

Haxell

adding regular types to Haskell

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- Original goal – an Xduce style type system for a message passing language
- New goal – full regular type integration into Haskell
 - Smooth integration of regular and algebraic types
 - Support for infinite hedges

Other Work

- Similar to Cduce in pattern matching
 - Does not support negative patterns
- Very similar to XHaskell except
 - intend to integrate type checking into the compiler, rather than continue with generating types
 - intend to support parametric polymorphism
 - will require less type annotations
 - the compiler “knows” which subsumption checks need to be made

Notation

- `<foo>True</foo><bar>10</bar><bar>20</bar><baz>“some text”</baz>`
- `:foo[True] :bar[10] :bar[20] :baz[“some text”]`
- Type:
 - `(| :foo[Bool], :bar[Integer]*, :baz[String] |)`
 - types nailed down quite specifically
- Change all the element names from foo to oof, bar to rab, and baz to zab:
 - `flipName:: (| :foo[Bool], :bar[Integer]*, :baz[String] |) -> (| :oof[Bool], :rab[Integer]*, :zab[String] |)`
 - even though contents of foo, bar, and baz are irrelevant

Type variables

- Haskell type variables:
 - `(| :foo[t1], :bar[t2]*, :baz[t3] |)`
- `flipName:: (| :foo[t1], :bar[t2]*, :baz[t3] |) -> (| :oof[t1], :rab[t2]*, :zab[t3] |)`
- Add elements:
 - `addem (|:foo[_], :bar[x]*, :baz[_]|) = fold (+) x`
- Types:
 - `addem:: (| :foo[t1], :bar[t2]*, :baz[t3] |)-> Integer`
 - `addem::(Num t2) => (| :foo[t1], :bar[t2]*, :baz[t3] |)-> t2`

New types

- `Foo t1 t2 t3 = (| :foo[t1], :bar[t2]*, :baz[t3] |)`
- `Bar = Foo String Integer Bool`

Weak Matching

- resolve ambiguity through weak matching:
 - $(| (:foo[], t1) | (t2, :bar[]) |)$
 - $t2$ could match $:foo$ and $t1$ could match $:bar$
 - weak matching, we exclude $t2$ from starting with $:foo$ and thereby resolve the ambiguity

Pattern Matching

- `addem (|:foo[_], :bar[x@Integer]*, :baz[_]|) = fold (+) x`
- `addem (|:foo[_], :bar[x@Float]*, :baz[_]|) = fold (*) x`
- `addem (|:foo[_], :bar[x@Double]*, :baz[_]|) = fold (/) x`

Current state

- Extended syntax (hacked Haskell grammar)
- Translation to “standard” Haskell (may require some apparently obscure declarations)
- Subsumption checks are performed once (if “—O”) at runtime and then resolve to True
- Pattern matching closer to Cduce than Xduce
- Marshal/Unmarshal between Haskell’s types and regular types is possible as defined by user.

Implementation

- A `newtype` for every regular type either
 - Defined by programmer, or
 - Generated from a pattern in a match
- Each pattern also generates a tuple type corresponding to the bound variables
- Also instance of class `XMark` providing info for
 - Validation
 - Casting to/from the type
 - Pattern matching – pattern match returns Either tuple error-message
- Finally, a structure containing all regex's is generated for validating and pattern matching.

Example regex type

```
regex Envelope =  
  ( | :envelope[ :headers[Header**]?? ,  
    :body[Any**] ] | )
```

becomes

```
newtype Envelope = Envelope [Tree]  
Instance XMark Envelope () where  
  -- "type" a hedge as an Envelope  
  create x = Envelope x  
  -- remove type so x can be cast to another type  
  decreate (Envelope x) = x  
  -- type name to be passed to validator  
  value x = "Envelope"
```

Pattern example

```
f (|:integers[int@Integer]**
   :strings[str@String]**|) = ...
```

becomes

```
newtype Gensym0 = Gensym0 [Tree]
type Gensym1 = ([Integer], [String])
instance XMark Gensym0 Gensym1 where
  ... as above ...
  -- pattern variable names
  varList x = ["int", "str"]
  -- convert results of pattern match to a Gensym1 tuple
  toTuple _ (Right [int, str]) =
    Just (create int, create str)
  toTuple _ (Left errorMessage) = error errorMessage
```

Implementation, cont.

- Instances of Cast where trees must be cast from one type to another
 - given `foo::Foo -> Bar` and `bar::Bar, (foo bar)` requires `Bar <: Foo` check
 - given instance `Cast Bar Foo`,
 - `foo (cast bar)` does the right thing
 - One runtime check (`if (Bar <: Foo) then ...`) becomes (`if True then ...`) if compiled with `-O`
- but these instances must currently be added manually
 - programmer changes `(foo bar)` to `(foo (cast bar))`
 - compiler will list all the instances to be inserted.
 - next step is to automate this

Implementation, cont.

- **Serialize/Deserialize**

- `class Castor regtype algebraic section`

- Given `deserialize::section->algebraic` then for any `regtype`, given `cast::regtype-> section` there is `deserialize::regtype-> algebraic`

- `class Xmlable algebraic regtype`

- **Much of this is very similar to patterns in “Scrap Your Boilerplate”**

- in particular, our “cast” is a slight generalization of the one in that paper.

Future Goals

- Compile time type checking inside GHC
 - remove most run time checks and generated code
- Default (de)serialization for algebraic types with user override
- Parametric polymorphism and type inference supporting infinite hedges

laziness and (non)determinism

Suppose we call a function with an infinite hedge:

$$f \left(\left(\left[\text{:int}[x] \right] \right) \mid x \leftarrow [1..] \right)$$

Suppose we have patterns:

- $f \left(\left[\text{:int}[a@Int]? \right], \left(\left[\text{:int}[b@Int] \right], \left[\text{:int}[c@Int] \right] \right)^* \mid \right) = \dots$
no assignment until finished and will diverge while matching
- $f \left(\left[\text{:int}[a@Int]? \right] \mid \right) = \dots$
 $f \left(\left(\left[\text{:int}[b@Int] \right], \left[\text{:int}[c@Int] \right] \right)^+ \mid \right) = \dots$
 $f \left(\left[\text{:int}[a@Int] \right], \left(\left[\text{:int}[b@Int] \right], \left[\text{:int}[c@Int] \right] \right)^+ \mid \right) = \dots$
each deterministic, but cannot be distinguished in bounded time, so it will diverge
- $f \left(\left(\left(\left[\text{:int}[b@Int] \right], \left[\text{:int}[c@Int] \right] \right)^*, \left[\text{:int}[a@Int]? \right] \right) \mid \right) = \dots$
matches in bounded time, allows processing of b and c, but accessing a will diverge

- Preference for the last alternative – matching and assignment should happen in bounded time, as in ordinary Haskell.

What's a hedge?

- Hedge

```
([:foo[10 12 13] :bar[:foo["abcde", :bar[]]]) =>  
<foo>10 12 13</foo> <bar><foo>abcde</foo><bar/></bar>
```

- Regular expression type

```
regex Foo = ([:foo[Integer** | String], :bar[Foo??]|)  
[namespace]:name for tag
```

“|” separates choice, “,” separates sequence, will add “&” for unordered

- Pattern (notice the variable bindings)

```
f(|:foo[(intList@Integer)** | astr@String], _|) =  
  (strval,intval)
```

where

```
strval= if ((length astr)>0)  
         then (head astr) else "",  
intval = foldl (+) 0 intList)
```