

# An informal description of Standard OIL and Instance OIL

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## Abstract

As the web becomes increasingly ubiquitous, more humans and computer programs need to

interact with data on websites. Seamless exchange of information has become the key issue for successful web deployments. Ontologies provide a way of capturing a shared understanding of terms that can be used by humans and programs to aid in information exchange. The ontology needs to be specified in some language. This paper introduces the newest version of OIL – the ontology inference layer – a proposal that presents a layered approach to a standard ontology language. It includes modelling constructs found in many widely knowledge representation languages, is compatible with RDFS, and includes a precise semantics for describing term meanings (and thus also for describing implied information).

This language has also been used as the basis for defining DAML-Ont, the ontology markup language from the DARPA sponsored DAML initiative (DAML stands for DARPA Agent Markup Language, see <http://www.daml.org>). A version of the DAML language, with working name DAML-OIL, was proposed in a message to the rdf-logic mailing list (<http://lists.w3.org/Archives/Public/www-rdf-logic/2000Nov/0094.html>) and is very close to Standard OIL as described below.

## Introduction

Ontologies have been gaining popularity as a method of providing a specification of a controlled vocabulary. Simple knowledge organizations such as Yahoo's taxonomy provide notions of generality and term relations but classical ontologies attempt to capture precise meanings of terms. In order to specify meanings, an ontology language must be used. OIL is an effort to produce a layered architecture for specifying ontologies. The language has been designed so that

- it provides the modelling primitives commonly used in frame-based ontologies
- it has a simple, clean, and well defined semantics based on description logics
- automated reasoning support (i.e., consistency and subsumption checking) may be specified and provided in a computationally efficient manner

It is expected that the core language will be extended to handle additional representation and reasoning needs. This initial presentation includes Standard OIL – a language intended to capture the necessary mainstream modelling primitives that both provide adequate expressive power and are well understood thereby allowing the semantics to be precisely specified and complete inference to be viable. While Standard OIL includes modelling constructs that allow individual fillers to be specified in term definitions, a full-fledged database capability is not included. Instance OIL includes a thorough individual integration. Heavy OIL includes additional representational (and reasoning) capabilities.

This document provides an informal description of the modelling primitives of the OIL dialects "OIL- Core" and "Instance OIL". It only gives a compact and informal description of these languages, plus a simple illustrative example. For more discussion and motivation, see the papers at [http:// www.ontoknowledge.org/oil/papers.shtml](http://www.ontoknowledge.org/oil/papers.shtml).

OIL's machine readable syntax is defined as an XML DTD (<http://www.ontoknowledge.org/oil/dtd/>), an XML Schema definition (<http://www.ontoknowledge.org/oil/xml-schema/>) and an RDF Schema definition (<http://www.ontoknowledge.org/oil/rdf-schema/>). To improve human readability OIL, also has a more compact pseudo syntax where keywords are indicated by **bold faced** text, and grouping of sub-content is indicated by indentation. The formal definition of this human-readable syntax can be found at <http://www.ontoknowledge.org/oil/syntax>. In this document, we give an informal

description of this human-readable syntax, since it is the easiest way to get acquainted with the modelling primitives of the language. We refer to the above URL's for the formal definitions.

## An informal description of Standard OIL Highlights

This section is for the casual user of OIL or the reader who is interested in a quick introduction to the most common important features of OIL. If the next section will be read, this section may be skipped.

An OIL ontology contains descriptions of classes, slots, and individuals. Classes are collections of objects. They are unary predicates such as *person*. Classes may be related to other classes by stating that one is a subclass of another – i.e., *person* is a subclass of *mammal*. Anything that is an instance of a *person* must also be an instance of a *mammal*. In this document, keywords are specified in bold. We could define *person* below:

```
class-def person  
    subclass-of mammal
```

Classes will typically also contain information about how their members relate to other objects. Slots are binary relations. They may also be related to each other via the notion of subrelation. For example, the slot *daughter-of* could be defined as follows:

```
slot-def has-daughter  
    subslot-of has-child
```

This defines a slot *has-daughter* that is a subslot of *has-child*, i.e., every pair  $(x,y)$  that is an instance of *has-daughter* must also be an instance of *has-child*. Slots may also be constrained through the use of slot-constraints. One can limit the type of a slot filler by stating that it must be of a particular type. For example, one can require that the filler of a *has-daughter* slot must be a female as follows:

```
slot-constraint has-daughter  
    value-type female
```

There are a number of kinds of slot constraints. The number of fillers a slot can have may be specified through the notion of cardinality. For example, one might define a class of people who have either one or two daughters who are doctors as follows:

```
slot-constraint has-daughter  
    max-cardinality 2 doctor  
    min-cardinality 1 doctor
```

If one defines the same cardinality for the minimum and maximum number, a shorthand notation is provided:

```
slot-constraint has-daughter  
    cardinality 1 doctor
```

An enumerated set of possible fillers may also be defined. For example, it is useful to list a set of possible colors for wine might be stated as:

```
slot-constraint has-color  
    value-type (one-of Red White Rose)
```

The enumerated set construct has implications on cardinality. Since the filler of the has-color slot above needs to be a member of the set containing three elements, then there is an inferred maximum cardinality of 3. There is no minimum cardinality implied from this description, since it is only saying that *if* there is a filler, then it must be an instance of the enumerated set.

Particular fillers may also be specified. For example, a white wine may be explicitly defined as something that is a wine whose color slot is filled with the color white.

```
class-def white-wine
  subclass-of wine
  slot-constraint has-color
  has-filler white
```

Filler specification also interacts with the cardinality restrictions. In the example above, since one filler is specified, there is an implied minimum cardinality of 1 on the has-color slot.

Term descriptions may also be combined by using Boolean connectives. Thus it is possible to use the connectors: and, or, and not to combine descriptions. **And** takes two or more expressions and conjoins them, **or** takes two or more expressions and treats them as a disjunction, and **not** takes a single expression that is negated. Expressions may be recursive, thus arbitrarily complex Boolean expressions may be formed:

```
(white-wine and (napa-wine or (not sweet-wine)))
```

For a more complete (informal) introduction to the language, the next section should be read. If a formal description of the language is desired, a grammar specification, semantics specification, and longer papers are available from: [http:// www.ontoknowledge.org/oil/](http://www.ontoknowledge.org/oil/) .

## **An informal description of Standard OIL**

An OIL ontology is a structure made up of several components, some of which may themselves be structures, some of which are optional, and some of which may be repeated. We will write **component**<sup>?</sup> to indicate an optional component, **component**<sup>+</sup> to indicate a component that may be repeated one or more times (i.e., that must occur at least once) and **component**<sup>\*</sup> to indicate a component that may be repeated zero or more times (i.e., that may be completely omitted). In general we assume that identifiers (for classes, roles etc) do not coincide with keywords of the language (such as **slot-constraint**, **or**, etc). However, this section is intended to be descriptive rather than precise: for a precise formal definition please refer to the URL's mentioned above.

An OIL ontology is delineated by the keywords **begin-ontology** and **end-ontology**, and consists of the actual *ontology definition*, defining a particular ontological vocabulary, preceded by an *ontology container*, which is concerned with describing features of such an ontology, like author, name, subject, etc. For representing metadata of ontologies, we make use of the Dublin Core Metadata Element Set (Version 1.1) [Dublin Core] standard. We will discuss both elements of an ontology specification in OIL. We start with the ontology container and will then discuss the backbone of OIL, the ontology definition.

### **Ontology Container**

We adopt the components as defined by the Dublin Core Metadata Element Set, Version 1.1 for

the ontology container part of OIL. Although every element in the Dublin Core set is optional and repeatable, in OIL some elements are required or have a predefined value. Required elements are written as element+. Some of the elements can be specialized with a qualifier which refines the meaning of that element. In our shorthand notation we will write element.qualifier. The precise syntax based on RDF is given in [Miller et al., 1999], and in the appendix. Here we provide our pseudo-XML syntax explained above.

**title**<sup>+</sup> The name of the ontology, e.g., "African animals".

**creator**<sup>+</sup> The name of an agent (i.e., a person, a group of persons, or a software agent) that created the ontology.

**subject**<sup>\*</sup> Keywords or classification code describing the subject the ontology deals with.

**description** Natural language text describing the content of the ontology, e.g., "A didactic example ontology describing African animals". Besides this description, there is one special description element required, which has the **release** qualifier:

**description.release** The version of the ontology (a number), e.g. 1.01.

**publisher**<sup>\*</sup> Defining the entity that is responsible for making the resource available.

**contributor**<sup>\*</sup> The name of an agent (i.e., a person, a group of persons, or a software agent) that helped to create the ontology.

**date**<sup>\*</sup> The date the ontology has been created, modified, or made available (see ISO 8601 for format instructions).

**type**<sup>+</sup> The nature of the resource. A predefined and required value is *ontology*, although this value is not yet in the Working Draft of the resource types [Guenther, 1999].

**format**<sup>\*</sup> The digital manifestation of the resource, recommended as a value is the MIME type of the resource, i.e. "text/xml".

**identifier**<sup>+</sup> The **URI** of the ontology.

**source**<sup>\*</sup> Optional references (**URI**) to sources from which the ontology is derived. E.g., a reference to a plain text description of the domain on which the ontology is based.

**language**<sup>+</sup> The language of the ontology. Obviously, one predefined and required value is "OIL". Other elements can contain the language of the content of the ontology, according to RFC 1766.

**relation**<sup>\*</sup> A list of references to other OIL ontologies. It is recommended to list all ontologies that are imported in the definition section with a **hasPart** qualifier. Other possible and meaningful qualifiers are **replaces**, **isReplacedBy**, **requires** and **isRequiredBy**. For example, to list an imported ontology, we write: **relation.hasPart** "http://www.ontosrus.com/animals/jungle.onto".

**rights**<sup>\*</sup> Information about rights held in and over the ontology.

### **Ontology definition**

Apart from various header fields encapsulated in its container, an OIL ontology consists of a set of definitions, optionally preceded by an important statement:

**import**<sup>?</sup> A list of one or more references to other OIL modules that are to be included in this

ontology. Each reference consists of a **URI** specifying where the module is to be imported from, e.g., "<http://www.ontosrus.com/animals/jungle.onto>". XML schemas and OIL provide the same (limited) means for composing specifications. Specifications can be included and the underlying assumption is that names of different specifications are different (via different prefixes).

**definition\*** Zero or more class definitions (**class-def**), axioms (**disjoint**, **covered**, **disjoint-covered**, **equivalent**), and slot definitions (**slot-def**), the structure of which will be described below.

A class definition (**class-def**) associates a class name with a class description. A **class-def** consists of the following components:

**type**<sup>?</sup> The type of definition. This can be either **primitive** or **defined**; if omitted, the type defaults to **primitive**. When a class is **primitive**, its definition (i.e., the combination of the following **subclass-of** and **slot-constraint** components) is taken to be a necessary but not sufficient condition for membership in the class. For example, if the primitive class **elephant** is defined to be a sub-class of **animal** with a slot constraint stating that *skin-color* must be **grey**, then all instances of **elephant** must necessarily be animals with grey skin, but there may be grey-skinned animals that are not instances of **elephant**. When a class is **defined**, its definition is taken to be a necessary *and* sufficient condition for membership of the class. For example, if the defined class **carnivore** is defined to be a sub-class of **animal** with a slot constraint stating that it *eats* meat, then all instances of **carnivore** are necessarily meat eating animals, and every meat eating animal is also an instance of **carnivore**.

**name** The name of the class (a string).

**documentation**<sup>?</sup> Some documentation describing the class (a string).

**subclass-of**<sup>?</sup> A list of one or more *class-expressions*, the structure of which will be described below. The class being defined in this **class-def** must be a sub-class of each of the class-expressions in the list.

**slot-constraint**<sup>\*</sup> Zero or more slot-constraints, a special kind of class-expression, the structure of which will be described below (note that a slot-constraint defines a class). The class being defined in this **class-def** must be a subclass of each slot-constraint.

A class-expression can be either a class name (some of which are *built-in*), an *enumerated-class*, a **slot-constraint**, or a boolean combination of class expressions using the operators **and**, **or** and **not**. The structure of these boolean combinations is as follows (note that they must be parenthesized):

**and:** A list of two or more class expressions that is to be treated as a conjunction. For example:

(meat **and** fish)

defines the class whose instances are all those individuals that are instances of both the class **meat** and the class **fish**.

**or:** A list of two or more class expressions that is to be treated as a disjunction. For example:

(meat **or** fish)

defines the class whose instances are all those individuals that are instances of either the class **meat** or the class **fish**.

**not:** An expression taking as a parameter a single class expression that is to be negated. For example,

(**not meat**)

defines the class whose instances are all those individuals that are not instances of the class **meat**.

Note that class expressions are recursively defined, so that arbitrarily complex expressions can be formed. For example

(**not (meat or fish)**)

defines the class whose instances are all those individuals that are not an instances of either the class **meat** or the class **fish**.

The built-in class names consist of **top**, **thing** and **bottom**. The meaning of these classes is pre-defined. **top** and **thing** are alternative names for the most general class. Every class is a sub-class of **top**, and every individual is an instance of **top**. **bottom** is the name of the least general (sometimes called empty or inconsistent) class. **bottom** is a sub-class of every class, and no individual is an instance of **bottom**.

An enumerated-class is a class that is defined by enumerating its instances. An enumerated-class consists of the key-word **one-of** followed by one or more individual names, with the whole expression being enclosed in parenthesis. For example,

(**one-of Leo Willie**)

defines the class whose instances are **Leo** and **Willie**.

In some situations it is possible to use a *concrete-type-expression* instead of a class-expression (e.g., in slot restrictions). A concrete-type-expression defines a range over some data type. Two data types are currently supported: **integer** and **string**. The expression **integer** defines the range of all integers (i.e., -infinity to +infinity) and the expression **string** defines the range of all strings. Sub-ranges can be defined using the expressions (**min x**), (**max x**), (**greater-than x**), (**less-than x**), (**equal x**) and (**range x y**), where both **x** and **y** are either integers or strings. Finally, expressions (of the same type) can be combined using the operators **and**, **or** or **not** as in class expressions. For example,

(**min 21**)

defines the data type consisting of all the integers greater than or equal to 21,

(**less-than 100**)

defines the data type consisting of all the integers less than 100 (i.e., less than or equal to 99),

(**greater-than "abc"**)

defines the data type consisting of all the strings that (lexically) succeed "abc" (e.g., "abd"),

(**or (equal "red") (equal "green") (equal "blue")**)

defines the data type consisting of the strings "red", "green" and "blue",

(**range 1 10**)

defines the data type consisting of all the integers greater than or equal to 1 and less than or equal to 10 and is equivalent to,

((**min 1**) **and** (**max 10**))

and,

(**equal "xyz"**)

defines the data type consisting of the string "xyz" and is equivalent to,

((**min "xyz"**) **and** (**max "xyz"**)).

A **slot-constraint** is a list of one or more constraints (restrictions) applied to a *slot* (sometimes called a *role* or an *attribute*). A slot is a binary relation (i.e., its instances are pairs of individuals), but a slot-constraint is actually a class definition — its instances are those individuals that satisfy the constraint(s). For example, if the pair (**Leo**, **Willie**) is an instance of

the slot **eats**, **Leo** is an instance of the class **lion** and **Willie** is an instance of the class **wildebeest**, then **Leo** is also an instance of the **has-value** constraint **wildebeest** applied to the slot **eats**. A **slot-constraint** consists of the following components:

**name** A slot name (a string). The slot is a binary relation that may or may not be defined in the ontology. If it is not defined, then it is assumed to be a binary relation with no globally applicable constraints, i.e., any pair of individuals could be an instance of the slot.

**has-value**? A list of one or more expressions (either class-expressions or concrete-type-expressions). Every instance of the class defined by the slot-constraint must be related via the slot relation to an instance of each expression in the list. For example, the **has-value** constraint:

**slot-constraint eats**

**has-value zebra wildebeest**

defines the class each instance of which **eats** some instance of the class **zebra** and some instance of the class **wildebeest**. Note that this does not mean that instances of the slot-constraint eat *only* **zebra** and **wildebeest**: they may also be partial to a little **gazelle** when they can get it. The **has-value** constraint:

**slot-constraint colour**

**has-value "red"**

defines the class each instance of which has the **colour** "red" (a string). Note that in the absence of some other constraint (e.g., a cardinality constraint), more than one **colour** may be specified. **has-value** expresses the *existential quantifier of Predicate logic*: for each instance of the class, there exists at least one value for this slot that fulfils the range restriction.

**value-type**? A list of one or more expressions (either class-expressions or concrete-type-expressions). If an instance of the class defined by the slot-constraint is related via the slot relation to some individual **x**, then **x** must be an instance or data value of each expression in the list. For example, the **value-type** constraint:

**slot-constraint eats**

**value-type meat**

defines the class each instance of which **eats** nothing that is not **meat**. Note that this does not mean that instances of the slot-constraint eat anything at all. The **value-type** constraint:

**slot-constraint age**

**value-type (min 21)**

defines the class each instance of which does not have an **age** less than 21. Note that this does not mean that instances of the slot-constraint have any age at all. **value-type** expresses the *universal (for-all) quantifier of Predicate logic*: for each instance of the class, every value for this slot must fulfill the range restriction.

**has-filler**? A list of one or more individual names or data values (integers or strings). Every instance of the class defined by the slot-constraint must be related via the slot relation to each individual and data value in the list. For example, the **has-filler** constraint:

**slot-constraint friend**

**has-filler Zoe Willie**

defines the class each instance of which is a **friend** of both **Zoe** and **Willie**. Note that this is equivalent to the **has-value** constraint:

**slot-constraint friend**

**has-value (one-of Zoe) (one-of Willie)**

The **has-filler** constraint:

**slot-constraint age**

**has-filler 21**

defines the class each instance of which has an **age** of 21. Note that this is equivalent to the **has-**

**value** constraint:

**slot-constraint** *age*  
**has-value** (equal 21)

It is also worth reemphasizing that in the absence of other constraints (e.g., *age* being a functional slot), there is nothing to prevent instances of this class having more than one *age*.

**max-cardinality**? A non-negative integer  $n$  followed by an expression (either a class-expression or a concrete-type-expression). An instance of the class defined by the slot-constraint can be related to at most  $n$  distinct instances or data values of the expression via the slot relation. The expression can be omitted, in which case an instance of the class defined by the slot-constraint can be related to at most  $n$  distinct individuals or data values (regardless of their class or type) via the slot relation. For example, the **max-cardinality** constraint:

**slot-constraint** *friend*  
**max-cardinality** 2 antelope

defines the class, each instance of which has at most 2 *friends* that are antelopes.

**min-cardinality**? A non-negative integer  $n$  followed by an expression (either a class-expression or a concrete-type-expression). An instance of the class defined by the slot-constraint must be related to at least  $n$  distinct instances or data values of the expression via the slot relation. The expression can be omitted, in which case an instance of the class defined by the slot-constraint must be related to at least  $n$  distinct individuals or data values (regardless of their class or type) via the slot relation. For example, the **min-cardinality** constraint:

**slot-constraint** *friend*  
**min-cardinality** 3 wildebeest

defines the class, each instance of which has at least 3 friends that are wildebeests. Note that conflicting cardinality constraints is one way in which logical inconsistencies can arise in an ontology. For example, a class to which both the above min-cardinality and max-cardinality constraints applied would be logically inconsistent (could have no instances) if the ontology correctly represented the fact that a wildebeest is a kind of antelope.

**cardinality**? A non-negative integer  $n$  followed (optionally) by an expression (either a **class-expression** or a **concrete-type-expression**). This is simply shorthand for a pair of **min-cardinality** and **max-cardinality** constraints, both with the same  $n$  and expression. For example,

**slot-constraint** *friend*  
**cardinality** 1 zebra

is equivalent to

**slot-constraint** *friend*  
**max-cardinality** 1 zebra  
**min-cardinality** 1 zebra

and defines the class, each instance of which has exactly 1 *friend* that is a zebra.

An axiom asserts some additional fact(s) about the classes in the ontology, for example that the classes *carnivore* and *herbivore* are disjoint (can have no instances in common). Valid axioms are:

**disjoint** A list of two or more class expressions. All of the class expressions in the list are pairwise disjoint, i.e., there can be no individual that is an instance of more than one of the class expressions in the list. For example,

**disjoint** carnivore herbivore

states that no individual can be an instance of both carnivore and herbivore.

**covered** A class expression followed by a list of one or more class expressions that cover it. Every instance of the first class expression is also an instance of at least one of the

class expressions in the list. For example,

**covered animal by carnivore herbivore omnivore mammal**  
states that every instance of **Animal** is also an instance of at least one of **carnivore herbivore omnivore** or **mammal**.

**disjoint-covered** A class expression followed by a list of one or more class expressions that cover it and that are also pairwise disjoint. Every instance of the first class expression is also an instance of exactly one of the class expressions in the list. For example,

**disjoint-covered animal by carnivore herbivore omnivore**  
states that every instance of **animal** is also an instance of exactly one of **carnivore herbivore** or **omnivore**.

**equivalent** A list of two or more class expressions. All of the class expressions in the list are equivalent (i.e., have the same instances). For example,

**equivalent wildebeest gnu**  
states that an individual is **wildebeest** if and only if it is a **gnu** (i.e., **wildebeest** and **gnu** are synonyms).

A slot definition (**slot-def**) associates a slot name with a slot description. A slot description specifies global constraints that apply to the slot relation, for example that it is a transitive relation. A **slot-def** consists of the following components:

**name** The name of the slot (a string).

**documentation**<sup>?</sup> Some documentation describing the slot (a string).

**subslot-of**<sup>?</sup> A list of one or more **slots**. The slot being defined in this **slot-def** must be a sub-slot of each of the slots in the list. For example,

**slot-def daughter-of**  
**subslot-of child-of**

defines a slot **daughter-of** that is a subslot of **child-of**, i.e., every pair  $(x,y)$  that is an instance of **daughter-of** must also be an instance of **child-of**.

**domain**<sup>?</sup> A list of one or more class-expressions. If the pair  $(x,y)$  is an instance of the slot relation, then  $x$  must be an instance of each class-expression in the list. For example,

**slot-def eats**  
**domain animal**

defines a slot **eats** such that any individual that **eats** another individual must be an instance of **animal**.

**range**<sup>?</sup> A list of one or more expressions (either class-expressions or concrete-type-expressions). If the pair  $(x,y)$  is an instance of the slot relation, then  $y$  must be an instance or data value of each class-expression or concrete-type-expression in the list. For example,

**slot-def friend**  
**range animal**

defines a slot **friend** such that any individual that is a **friend** of another individual must be an instance of **animal**, and

**slot-def age**  
**range (min 0)**

defines a slot **age** such that if the pair  $(x,y)$  is an instance of **age**, then  $y$  must be a non-negative **integer**. Note that it is good practice to specify the range data type of a slot that is to be used for data values.

**inverse**<sup>?</sup> The name of a slot **S** that is the inverse of the slot being defined. If the pair  $(x,y)$  is an

instance of the slot *S*, then  $(y,x)$  must be an instance of the slot being defined. For example,

**slot-def** *eats*

**inverse** *eaten-by*

defines the inverse of the slot *eats* to be the slot *eaten-by*, i.e., if *x eats y* then *y is eaten-by x*.

**properties**? A list of one or more properties of the slot. Valid properties are:

**transitive** The slot is transitive, i.e., if both  $(x,y)$  and  $(y,z)$  are instances of the slot, then  $(x,z)$  must also be an instance of the slot. For example,

**slot-def** *bigger-than*

**properties** *transitive*

defines the slot *bigger-than* to be transitive, so if *Jumbo* the *elephant* is *bigger-than* *Robbie* the *rhino*, and *Robbie* the *rhino* is *bigger-than* *Walter* the *warthog*, then *Jumbo* must be *bigger-than* *Walter*. Note that no slot can be both transitive and functional.

**symmetric** The slot is symmetric, i.e., if  $(x,y)$  is an instance of the slot, then  $(y,x)$  must also be an instance of the slot. For example,

**slot-def** *lives-with*

**properties** *symmetric*

defines the slot *lives-with* to be symmetric, so if *Zoe* the *zebra* *lives-with* *Willie* the *wildebeest*, then *Willie* also *lives-with* *Zoe*.

**functional** The slot is functional, i.e., if  $(x,y)$  is an instance of the slot, then there is no *z* such that  $(x,z)$  is an instance of the slot and *y* is not equal to *z*. For example,

**slot-def** *has-mother*

**properties** *functional*

defines the slot *has-mother* to be functional. Note that no slot can be both functional and transitive.

## An informal description of Instance OIL

Instance OIL is a strict superset of Standard OIL. Instance OIL adds to Standard OIL the possibility to define instances of classes and roles, using the following two constructions:

An **instance-of** statement asserts that an individual is an instance of a class or classes. It consists of an individual name (a string) followed by one or more class-expressions. The individual must be an instance of each of the class-expressions in the list. For example,

**instance-of** *Zoe* *zebra*

states that *Zoe* is a *zebra*.

An **related** statement asserts that an individual is related to another individual or data value via a slot relation. It consists of the slot name and an individual name followed by either a second individual name or a data value. The first individual must be related to the second individual or data value via the slot relation. For example,

**related** *has-mother* *Zachariah* *Zoe*

states that *Zoe* is the *mother* of *Zachariah*, and

**related** *age* *Zoe* 35

states that *Zoe* is *age* 35.

## An example OIL ontology

The following example of an OIL ontology illustrates some of the key features of the language. The ontology is intended purely for didactic purposes and is not to be taken as an example of

good modeling practice.

### ontology-container

**title** "African animals"  
**creator** "Ian Horrocks"  
**subject** "animal, food, vegetarians"  
**description** "A didactic example ontology describing African and Asian animals"  
**description.release** "1.01"  
**publisher** "I. Horrocks"  
**type** "ontology"  
**format** "pseudo-xml"  
**format** "pdf"  
**identifier** "http://www.cs.vu.nl/~dieter/oil/TR/oil.pdf"  
**source** "http://www.africa.com/nature/animals.html"  
**language** "OIL"  
**language** "en-uk"  
**relation.hasPart** "http://www.ontosRus.com/animals/jungle.onto"

### ontology-definitions

**slot-def** *eats*  
    **inverse** *is-eaten-by*  
**slot-def** *has-part*  
    **inverse** *is-part-of*  
    **properties** transitive  
**slot-def** *comes-from*  
**slot-def** *age*  
    **range** (min 0)  
    **properties** functional  
**slot-def** *weight*  
    **range** (min 0)  
    **properties** functional  
**slot-def** *colour*  
    **range** string  
    **properties** functional  
**class-def** animal  
**class-def** plant  
**disjoint** animal plant  
**class-def** tree  
    **subclass-of** plant  
**class-def** branch  
    **slot-constraint** *is-part-of*  
        **has-value** tree  
**class-def** leaf  
    **slot-constraint** *is-part-of*  
        **has-value** branch  
**class-def** defined carnivore  
    **subclass-of** animal  
    **slot-constraint** *eats*  
        **value-type** animal  
**class-def** defined herbivore  
    **subclass-of** animal  
    **slot-constraint** *eats*

```

    value-type (plant or (slot-constraint is-part-of has-value plant))
disjoint carnivore herbivore
class-def giraffe
  subclass-of animal
  slot-constraint eats
  value-type leaf
class-def lion
  subclass-of animal
  slot-constraint eats
  value-type herbivore
class-def tasty-plant
  subclass-of plant
  slot-constraint eaten-by
  has-value herbivore carnivore
class-def elephant
  subclass-of animal
  slot-constraint eats
  value-type plant
  slot-constraint colour
  has-filler "grey"
class-def defined adult-elephant
  subclass-of elephant
  slot-constraint age
  has-value (min 20)
covered adult-elephant by (slot-constraint weight has-value (range 5000 8000))
class-def defined african-elephant
  subclass-of elephant
  slot-constraint comes-from
  has-filler Africa
class-def defined indian-elephant
  subclass-of elephant
  slot-constraint comes-from
  has-filler India
disjoint-covered elephant by african-elephant indian-elephant
class-def kenyan-elephant
  subclass-of elephant
disjoint kenyan-elephant indian-elephant
class-def defined african-animal
  subclass-of animal
  slot-constraint comes-from
  has-value ((one-of Africa) or (slot-constraint is-part-of has-filler Africa))
class-def defined asian-animal
  subclass-of animal
  slot-constraint comes-from
  has-value ((one-of Asia) or (slot-constraint is-part-of has-filler Asia))
class-def defined large-animal
  subclass-of animal
  slot-constraint weight
  has-value (min 1000)
class-def continent
class-def country

```

**instance-of** Africa continent  
**instance-of** Asia continent  
**instance-of** India country  
**related** *is-part-of* India Asia

Some points to note in the above ontology are:

- The class **carnivore** is a defined class, and **lion** can be recognized as a sub-class of **carnivore** because of its definition.
- The class **herbivore** is a defined class, and **giraffe** can be recognized as a sub-class of **herbivore** because of its definition. However, in this case the inference is a little more complex and is only valid because *has-part* is transitive and *is-part-of* is the inverse of *has-part*.
- The class **tasty-plant** is inconsistent. This is because **tasty-plant** is a kind of **plant** that is eaten by both **herbivores** and **carnivores**, but we have already stated that **carnivore** eat only **animals**, and that **animal** and **plant** are disjoint.
- The class **adult-elephant** is defined to be all elephants aged 15 (an integer) and over. An **adult-elephant** is also asserted to have a *weight* between 1500 and 3000 (an integer range). As a result, **adult-elephant** can be recognized as a sub-class of **large-animal**.
- The class **kenyan-elephant** is disjoint from **indian-elephant**. Moreover, **indian-elephant** and **african-elephant** form a disjoint covering of **elephant**. As a result, **kenyan-elephant** can be recognized as a sub-class of **african-elephant**.
- **Africa**, **Asia** and **India** are individuals, and **India** is *part-of* **Asia**. As a result, **indian-elephant** can be recognized as a sub-class of **asian-animal**.

## Current Limitations of OIL

Our starting point was to define a core language with the intention that additional (and possibly important) features be defined as a set of extensions (still with clearly defined semantics). Modelers will be free to use these language extensions, but it must be clear that this may compromise reasoning support. This seems to us a better solution than trying to define a single "all things to all men" language like Ontolingua. In this section we briefly discuss a number of features which are available in other ontology modeling languages and which are not, or not yet, included in OIL. For each of these features we briefly explain why we chose them, and mention future prospects where relevant.

**Default reasoning:** Although OIL does provide a mechanism for inheriting values from super-classes, such values cannot be overwritten. As a result, such values cannot be used for the purpose of modeling default values. If an attempt is made at "overwriting" an inherited attribute value, this will simply result in inconsistent class definitions which have an empty extension. For example, if we define the class "CS professor" with attribute "gender" and value "male", and we subsequently define a subclass for which we define the gender attribute as "female", this subclass will be inconsistent and have an empty extension (assuming that "male" and "female" are disjoint).

**Rules/Axioms:** As discussed above, only a fixed number of algebraic properties of slots can be expressed in OIL, plus a number of axioms relating classes (disjoint covers etc.). These axioms are sufficient for describing arbitrary subsubsumption relations (rules) or properties that must hold for all the items in the ontology. For example,

**covered Class-expression1 by Class-expression2**

states that every instance of Class-expression1 is also an instance of Class-expression2, while

**covered thing by Class-expression**

states that the properties defined in Class-expression must hold for every instance of every class in the ontology. Such a powerful feature is undoubtedly useful, e.g., to enforce the correct interpretation of concepts across multiple integrated ontologies. However, there may also be a requirement for other kinds of rules/axioms with different semantic interpretations.

**Further relation properties:** The properties that can be specified for relations in OIL are currently restricted to *inverse*, *functional*, *transitivity* and *symmetry*. Other reasonable candidates are *reflexivity*, *irreflexivity*, *antisymmetry*, *asymmetry*, *linearity* ( $aRb \ bRa$  for any pair  $a,b$ ), *connectivity* ( $aRb$  or  $a=b$  or  $bRa$  for any pair  $a,b$ ), *partial order* and *total order*. (Note that some of these can be defined in terms of each other). It would also be useful to be able to define *composite* relations.

**Modules:** OIL contains a very simple construction to modularize ontologies. In fact, this mechanism is identical to the namespace mechanism in XML and XML schema. It amounts to a textual inclusion of the imported module, where name-clashes are avoided by prefixing every imported symbol with a unique prefix indicating its original location. However, much more elaborate mechanisms would be required for the structured representation of large ontologies. Means of renaming, restructuring, and redefining imported ontologies must be available. Future extensions will cover parameterized modules, signature mappings between modules, and restricted export interfaces for modules.

**Limited Second-order expressivity:** Many existing languages for ontologies (KIF, CycL, Ontolingua) include some form of reification mechanism, which allows the treatment of statements of the language as objects in their own right, thereby making it possible to express statements over these statements. A full second order extension would be clearly undesirable (even unification is undecidable in full 2nd order logic). However, much weaker second order constructions already provide much if not all of the required expressivity without causing any computational problems (in effect, they are simply 2nd order syntactic sugar for what are essentially first order constructions). A precise characterization of such expressivity is required in a future extension of OIL. OIL is currently very restricted. Only classes are provided, not meta-classes or individuals.