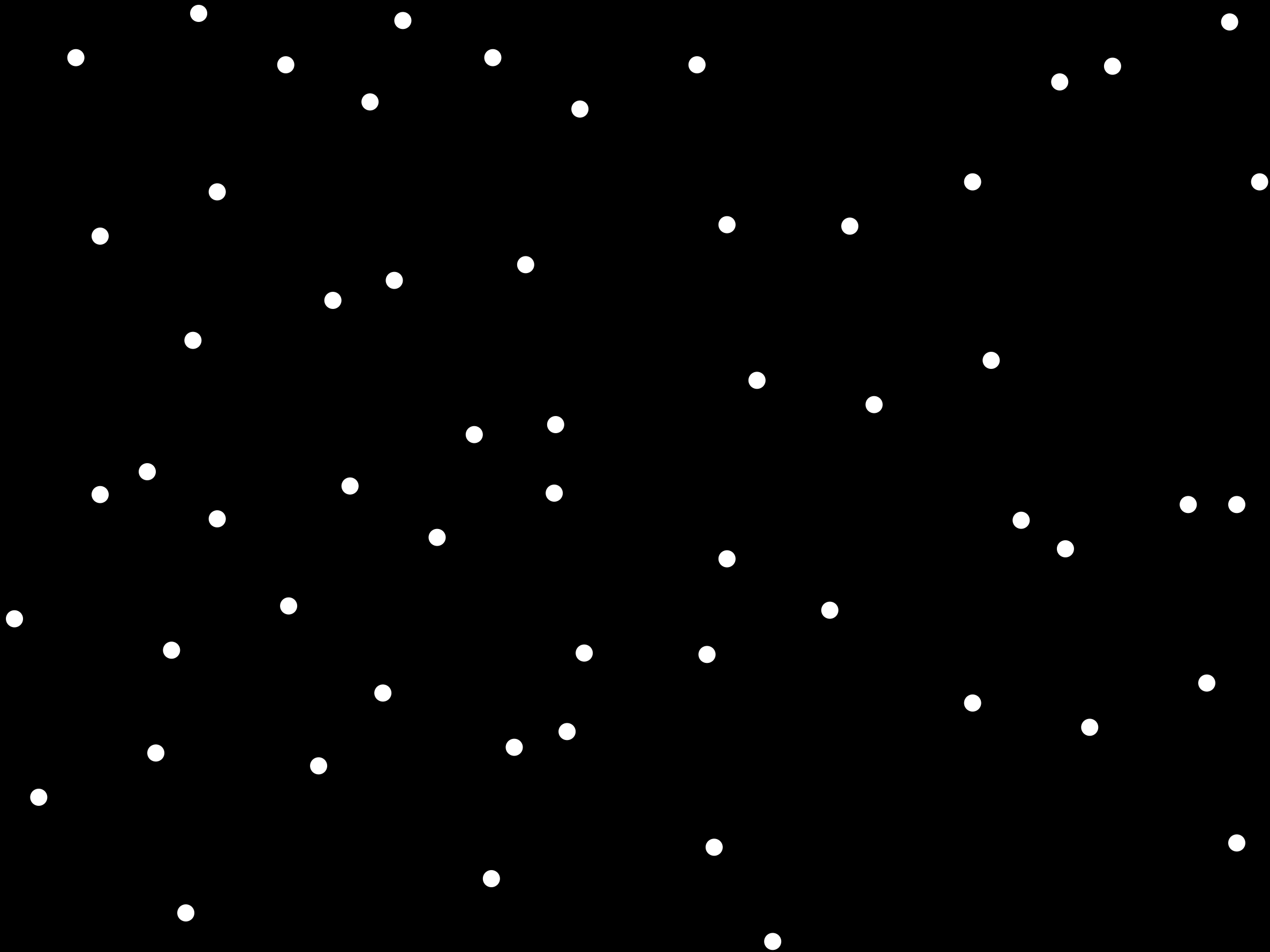
$\llbracket a \rrbracket$ $\llbracket b \rrbracket$ $\llbracket c \rrbracket$
 p_1 p_2 p_3

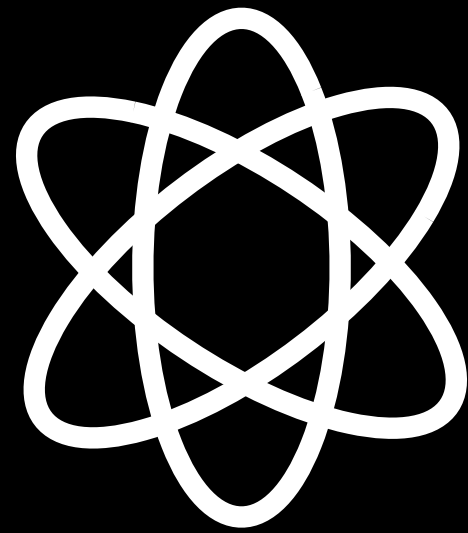
Wait a minute...

gcd is $O(n^2)$

mod is $O(n^2)$

How is this more efficient?





Flow sets are sparse.

99% of flow sets: < 5 values

Median flow set: 2 values

Primes are dense.

U

U \ln *U*

1,000,000 abstract values?

23 bit prime

Most flow sets fit in a word.

Most of the time, $n = 1$.

If not, great locality.

4-6%

^{2x} 5x 8x

1000x

Programming

is about

making choices.

3 E's

Elegance

Efficiency

Efficacy

Programmers:

Pick any two

Functional

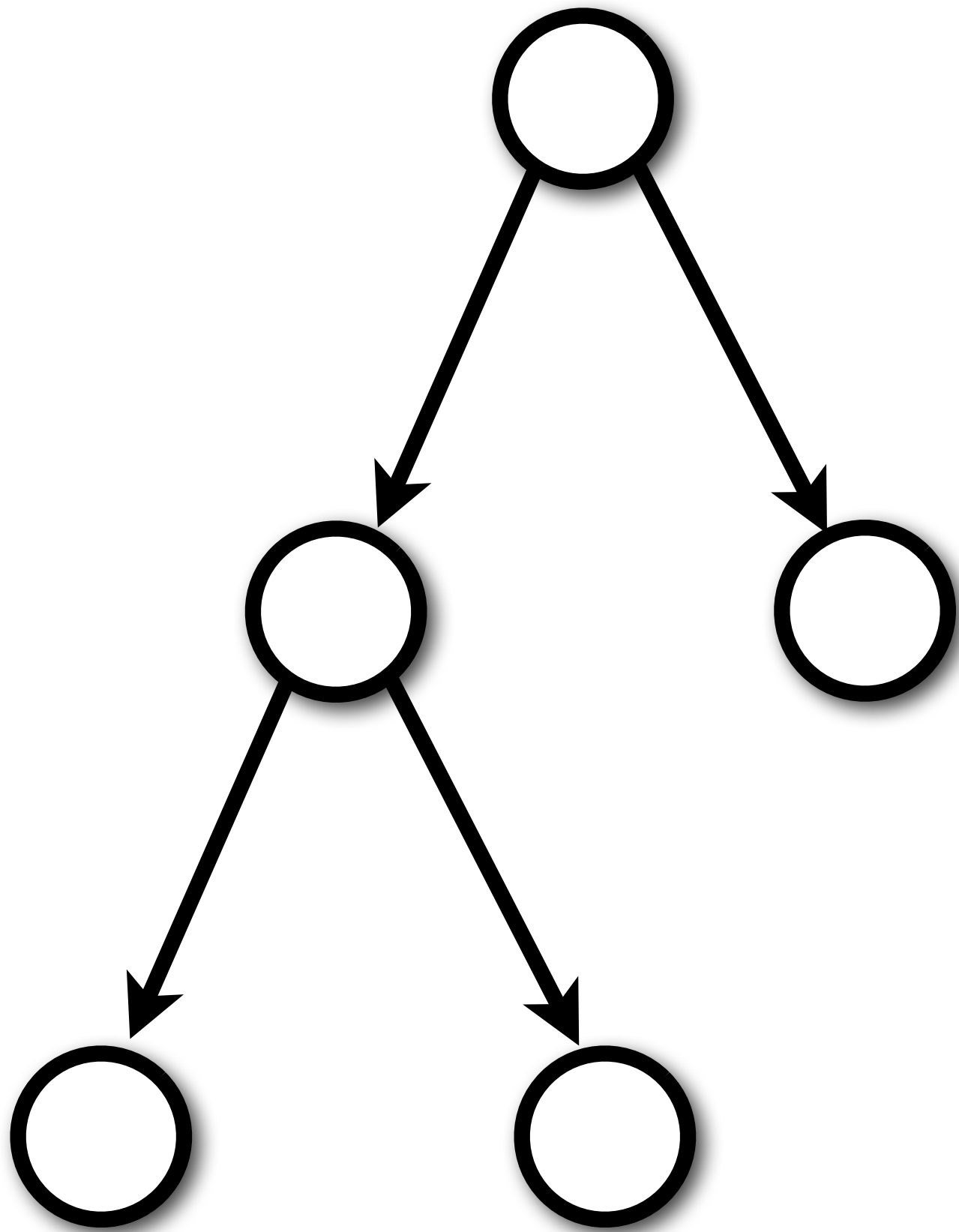
Programmers:

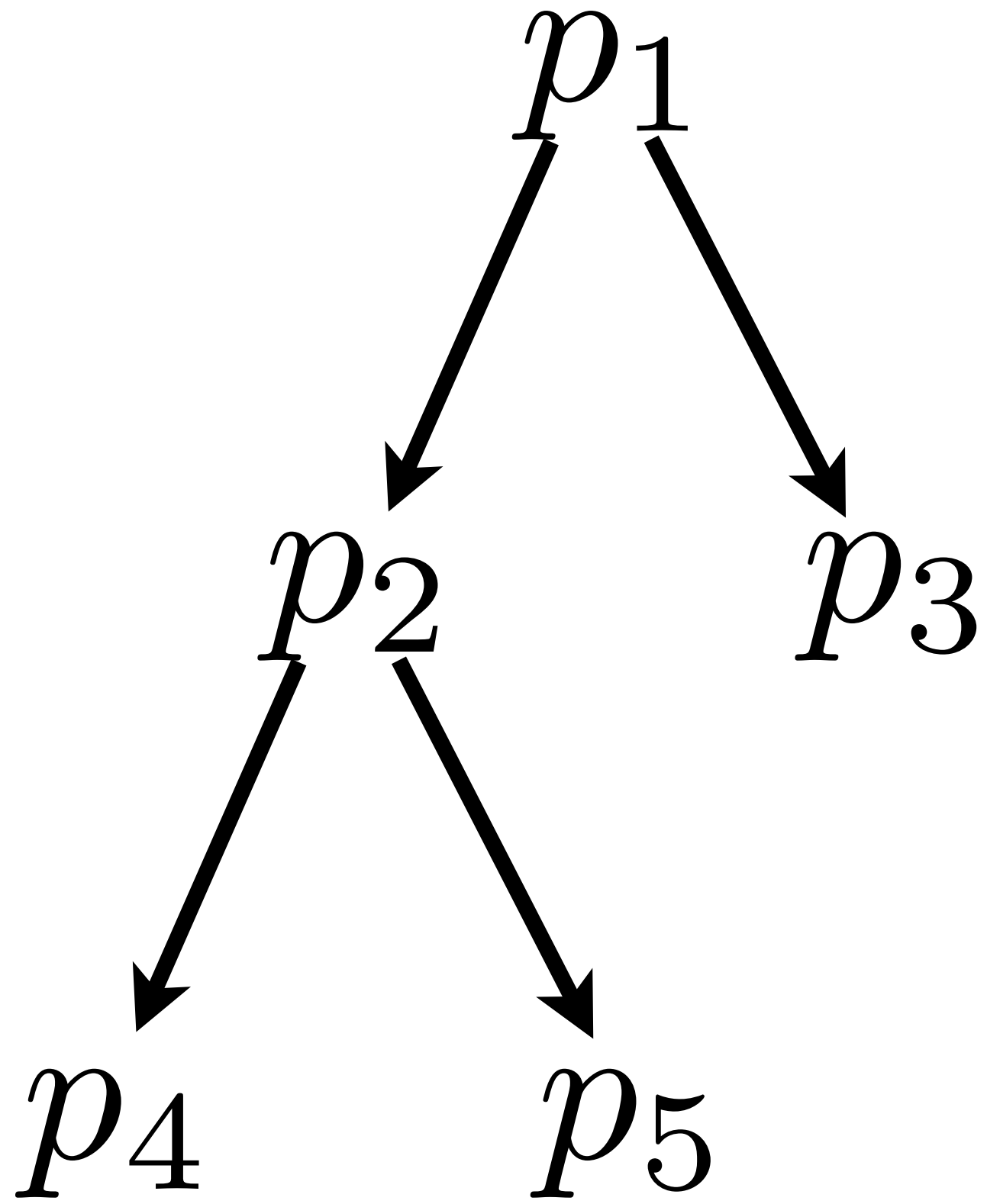
Pick any three

Questions?

Algebraic data types?

deriving (Hashable)





*p*₁ *p*₂ *p*₃ *p*₄ *p*₅

