Concurrent Orchestration in Haskell

galo

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Code Puzzle

```
foo :: (a \rightarrow s \rightarrow s) \rightarrow s \rightarrow 0rc a \rightarrow 0rc s
foo f s p = do a <- newMVarM s
```

x <- p
v <- takeMVarM a
let w = f x v
putMVarM a w
return w</pre>

This code implements a well-known idiom as we go on, try to figure out what it is...

Outline

- Concurrent scripting
- Laws
- Thread management



Testing Xen Virtual Machines



- Tester talks with each of the VMs concurrently
- Many possible behaviors are "correct" / "incorrect"
- Timeouts, VMs dying, etc.
- Subtle concurrency bugs in test framework

fplang :: Orc String
fplang = return "Haskell" <|> return "ML" <|> return "Scheme"





```
metronome :: Orc ()
```

metronome = return () <|> (delay 2.5 >> metronome)







least x y = if price x < price y then x else y threshold x = guard (price x < 300) >> return x



queens = fmap show (extend [])
<|> return ("Computing 8-queens...")

```
extend :: [Int] -> Orc [Int]
extend xs = if length xs == 8
    then return xs
    else do
        j <- listOrc [1..8]
        guard $ not (conflict xs j)
        extend (j:xs)</pre>
```



```
conflict :: [Int] -> Int
conflict = ...
```

listOrc :: [a] -> Orc a
listOrc = foldr (<|>) stop . map return

*Main> printOrc (queens)
Ans = "Computing &-queens..."
Ans = "[5,7,1,3,8,6,4,2]"
Ans = "[5,2,4,7,3,8,6,1]"
Ans = "[6,4,2,8,5,7,1,3]"
Ans = "[5,3,8,4,7,1,6,2]"
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:

*Main> printOrc (queens)
Ans = "Computing &-queens..."
Ans = "[4,2,7,3,6,8,5,1]"
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Ans = "[3,6,4,2,8,5,7,1]"
Ans = "[2,7,3,6,8,5,1,4]"

Orc Example



```
baseball :: Orc (String,String)
baseball = do
```

team <- prompt "Name a baseball team"
 `after` (12, return "Yankees")
 <!> prompt "Name another team"
 `notBefore` 10
 <!> (delay 8 >> return "Mariners")

```
agree <- prompt ("Do you like "++team++"?")
`after` (20, guard (team/="Mets") >> return "maybe")
```

return (team, agree)



baseball :: Orc (String,String) baseball = do

Name a baseball team Mets_

Name another team

agree <- prompt ("Do you like "++team++"?")
`after` (20, guard (team/="Mets") >> return "maybe")

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Name a baseball team Mets_

Name another team

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Do you like Mariners?

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Do you like Mariners?

Do you like Mets? ______galois

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putMVarM a w
return w</pre>



Orc Code

% printOrc (scan (+) 0 \$ listOrc [1,2,3,4,5])



Orc Code

% printOrc (scan (+) 0 \$ listOrc [1,2,3,4,5]) Ans = 1

Ans = 3

Ans = 6

Ans = 11

Ans = 15

%



Layered Implementation

	•	 Layered implementation — layered semantics
Orc Scripts		 Properties at one level depend on properties at the level below
Orc Monad	multiple results	
HIO Monad	thread control	 What properties should Orc terms satisfy?
IO Monad	external effects	-Hence, what properties should
		be built into HIO?

Transition Semantics

 Unresolved question: what laws should the basic operations of the IO monad satisfy?
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Key Definitions

type Orc a = (a -> HIO ()) -> HIO ()

return x = $\setminus k \rightarrow k x$ p >>= h = $\setminus k \rightarrow p (\setminus x \rightarrow h x k)$ p <|> q = $\setminus k \rightarrow fork (p k) >> q k$

stop = $k \rightarrow return$ ()

runOrc $p = p (\langle x - \rangle return ())$



Bind

type Orc a = (a -> HIO a) -> HIO a

return x =
$$k \rightarrow k$$
 x
p >>= h = $k \rightarrow p (x \rightarrow h x k)$
p <|> q = $k \rightarrow fork (p k) >> q k$



=



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=



Par

type Orc a = (a -> HIO a) -> HIO a

return x =
$$k \rightarrow k$$
 x
p >>= h = $k \rightarrow p (x \rightarrow h x k)$
p <|> q = $k \rightarrow p k \rightarrow q k$







- Give p a continuation that will store its result
- Return the "value" that accesses that result for the then current continuation



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eagerly :: Orc a -> Orc (Orc a) eagerly $p = \langle k - \rangle$ do r <- newEmptyMVarM</pre> forkM (p `saveOnce` (r)) k (\k' -> readMVarM r >>= k') saveOnce :: Orc a -> (MVar a) -> HIO () p `saveOnce` (r) = do p (\x -> putMVarM r x)

- Give p a continuation that will store its result (but once only even if duplicated)
- Return the "value" that accesses that result for the then current continuation

Eagerly

p `saveOnce` (r) = do

ticket <- newMVarM ()</pre>

- p (\x -> takeMVarM ticket >> putMVarM r x
 - Give p a continuation that will store its result (but once only even if duplicated)
 - Return the "value" that accesses that result for the then current continuation

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)

Eagerly

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eagerly :: Orc a -> Orc (Orc a) eagerly $p = \langle k - \rangle$ do r <- newEmptyMVarM</pre> e <- newLocality local e \$ forkM (p `saveOnce` (r,e)) k (\k' -> readMVarM r >>= k') saveOnce :: Orc a -> (MVar a,Locality) -> HIO () p `saveOnce` (r,e) = do ticket <- newMVarM ()</pre>

- p (\x -> takeMVarM ticket >> putMVarM r x >> close e)
 - Give p a continuation that will store its result (but once only even if duplicated)
 - Return the "value" that accesses that result for the then current continuation
 - Thread management can be carried over too



```
sync :: (a->b->c) -> Orc a -> Orc b -> Orc c
sync f p q = do
po <- eagerly p
qo <- eagerly q
return f <*> po <*> qo
```

```
notBefore:: Orc a -> Float -> Orc a
p `notBefore` w = sync const p (delay w)
```

```
    Entering the handle waits 
for the result
```

- Synchronization
- cut



```
sync :: (a->b->c) -> Orc a -> Orc b -> Orc c
sync f p q = do
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```
• Entering the handle waits for the result
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- Synchronization
- cut

```
notBefore:: Orc a -> Float -> Orc a
```

p `notBefore` w = sync const p (delay w)

cut:: Orc a -> Orc a
cut p = do
 po <- eagerly p
 po</pre>





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p `notBefore` w = sync const p (delay w)
```

- Entering the handle waits for the result
- Synchronization
- cut

```
cut:: Orc a -> Orc a
cut p = do
    po <- eagerly p
    po</pre>
```

cut:: Orc a -> Orc a cut = join . eagerly



Orc Laws

Left-Return: (return x >>= k) = k xRight-Return: (p >>= return) = pBind-Associativity: ((p >>= k) >>= h) = (p >>= (k >=> h))

Stop-Identity:p < l > stop = pPar-Commutativity:p < l > q = q < l > pPar-Associativity:p < l > (q < l > r) = (p < l > q) < l > r

Left-Zero: (stop >>= k) = stopPar-Bind: ((p <|> q) >>= k) = ((p >>= k) <|> (q >>= k))

Non-Laws

Bind-Par?: $p \implies (\x -> h x < |> k x) = (p \implies h) < |> (p \implies k)$ **Right-Zero?:** $p \implies$ stop = stop



Non-Laws

Bind-Par?: $p >>= (\x -> h x <|> k x) = (p >>= h) <|> (p >>= k)$ **Right-Zero?:** p >> stop = stop

p `until` done = cut (silent p <|> done)
silent p = p >> stop



Non-Laws

Bind-Par?: $p \implies (\x -> h x < |> k x) = (p \implies h) < |> (p \implies k)$ **Right-Zero?:** $p \implies$ stop = stop

```
p `until` done = cut (silent p <|> done)
silent p = p >> stop
```

```
hassle = (metronome >> email("Simon", "Hey!"))
`until`
    (delay 60 >> return ())
```



Eagerly Laws

Eagerly-Par: eagerly $p \rightarrow = (x \rightarrow k \times \langle p \rangle) = (eagerly p \rightarrow k) \langle p \rangle$

Eagerly-Swap:

- do y <- eagerly p = do x <- eagerly q
 - x <- eagerly q y <- eagerly p
 - return (x,y) return (x,y)

Eagerly-IO: eagerly (ioOrc m) >> p = (ioOrc m >> stop) <|> p

Val

val :: Orc a -> Orc a

val p = $k \rightarrow do$

- r <- newEmptyMVarM</pre>
- e <- newLocality

local e \$ forkM (p `saveOnce` (r,e))

k (unsafePerformIO \$ readMVarM r)



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saveOnce :: Orc a -> (MVar a,Locality) -> HIO ()

ticket <- newMVarM ()</pre>

p (\x -> takeMVarM ticket >> putMVarM r x >> close e)

• The implementation of val (the alternative that uses lazy thunks) is almost identical

Example

```
quotesVal :: Query -> Query -> Orc Quote
```

```
quotesVal srcA srcB = do
```

quoteA <- val \$ getQuote srcA</pre>

```
quoteB <- val $ getQuote srcB</pre>
```

cut (publish (least quoteA quoteB)

<l> (threshold quoteA)

<l> (threshold quoteB)

```
<l> (delay 25 >> (publish quoteA </> publish quoteB))
```

<l> (delay 30 >> return noQuote))

```
publish :: NFData a \Rightarrow a \rightarrow 0rc a
publish x = deepseq x $ return x
```

- Good: use the lazy values directly
- Bad: have to be careful about evaluation

HIO Monad



- Don't want the programmer to have to do explicit thread management
 - Nested groups of threads
- Want richer equational theory than IO
 - e.g. by managing asynchronous exceptions

HIO Monad

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- Don't want the programmer to have to do explicit thread management
 - Nested groups of threads
- Want richer equational theory than IO
 - e.g. by managing asynchronous exceptions

first :: Int -> Orc a -> Orc a
first n p = do
vals <- newEmptyMVarM
end <- newEmptyMVarM
echo n vals end
 <!> silent (generate p vals end)

```
generate p vals end =
```

(p >>= putMVarM vals) `until` takeMVarM end

```
echo n vals end = loop n
where loop 0 = silent $ putMVarM end ()
loop n = do x <- takeMVarM vals
return x <l> loop (n-1)
```

- Use MVars to communicate
- Use `until` to kill-off work when finished

Orc Example



Standard function:

filterM _ [] = return []
filterM p (x:xs) = do
 b <- p x
 ys <- filterM p xs
 return (if b then x:ys else ys)</pre>

Final Fun

baz :: [a] -> Orc [a]
baz xs = filterM pred xs

This code implements a well-known function what is it?

pred x = return False <l> return True

