

Graphs, foams, homology, polytopes, and abstract tensors

J. Scott Carter

University of South Alabama

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Acknowledgements

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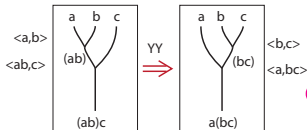
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Section 1

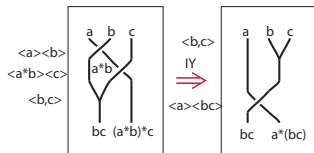
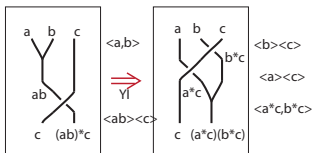
in which systems of abstract tensor equations are formulated and a solution is proposed.

3D moves

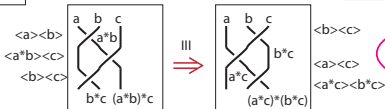
$$\langle a, b \rangle + \langle ab, c \rangle = \langle b, c \rangle + \langle a, bc \rangle$$



$$\langle a \rangle \langle b \rangle + \langle a^* b \rangle \langle c \rangle + \langle b, c \rangle = \langle b, c \rangle + \langle a \rangle \langle bc \rangle$$



$$\langle a, b \rangle + \langle ab \rangle \langle c \rangle = \langle b \rangle \langle c \rangle + \langle a \rangle \langle c \rangle + \langle a^* c, b^* c \rangle$$



$$\langle a \rangle \langle b \rangle + \langle a^* b \rangle \langle c \rangle = \langle a \rangle \langle c \rangle + \langle a^* c \rangle \langle b^* c \rangle$$

In 3-D equations

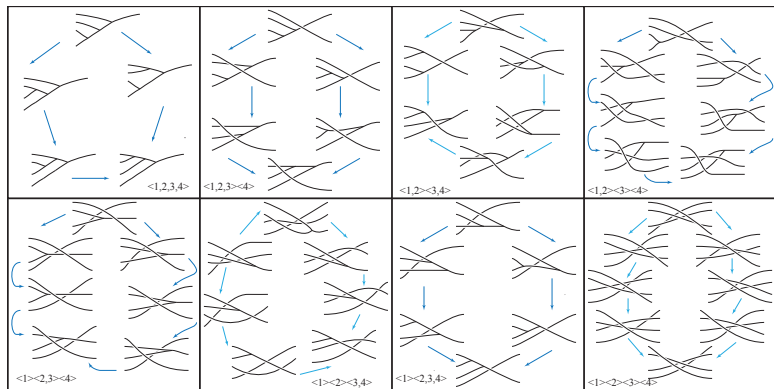
A	$Y_j^{\ell,n} Y_i^{j,k} = Y_p^{n,k} Y_i^{\ell,p}$
YI	$Y_d^{a,b} X_{e,f}^{d,c} = X_{c,h}^{b,c} X_{e,g}^{a,c} Y_f^{g,h}.$
IY	$X_{d,g}^{a,b} X_{h,f}^{g,c} Y_e^{d,h} = Y_i^{b,c} X_{e,f}^{a,i}$
III	$X_{d,e}^{a,b} X_{i,h}^{e,c} X_{f,g}^{d,i} = X_{j,k}^{b,c} X_{f,e}^{a,j} X_{g,h}^{e,k}$

Easy Solution

$$Y_c^{a,b} = \begin{cases} 1 & \text{if } c = ab, \\ 0 & \text{otherwise;} \end{cases}$$

$$X_{c,d}^{a,b} = \begin{cases} 1 & \text{if } c = b \text{ \& } d = b^{-1}ab, \\ 0 & \text{otherwise.} \end{cases}$$

4D moves



4-D Equations

YYY	$K_{p,q}^{m,n} K_{s,t}^{p,r} = K_{u,v}^{n,r} K_{s,w}^{m,u} K_{t,q}^{w,v}$
YYI	$M_{t,u,v}^{p,q} M_{x,y,z}^{u,r} K_{m,n}^{t,x} = K_{s,i}^{q,r} M_{m,y,j}^{p,s} M_{n,z,v}^{j,i}$
YY	$W_{i,j}^{p,q,r} M_{m,n,y}^{i,s} = M_{t,u,v}^{r,s} M_{m,x,z}^{q,t} W_{n,h}^{p,x,u} W_{y,i}^{h,z,v}$
YII	$\mathbb{X}_{t,x,v}^{p,q,r} M_{y,u,v}^{x,s} M_{m,n,z}^{t,i} = M_{j,x,g}^{r,s} M_{m,t,x}^{q,j} \mathbb{X}_{n,y,f}^{p,t,x} \mathbb{X}_{z,u,v}^{f,w,g}$
IYY	$K_{u,v}^{p,q} W_{i,j}^{v,r,s} W_{y,z}^{u,i,t} = W_{u,v}^{q,s,t} W_{y,i}^{p,r,u} K_{z,j}^{i,v}$
IYI	$M_{u,v,x}^{p,q} \mathbb{X}_{i,j,n}^{x,r,s} \mathbb{X}_{h,g,f}^{v,j,t} W_{y,z}^{u,i,h} = W_{u,v}^{q,s,t} \mathbb{X}_{y,i,j}^{p,r,u} M_{z,f,n}^{j,v}$
IYY	$W_{h,j}^{p,q,r} W_{i,n}^{j,s,t} \mathbb{X}_{x,y,z}^{h,i,u} = \mathbb{X}_{v,g,f}^{r,t,u} \mathbb{X}_{x,h,j}^{q,s,v} W_{y,i}^{p,h,g} W_{z,n}^{i,j,f}$
IIII	$\mathbb{X}_{v,x,y}^{p,q,r} \mathbb{X}_{z,e,n}^{y,s,t} \mathbb{X}_{k,h,j}^{x,e,u} \mathbb{X}_{g,f,i}^{v,z,k} = \mathbb{X}_{k,z,v}^{r,t,u} \mathbb{X}_{g,e,x}^{q,s,k} \mathbb{X}_{f,h,y}^{p,e,z} \mathbb{X}_{i,j,n}^{y,x,v}$

Solutions part 1

$$K_{r,s}^{p,q} = \begin{cases} 1 & \text{if } \exists a, b, c \text{ s.t.} \\ & \begin{cases} p = ((ab, c), 0|0), & q = ((a, b), 0|0), \\ r = ((a, bc), 0|0), & \& s = ((b, c), 0|0). \end{cases} \\ 0 & \text{otherwise;} \end{cases}$$

Solutions part 2

$$M_{p,q,r}^{i,j} = \begin{cases} 1 & \text{if } \exists a, b, c \text{ s.t.} \\ & \begin{cases} i = ((0, 0), ab|c), & j = ((a, b), 0|0), \\ p = ((a, b) \triangleleft c, 0|0), & q = ((0, 0), a|c), \\ \& r = ((0, 0), b|c), \end{cases} \\ 0 & \text{otherwise;} \end{cases}$$

Solutions part 3

$$W_{p,q}^{i,j,\ell} = \begin{cases} 1 & \text{if } \exists a, b, c \text{ s.t.} \\ & \begin{cases} i = ((b, c), 0|0), & j = ((0, 0), (a \triangleleft b)|c), \\ \ell = ((0, 0), a|b), \\ p = ((0, 0), a|(bc)), & \& q = ((b, c), 0|0), \end{cases} \\ 0 & \text{otherwise;} \end{cases}$$

Solutions part 4

$$\mathbb{X}_{s,t,u}^{p,q,r} = \begin{cases} 1 & \text{if } \exists a, b, c \text{ s.t.} \\ & \begin{cases} p = ((0, 0), b|c), & q = ((0, 0), (a \triangleleft b)|c), \\ r = ((0, 0), a|b), \\ s = ((0, 0), (a|b) \triangleleft c), & t = ((0, 0), a|c), \\ \& u = ((0, 0), b|c), \end{cases} \\ 0 & \text{otherwise;} \end{cases}$$

where, *e.g.*, $a \triangleleft c = c^{-1}ac$, $(a, b) \triangleleft c = (a \triangleleft c, b \triangleleft c)$,
and $(a|b) \triangleleft c = (a \triangleleft c)|(b \triangleleft c)$

Algebraic setting

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So when I write, *e.g.*

$$M_{p,q,r}^{i,j} = \begin{cases} 1 \text{ if } \exists a, b, c \text{ s.t.} \\ \quad \begin{cases} i = ((0, 0), ab|c), & j = ((a, b), 0|0), \\ p = ((a, b) \triangleleft c, 0|0), & q = ((0, 0), a|c), \\ \quad \& r = ((0, 0), b|c), \end{cases} \\ 0 \text{ otherwise;} \end{cases}$$

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the tensor M is a morph.

$$V \otimes V \xleftarrow{M} V \otimes V \otimes V.$$

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and

$$V \otimes V \otimes V \xleftarrow{X} V \otimes V \otimes V.$$

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Section 2

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The Space Y^n

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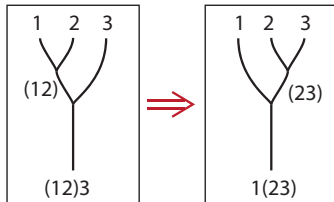
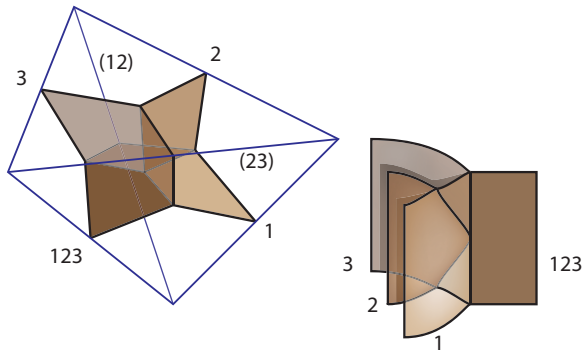
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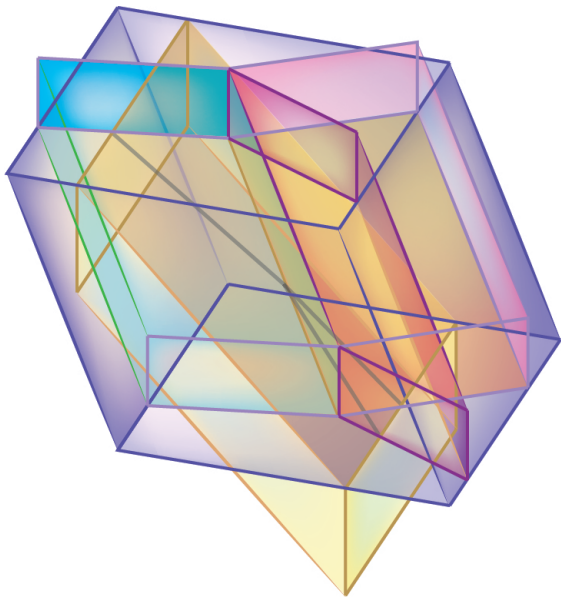
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Take $\Delta_j^n = \{\vec{x} \in \Delta^{n+1} : x_j = 0\}$. Embed a copy,
 $Y_j^{n-1} \subset \Delta_j^n$. Cone $\cup_{j=1}^{n+2} Y_j^{n-1}$ to the barycenter
 $b = \frac{1}{n+2}(1, 1, \dots, 1)$ of Δ^{n+1} .

$$Y^n = C \left(\cup_{j=1}^{n+2} Y_j^{n-1} \right).$$

The space Y^2



The space Y^3



Foam Definition

An n -**dimensional foam** is a compact top. sp. X for which each point $x \in X$ has a nbhd. $N(x)$ that is homeom. to a nbhd. M of a point in Y^n .

Local pictures of crossings

Now take

$$\left[\bigcup_{\ell=1}^k (\Delta^{j_1} \times \cdots \times Y^{j_{\ell-1}} \times \cdots \times \Delta^{j_k}) \subset \mathbb{R}^{n+1} \times \{\ell\} \right],$$

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and project this into \mathbb{R}^{n+1} . The factor ℓ is in the $(n+2)$ nd coordinate and represents the relative height of each Y .

$$\left[(Y^{j_1-1} \times \Delta^{j_2} \cdots \times \Delta^{j_\ell} \times \cdots \times \Delta^{j_k}) \subset \mathbb{R}^{n+1} \times \{1\} \right],$$

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These project to a 0-dimensional multiple point
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Section 3

in which a discussion of categorification is used to justify the ideas presented in the previous slide.

New school categorification

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Instead of equality among morphisms, posit 2-morphisms that satisfy their own set of relations. Climb the dimension ladder.

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$$\vec{x} = \sum_{j=1}^n x^j e_j \quad \text{where} \quad e_j = \begin{bmatrix} 0 \\ \vdots \\ 0 \\ 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix} \leftarrow \text{jth} \left. \begin{array}{l} \text{ } \\ \text{ } \\ \text{ } \\ \text{ } \\ \text{ } \\ \text{ } \\ \text{ } \end{array} \right\} \text{n rows}$$

Abstract
Tensor
Formalism

so superscripts are row indices.

Suppose V, W , etc. $f.\dim'l$ vec. sp. over F .

$$\vec{x} = \sum_{j=1}^n x^j e_j \quad \text{where} \quad e_j = \begin{bmatrix} 0 \\ \vdots \\ 0 \\ 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix} \leftarrow \begin{matrix} j\text{th} \\ \left. \vphantom{\begin{matrix} 0 \\ \vdots \\ 0 \\ 1 \\ 0 \\ \vdots \\ 0 \end{matrix}} \right\} n \text{ rows} \end{matrix}$$

Abstract
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so superscripts are row indices.

Write $\vec{x} = \boxed{x}$.

Suppose V, W , etc. $f.\dim'l$ vec. sp. over F .

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Abstract
Tensor
Formalism

so superscripts are row indices.

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Composition of linear maps is denoted by vertical juxtaposition.

$$T \xleftarrow{C} U \xleftarrow{B} W \xleftarrow{A} V \quad \begin{array}{|c} \boxed{C} \\ \boxed{B} \\ \boxed{A} \end{array} \begin{array}{l} \uparrow \\ T \\ V \end{array}$$

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$C(B(A(\vec{x})))$

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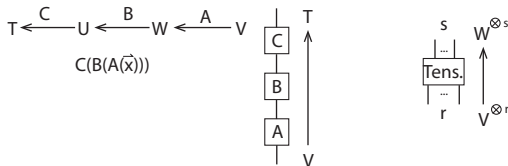
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Tensor
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Example: A Frobenius alg. is a v. sp. V tog. w/ linear maps

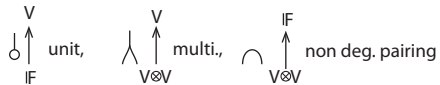
Example: A Frobenius alg. is a v. sp. V tog. w/ linear maps

$$\begin{array}{c} V \\ \uparrow \\ \circ \quad | \quad \text{unit,} \\ \text{IF} \end{array}$$

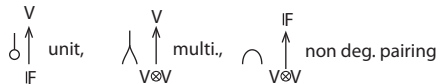
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$$\begin{array}{c} V \\ \circ \uparrow \\ \text{IF} \end{array} \text{ unit,} \quad \begin{array}{c} V \\ \wedge \uparrow \\ V \otimes V \end{array} \text{ multi.,}$$

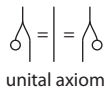
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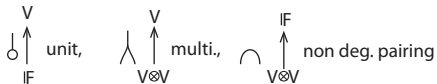
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$$\begin{array}{c} \circ \quad | \\ \wedge \\ \text{unit axiom} \end{array} = \begin{array}{c} | \\ \wedge \\ \circ \end{array} \quad \text{here } \begin{array}{c} | \\ \text{denotes the identity map} \\ | \\ V \\ \uparrow \\ V \end{array}$$

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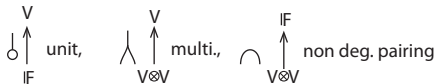


unital axiom



associativity

Example: A Frobenius alg. is a v. sp. V tog. w/ linear maps



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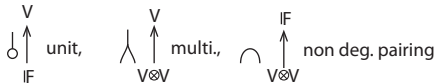


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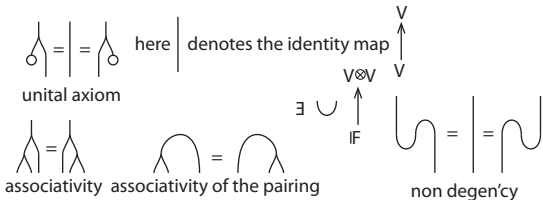


associativity associativity of the pairing

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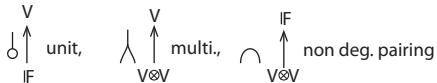
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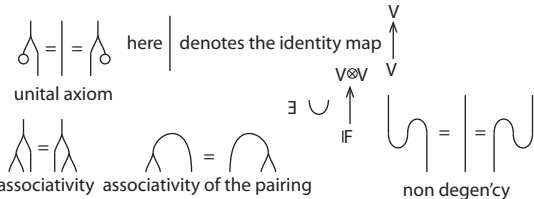
$$\begin{array}{c} \wedge = \wedge \\ \text{associativity} \end{array} \quad \begin{array}{c} \cup = \cup \\ \text{associativity of the pairing} \end{array} \quad \exists \quad \begin{array}{c} \cup \\ \uparrow \\ \text{IF} \end{array} \quad \begin{array}{c} \cup = | = \cup \\ \text{non degen'cy} \end{array}$$

We can use \cup and \cap to define comulti.

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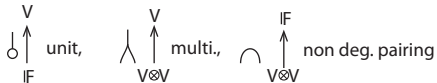
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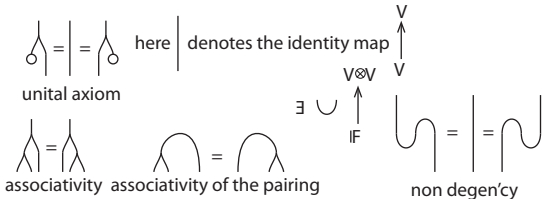
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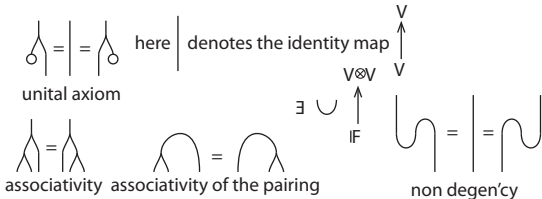
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It's remarkable that most of the axiomatics for the alg.coalg str. follows directly from the diagrammatics.

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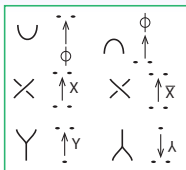
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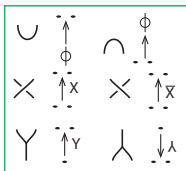
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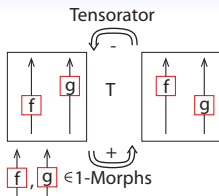
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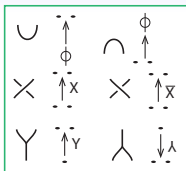


Generating 1-morphs

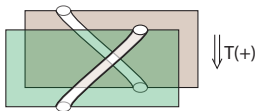
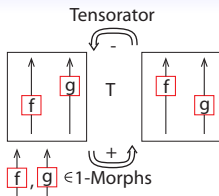


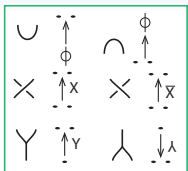
Generating 1-morphs



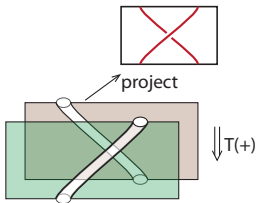
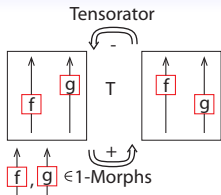


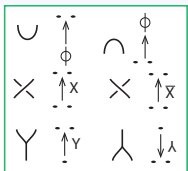
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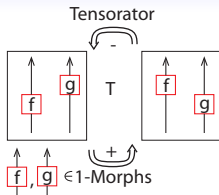


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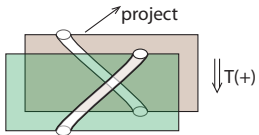


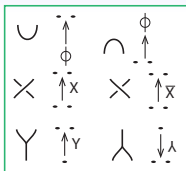
Tensorator

Relations (3-morphism)

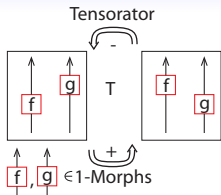


project

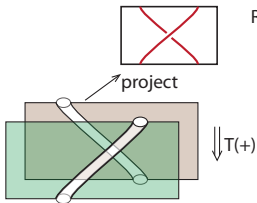
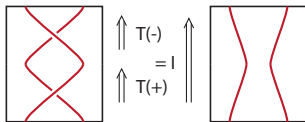


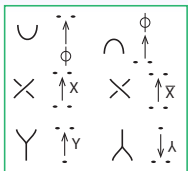


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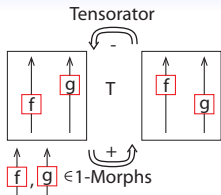


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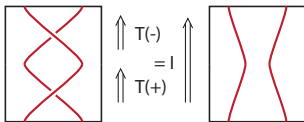


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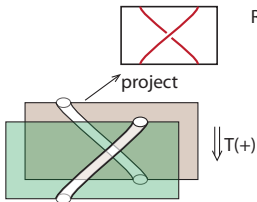


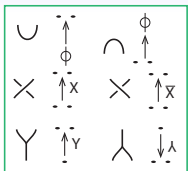
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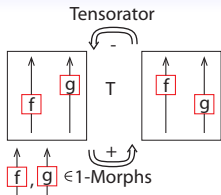


ditto for $T(+)$, $T(-)$

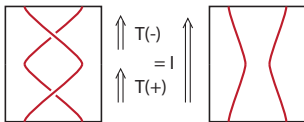




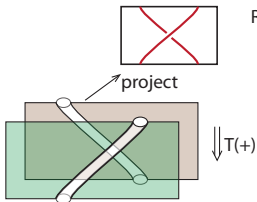
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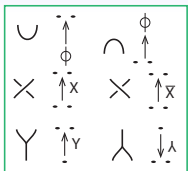
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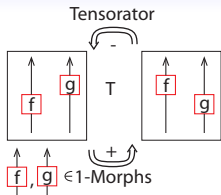
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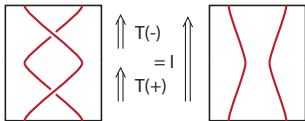
$T(\pm)$ is natural w.r.t. any 2-morph.



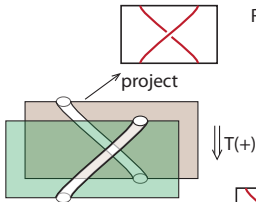
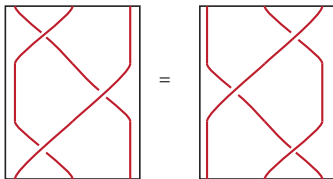
Generating 1-morphisms



Relations (3-morphism)

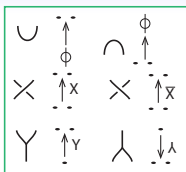


ditto for $T(+)\circ T(-)$

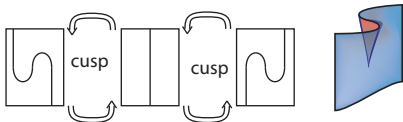
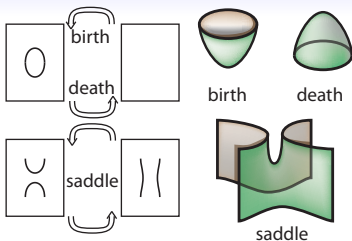


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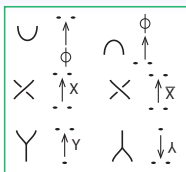
In part.:



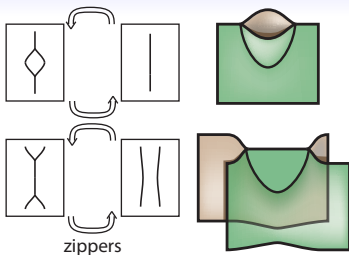
Generating 1-morphs



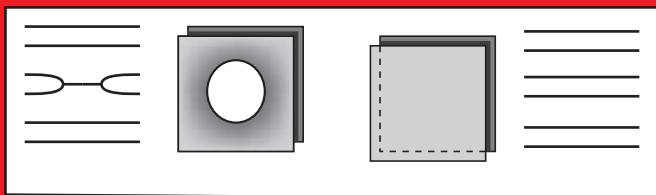
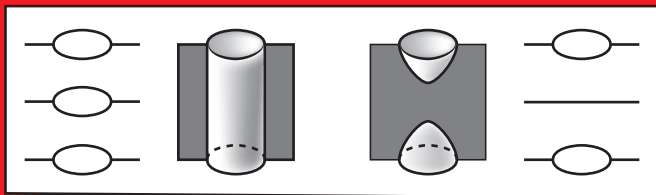
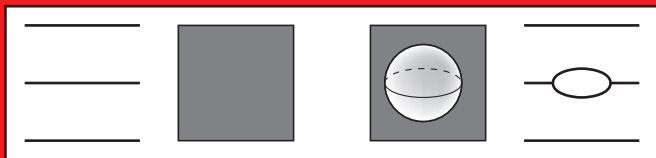
- The critical points evolve to be come folds.
(co-oriented away from optimal points)
- Several obvious relations hold among these 2-morphisms.
 - Including
 - canceling birth/saddle death/saddle
 - lips
 - beak-to-beak
 - swallow-tail
 - horizontal cusp

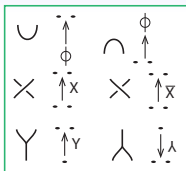


Generating 1-morphs

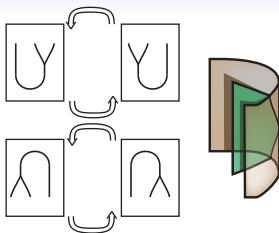


- The vertices of Y or λ evolve to form seams of the foam.
(co-oriented towards the single sheet)
- There are zig-zag moves that cancel a pair of zipper 2-morphisms.
- Under some circumstances, one might want to suppose bubble and saddle moves hold.
But, as is the case with birth followed by death,
or a pair of opposite saddles,
it is better to suppose that the moves
in the next slide are some type of 3-morphisms.

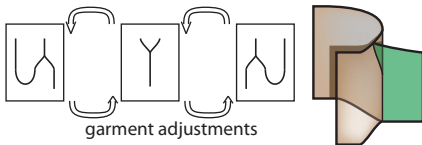




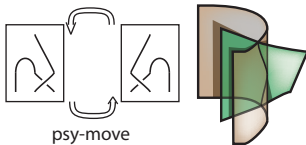
Generating 1-morphs



speedometers

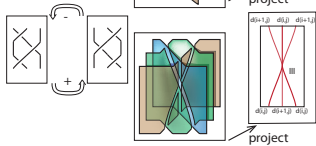
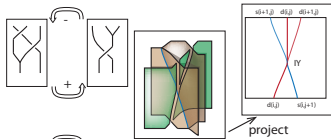
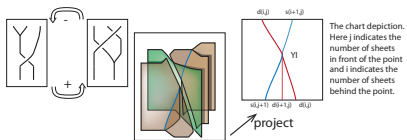
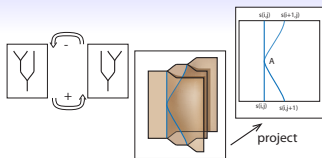


garment adjustments

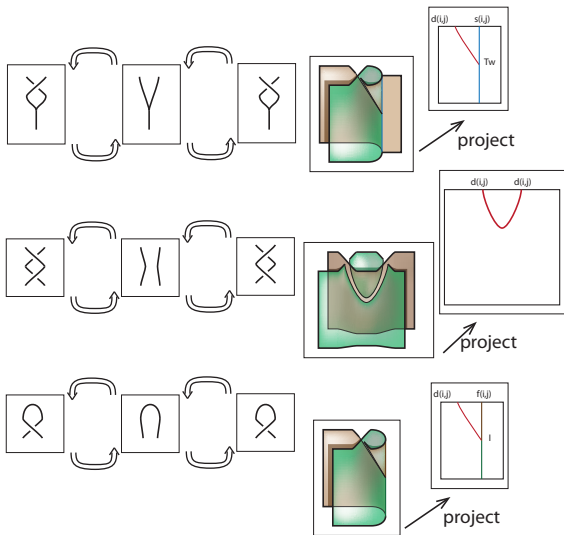


psy-move

- These 2-morphs indicate how folds, seams, double points can interact.
- Each move is invertible in both directions.



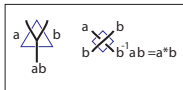
The broken surface representation of the resulting foam.



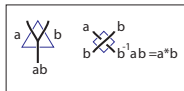
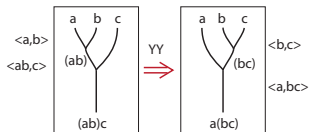
Section 4

in which an homology theory that encompasses both quandle and group homology is outlined.

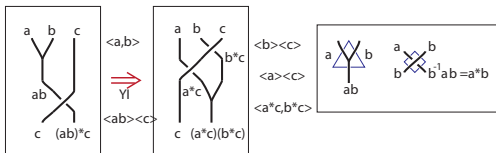
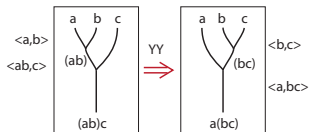
Fundamental group



$$(ab)c = a(bc)$$

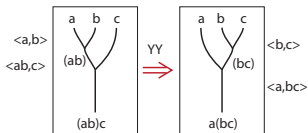


$$(ab)c = a(bc)$$

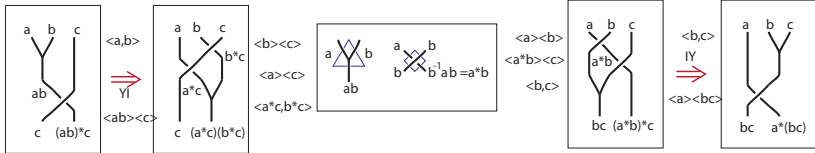


$$(ab)^*c = (a^*c)(b^*c)$$

$$(ab)c = a(bc)$$

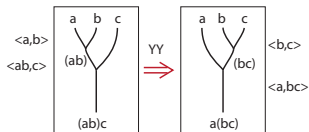


$$(a^*b)^*c = a^*(bc)$$

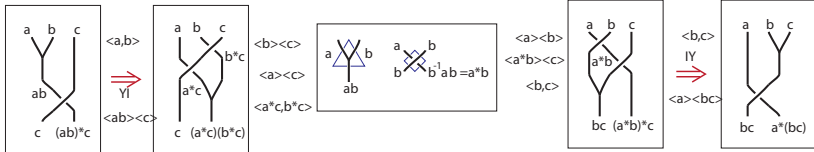


$$(ab)^*c = (a^*c)(b^*c)$$

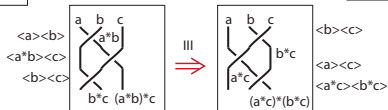
$$(ab)c = a(bc)$$



$$(a^*b)^*c = a^*(bc)$$



$$(ab)^*c = (a^*c)(b^*c)$$



$$(a^*b)^*c = (a^*c)^*(b^*c)$$

$$\text{YY} \quad (ab)c = a(bc),$$

$$\text{IY} : \quad (ab) \triangleleft c = (a \triangleleft c)(b \triangleleft c),$$

$$\text{YI} : \quad (a \triangleleft b) \triangleleft c = a \triangleleft (bc),$$

$$\text{III} : \quad (a \triangleleft b) \triangleleft c = (a \triangleleft c) \triangleleft (b \triangleleft c).$$

A quandle

satisfies three axioms that correspond to the Reidemeister moves:

$$I : \quad (\forall a) : \quad a \triangleleft a = a$$

$$II : \quad (\forall a, b)(\exists c) : \quad c \triangleleft b = a$$

$$III : \quad (\forall a, b, c) : \quad (a \triangleleft b) \triangleleft c = (a \triangleleft c) \triangleleft (b \triangleleft c).$$

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satisfies three axioms that correspond to the Reidemeister moves:

$$I : \quad (\forall a) : \quad a \triangleleft a = a$$

$$II : \quad (\forall a, b)(\exists c) : \quad c \triangleleft b = a$$

$$III : \quad (\forall a, b, c) : \quad (a \triangleleft b) \triangleleft c = (a \triangleleft c) \triangleleft (b \triangleleft c).$$

We are interested in how the group G and its associated quandle $\text{Conj}(G)$ interact.

Remark

There are related concepts for which the homology sketched below applies, *e.g.*:

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- G -families of quandles (IIJO)

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- G -families of quandles (IIJO)
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- Lebed's qualgebras

Homology of G -families of quandles was defined to study handle-body knots.

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Here we use, YY , YI , IY , and III to define the homological conditions.

Slicing

Cut the interval $[0, n]$ into integral pieces.

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$$\langle 1, 2, \dots, \ell_1 \rangle \langle \ell_1 + 1, \dots, \ell_1 + \ell_2 \rangle \cdots \\ \left\langle \sum_{i=1}^{j-1} \ell_i + 1, \dots, \sum_{i=1}^j \ell_i \right\rangle \cdots \left\langle \sum_{i=1}^k \ell_i + 1, \dots, n \right\rangle.$$

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Such a slice corresponds to a decomposition of the n -ball into a product of simplices. There are 2^{n-1} ways to cut.

Boundaries

$$\partial\langle j + 1, j + 2, \dots, j + k \rangle$$

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$$\begin{aligned} \partial\langle j + 1, j + 2, \dots, j + k \rangle \\ = \triangleleft(j + 1)\langle j + 2, \dots, j + k \rangle \end{aligned}$$

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$$\partial(PQ) = (\partial P)Q + (-1)^{\dim P} P(\partial Q).$$

Boundaries

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$$\partial(PQ) = (\partial P)Q + (-1)^{\dim P} P(\partial Q).$$

In part,

$$\partial \langle j+1 \rangle = \triangleleft (j+1) \lrcorner - \lrcorner.$$

Following Przytycki, one can observe that $\partial \circ \partial = 0$ in this context, if and only if

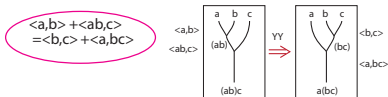
- $a(bc) = (ab)c$
- $a \triangleleft (bc) = (a \triangleleft b) \triangleleft c$
- $(ab) \triangleleft c = (a \triangleleft c)(b \triangleleft c)$
- $(a \triangleleft b) \triangleleft c = (a \triangleleft c) \triangleleft (b \triangleleft c)$

Sample computations 1

$$\begin{aligned}\partial\langle 1, 2, 3 \rangle &= \langle 2, 3 \rangle - \langle 1 \cdot 2, 3 \rangle \\ &\quad + \langle 1, 2 \cdot 3 \rangle - \langle 1, 2 \rangle\end{aligned}$$

Sample computations 1

$$\partial\langle 1, 2, 3 \rangle = \langle 2, 3 \rangle - \langle 1 \cdot 2, 3 \rangle \\ + \langle 1, 2 \cdot 3 \rangle - \langle 1, 2 \rangle$$

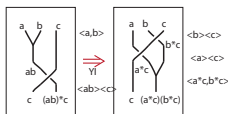


Sample computations 2

$$\begin{aligned}\partial\langle 1, 2 \rangle \langle 3 \rangle &= \partial\langle 1, 2 \rangle \langle 3 \rangle + \langle 1, 2 \rangle \triangleleft (3) - \langle 1, 2 \rangle \\ &= \langle 2 \rangle \langle 3 \rangle - \langle 1 \cdot 2 \rangle \langle 3 \rangle + \langle 1 \rangle \langle 3 \rangle \\ &\quad + \langle 1 \triangleleft 3, 2 \triangleleft 3 \rangle - \langle 1, 2 \rangle\end{aligned}$$

Sample computations 2

$$\begin{aligned}
 \partial\langle 1, 2 \rangle \langle 3 \rangle &= \partial\langle 1, 2 \rangle \langle 3 \rangle + \langle 1, 2 \rangle \triangleleft (3) - \langle 1, 2 \rangle \\
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 &\quad + \langle 1 \triangleleft 3, 2 \triangleleft 3 \rangle - \langle 1, 2 \rangle
 \end{aligned}$$



$$\begin{aligned}
 &\langle a, b \rangle + \langle ab \rangle \langle c \rangle \\
 &= \langle b \rangle \langle c \rangle + \langle a \rangle \langle c \rangle \\
 &\quad + \langle a^*c, b^*c \rangle
 \end{aligned}$$

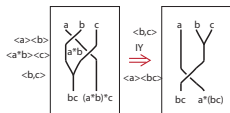
Sample computations 3

$$\begin{aligned}\partial\langle 1 \rangle \langle 2, 3 \rangle &= \langle 2, 3 \rangle - \langle 2, 3 \rangle - \langle 1 \rangle \partial \langle 2, 3 \rangle \\ &= -\langle 1 \triangleleft 2 \rangle \langle 3 \rangle + \langle 1 \rangle \langle 2 \cdot 3 \rangle - \langle 1 \rangle \langle 2 \rangle\end{aligned}$$

Sample computations 3

$$\begin{aligned}
 \partial\langle 1 \rangle \langle 2, 3 \rangle &= \langle 2, 3 \rangle - \langle 2, 3 \rangle - \langle 1 \rangle \partial \langle 2, 3 \rangle \\
 &= -\langle 1 \triangleleft 2 \rangle \langle 3 \rangle + \langle 1 \rangle \langle 2 \cdot 3 \rangle - \langle 1 \rangle \langle 2 \rangle
 \end{aligned}$$

$$\begin{aligned}
 \langle a \rangle \langle b \rangle + \langle a * b \rangle \langle c \rangle \\
 + \langle b, c \rangle = \langle b, c \rangle + \langle a \rangle \langle bc \rangle
 \end{aligned}$$

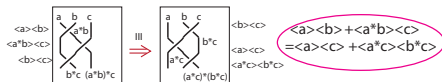


Sample computations

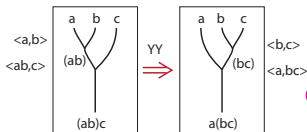
$$\begin{aligned}\partial\langle 1\rangle\langle 2\rangle\langle 3\rangle &= \langle 2\rangle\langle 3\rangle - \langle 2\rangle\langle 3\rangle - \langle 1\rangle\partial(\langle 2\rangle\langle 3\rangle) \\ &= -\langle 1\triangleleft 2\rangle\langle 3\rangle + \langle 1\rangle\langle 3\rangle \\ &\quad +\langle 1\triangleleft 3\rangle\langle 2\triangleleft 3\rangle - \langle 1\rangle\langle 2\rangle\end{aligned}$$

Sample computations

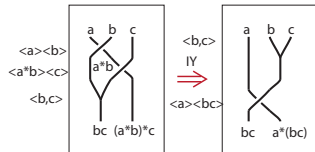
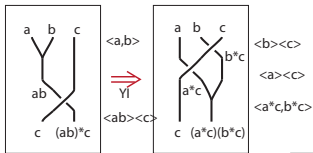
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 \partial\langle 1 \rangle \langle 2 \rangle \langle 3 \rangle &= \langle 2 \rangle \langle 3 \rangle - \langle 2 \rangle \langle 3 \rangle - \langle 1 \rangle \partial(\langle 2 \rangle \langle 3 \rangle) \\
 &= -\langle 1 \triangleleft 2 \rangle \langle 3 \rangle + \langle 1 \rangle \langle 3 \rangle \\
 &\quad + \langle 1 \triangleleft 3 \rangle \langle 2 \triangleleft 3 \rangle - \langle 1 \rangle \langle 2 \rangle
 \end{aligned}$$



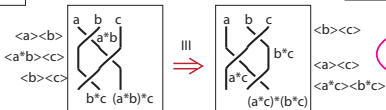
$$\langle a, b \rangle + \langle ab, c \rangle = \langle b, c \rangle + \langle a, bc \rangle$$



$$\langle a \rangle \langle b \rangle + \langle a^* b \rangle \langle c \rangle + \langle b, c \rangle = \langle b, c \rangle + \langle a \rangle \langle bc \rangle$$

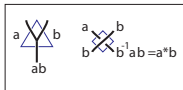


$$\langle a, b \rangle + \langle ab \rangle \langle c \rangle = \langle b \rangle \langle c \rangle + \langle a \rangle \langle c \rangle + \langle a^* c, b^* c \rangle$$

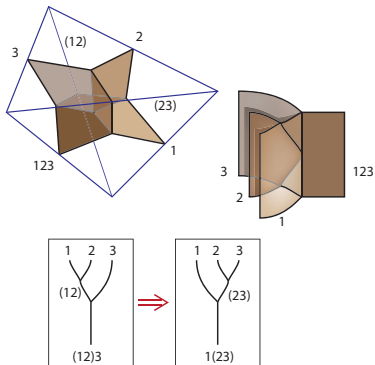


$$\langle a \rangle \langle b \rangle + \langle a^* b \rangle \langle c \rangle = \langle a \rangle \langle c \rangle + \langle a^* c \rangle \langle b^* c \rangle$$

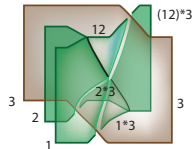
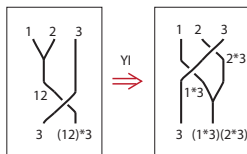
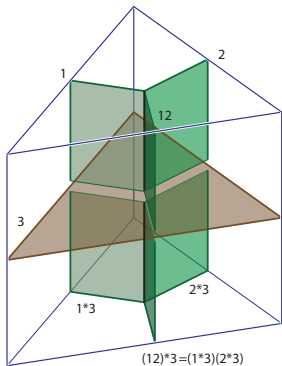
Triangles and Squares



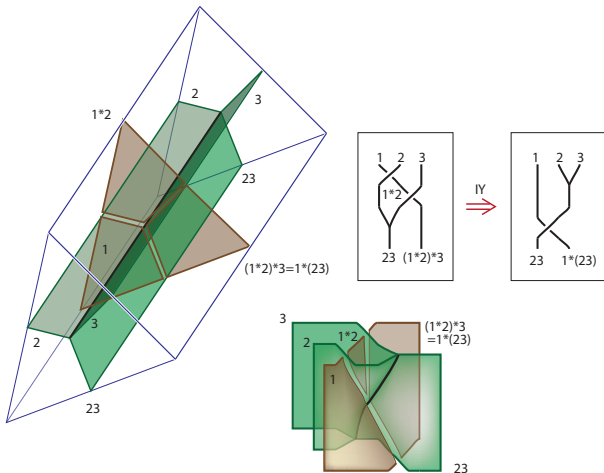
Tetrahedron



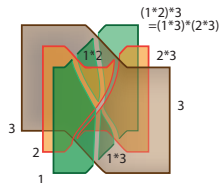
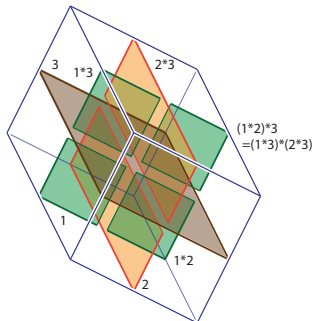
First Prism



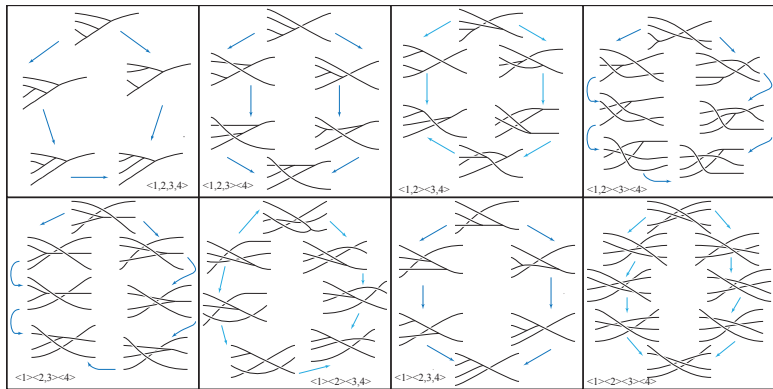
Second Prism



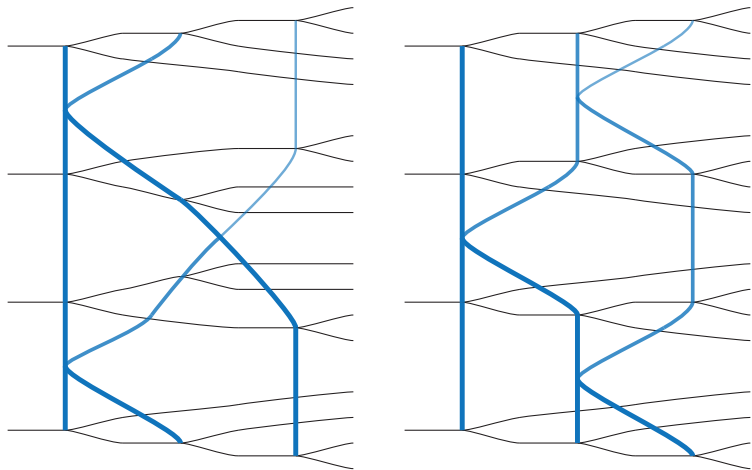
Cube



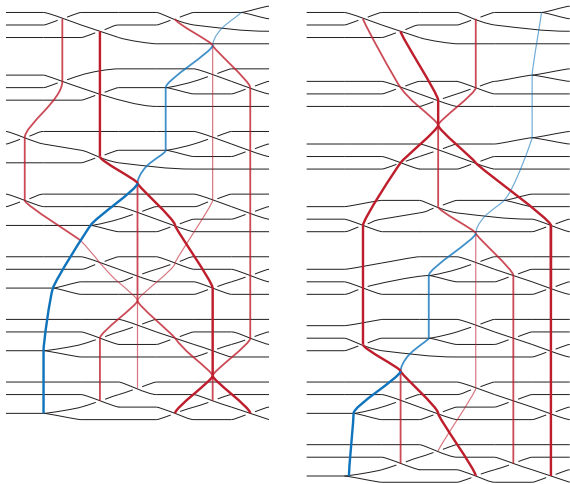
8 interesting moves



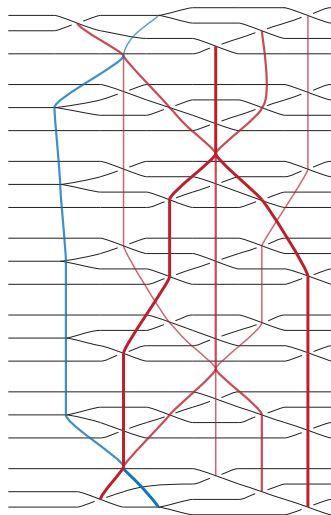
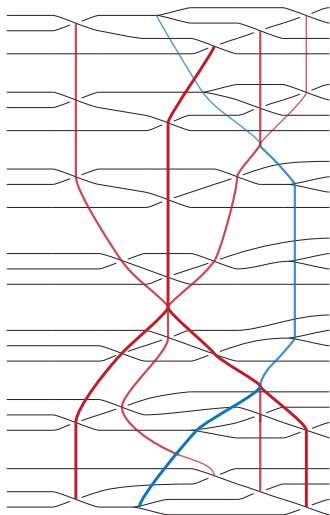
The YYY-move (Stasheff polytope)



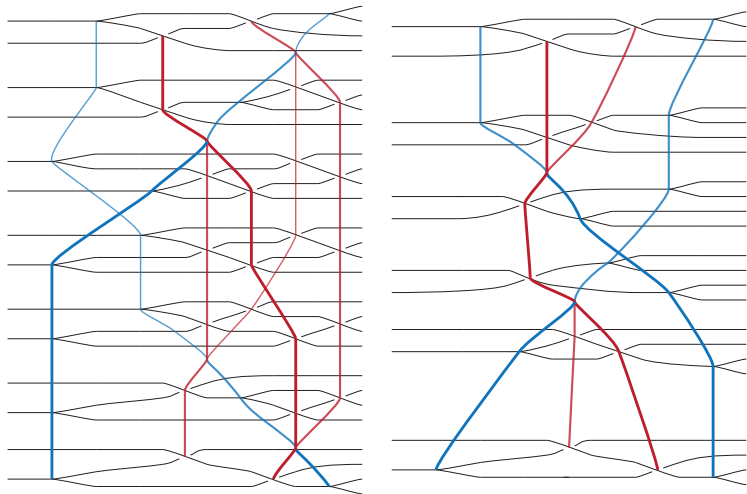
The YII-move



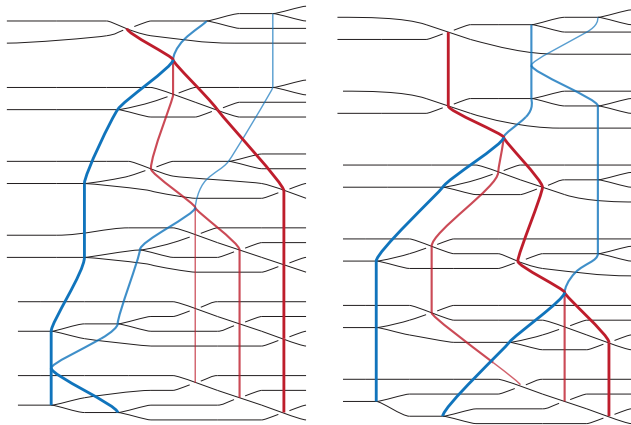
The IYI-move



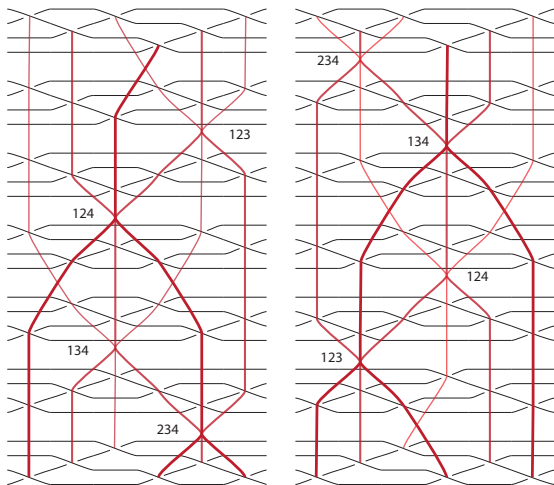
The YY-move



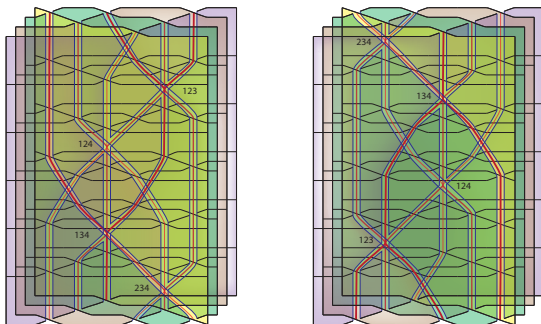
The YYI-move



The tetrahedral-move



The tetrahedral-move



Visual interlude

in which many (not all) of the Reidemeister/Roseman moves for 2-foams are indicated.

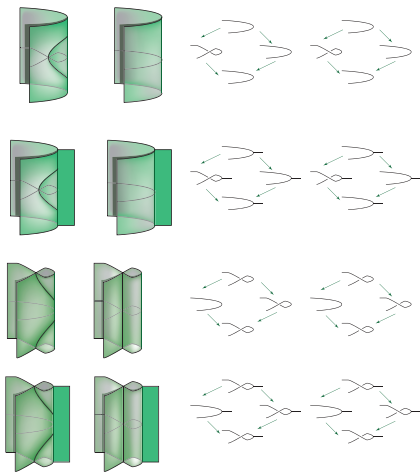
Visual interlude

in which many (not all) of the Reidemeister/Roseman moves for 2-foams are indicated. Categorical hint:

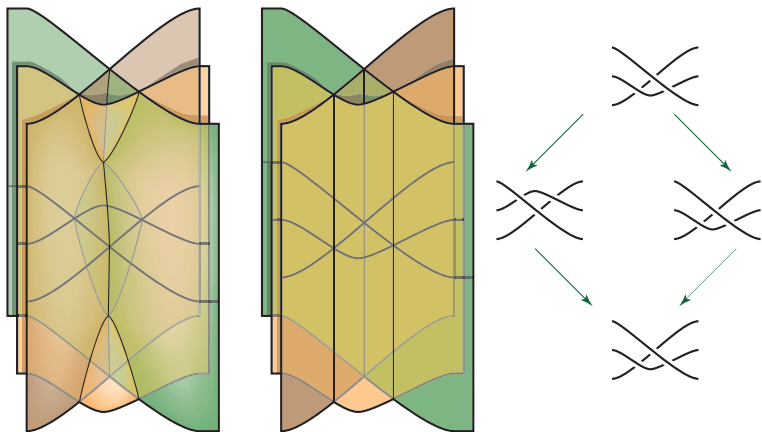
Visual interlude

in which many (not all) of the Reidemeister/Roseman moves for 2-foams are indicated. Categorical hint: when certain 2-morphisms are defined to be invertible, these moves arise.

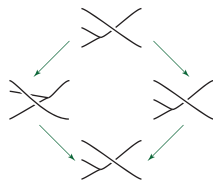
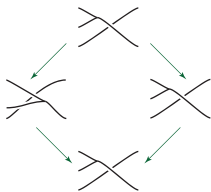
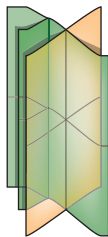
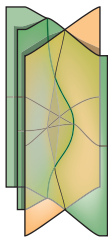
Critical points of the branch point set



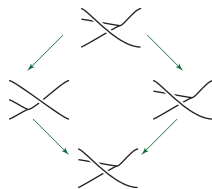
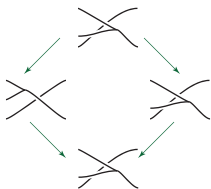
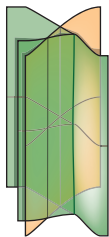
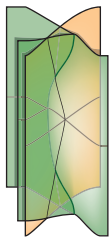
Critical points of the triple point set



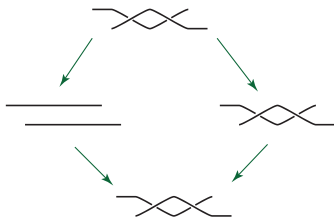
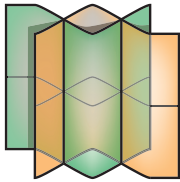
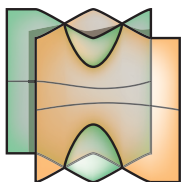
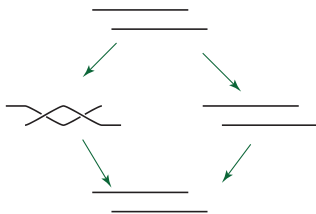
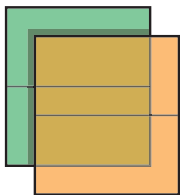
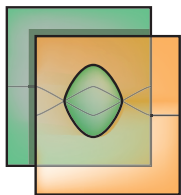
Critical points of the intersection set 1



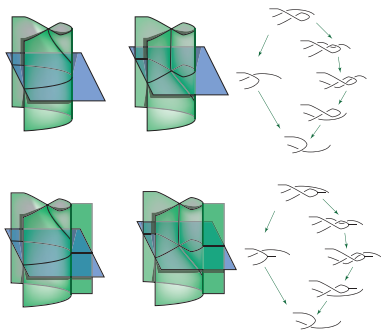
Critical points of the intersection set 2



Critical points of the double point set



Int. pts. $b/2$ branch/twist set and trnsvs.
sheet



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The analogues in one higher dimensions

of the generating chains in homology.

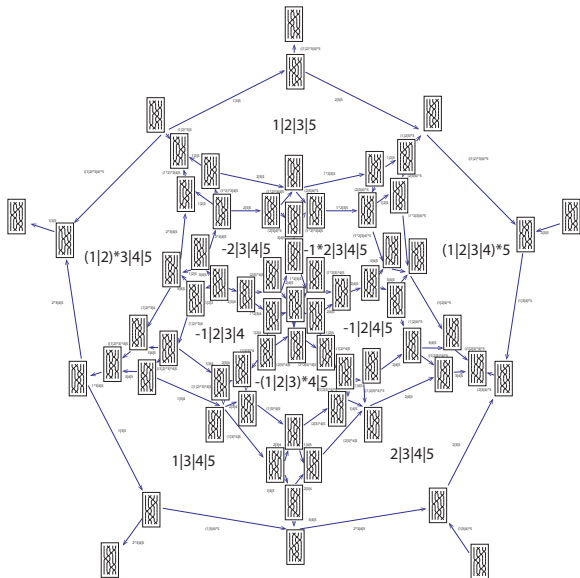
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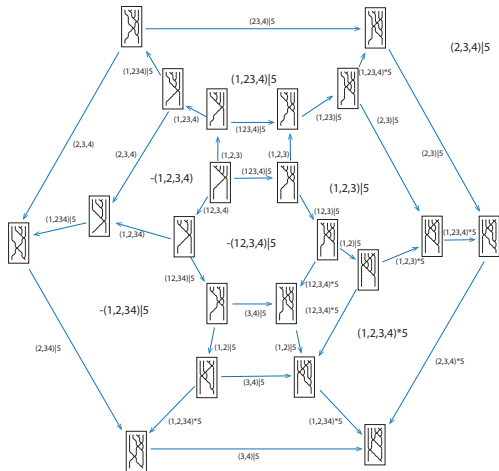
The analogues in one higher dimensions

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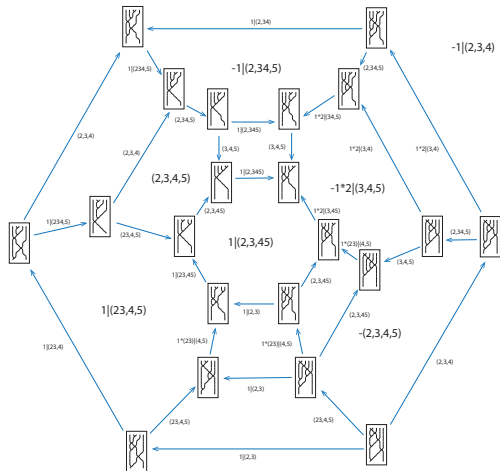
1|2|3|4|5



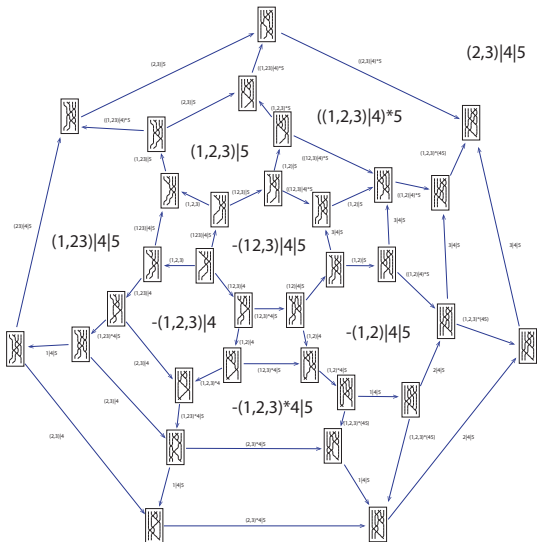
1234|5



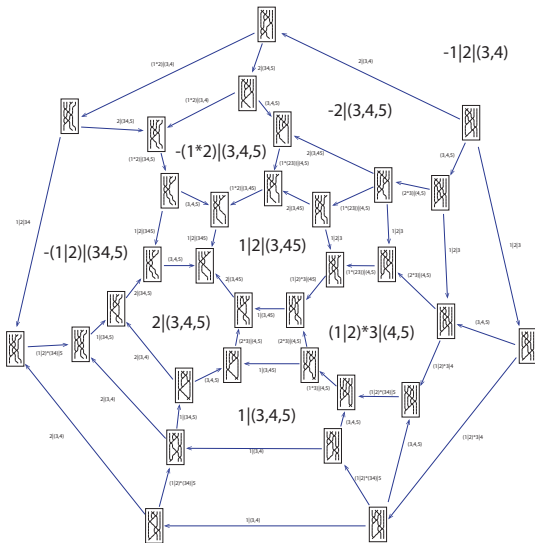
1|2345



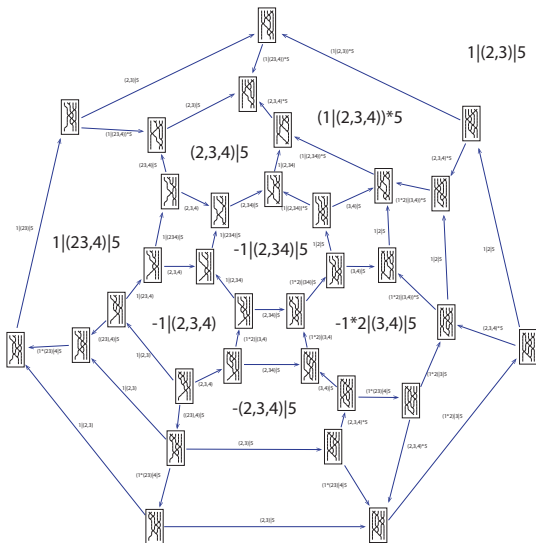
123|4|5



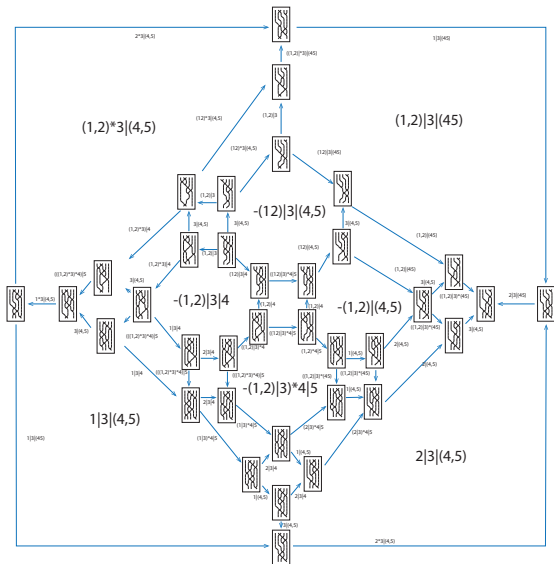
1|2|345



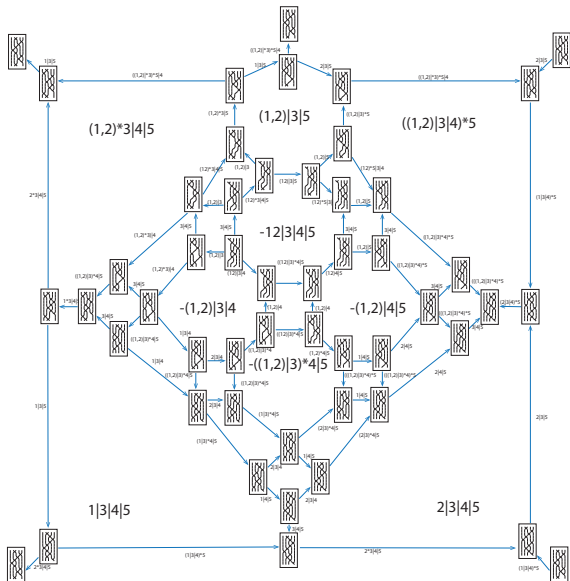
1|234|5



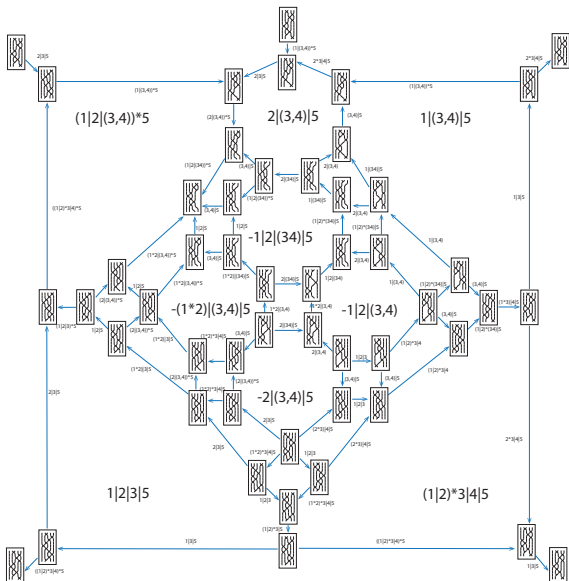
12|3|45



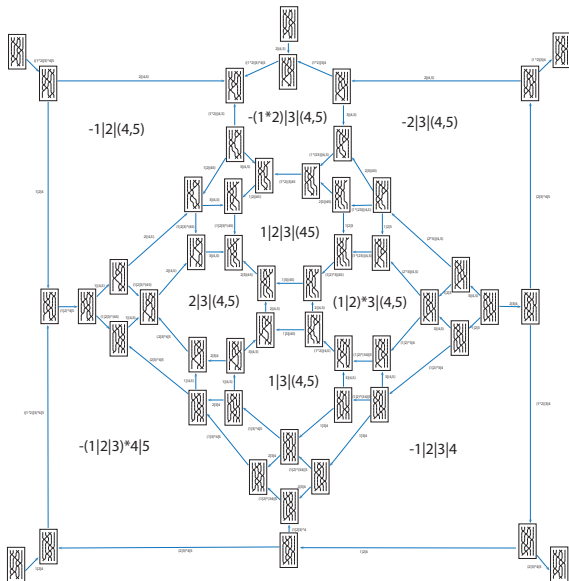
12|3|4|5



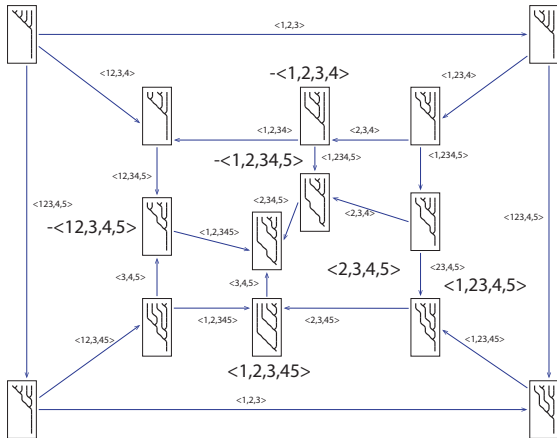
1|2|34|5



1|2|3|45



12345



Section 5

in which the polytopal duals to the generating chains are formulated.

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Permutahedron

In \mathbb{R}^n consider the convex hull of

$$\{(\sigma(1), \sigma(2), \dots, \sigma(n)) : \sigma \in \Sigma_n\}.$$

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$i = (i, i + 1)$. Hexagonal faces are $i(i + 1)i(i + 1)i(i + 1)$.

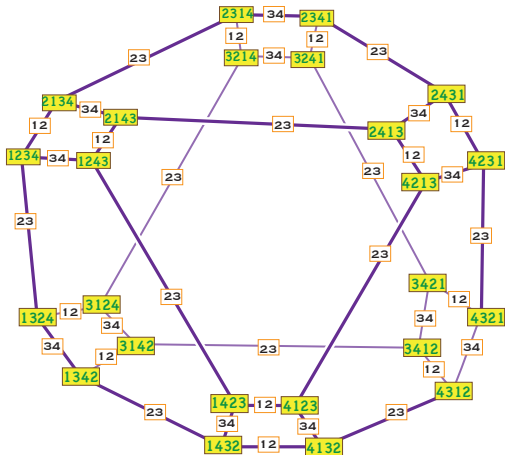
Permutahedron

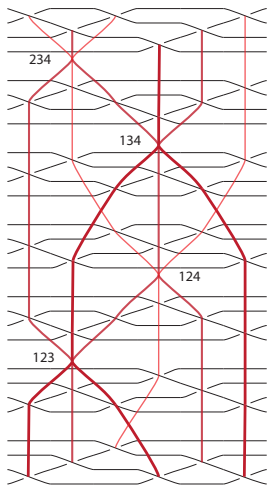
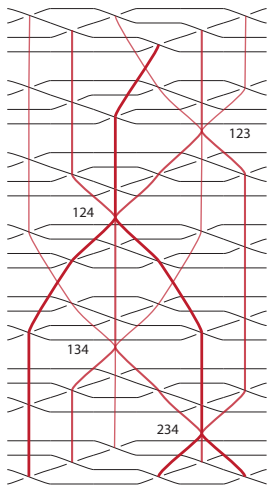
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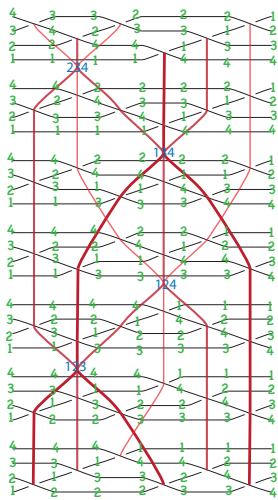
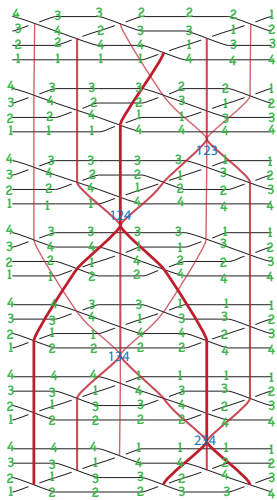
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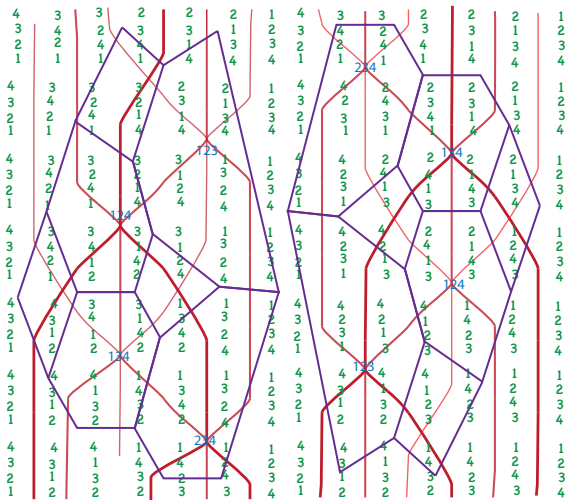
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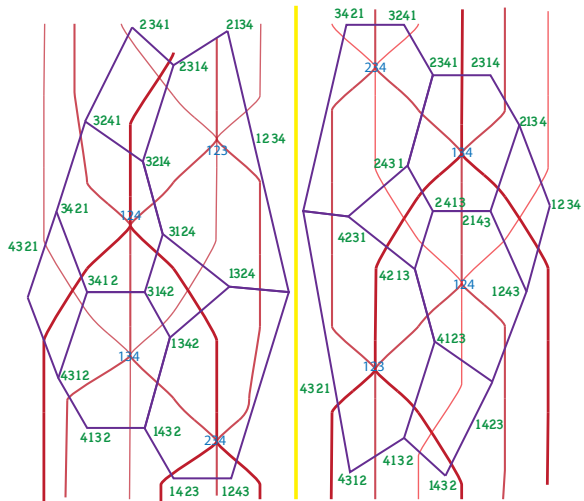
$\mathbf{i} = (i, i + 1)$. Hexagonal faces are $\mathbf{i}(\mathbf{i} + 1)\mathbf{i}(\mathbf{i} + 1)\mathbf{i}(\mathbf{i} + 1)$. *etc.*

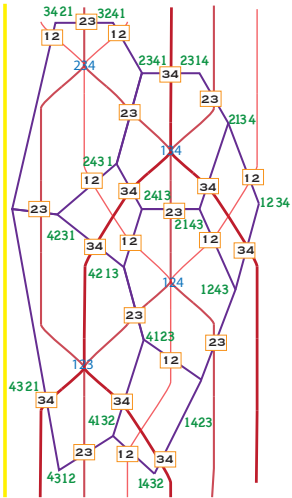
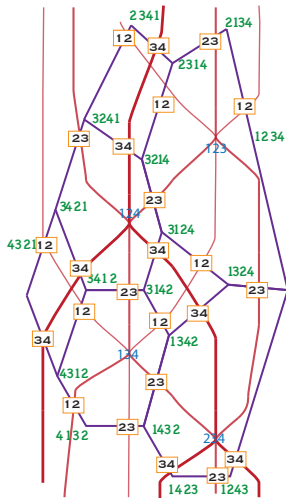


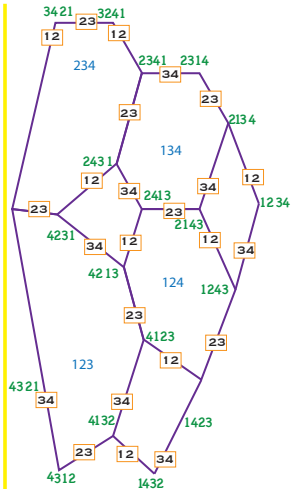
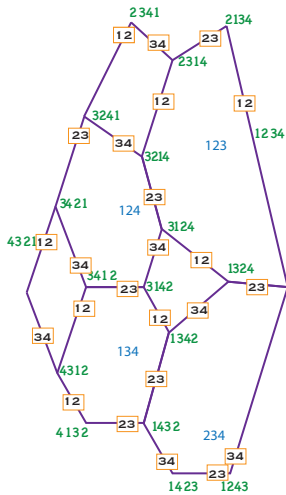


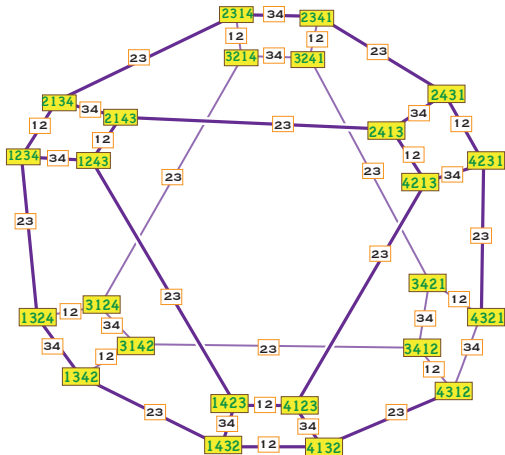




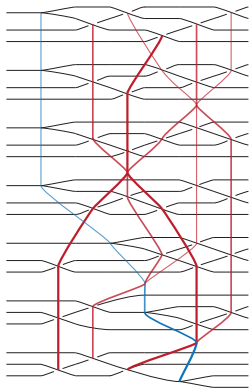




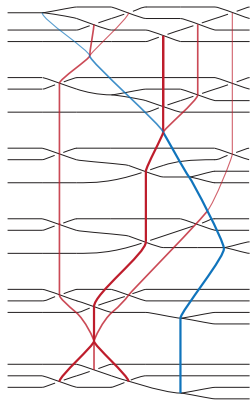


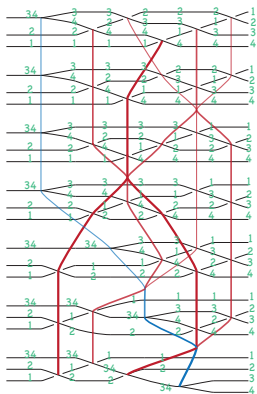


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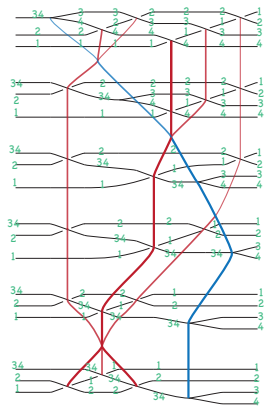


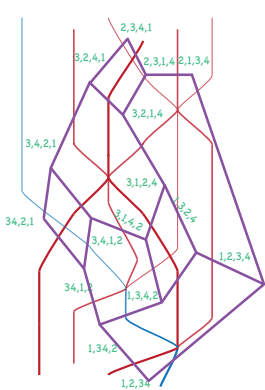
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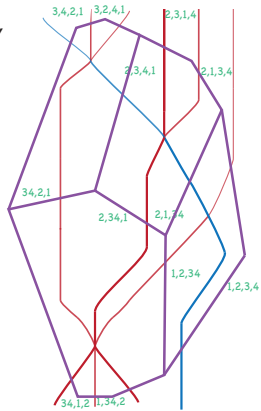


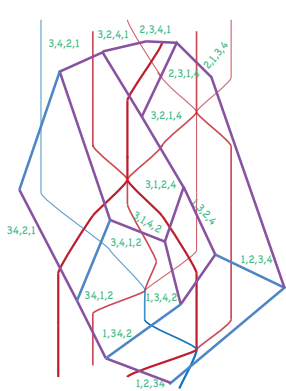
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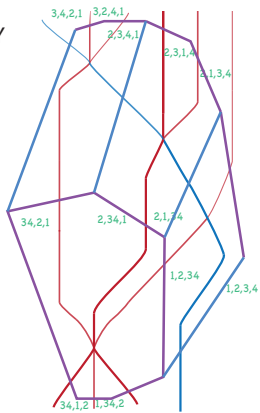


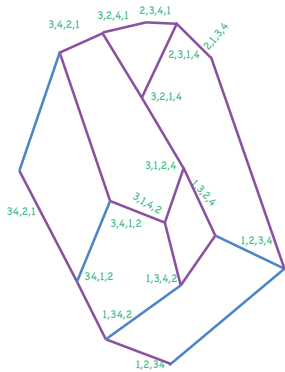
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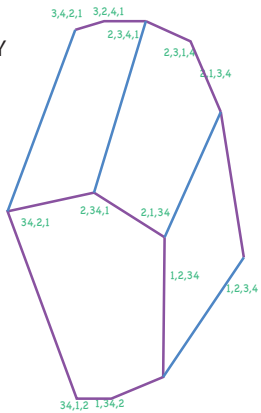


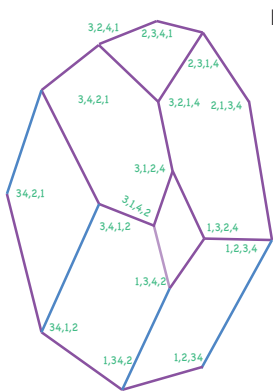
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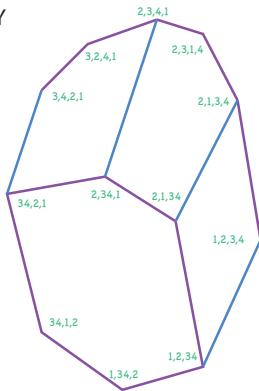


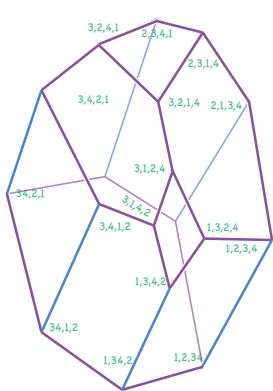
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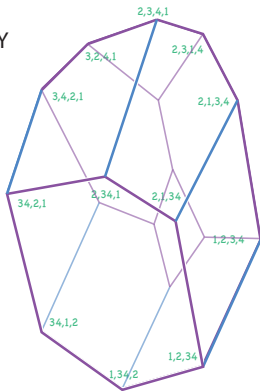


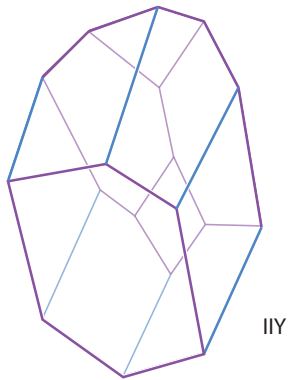
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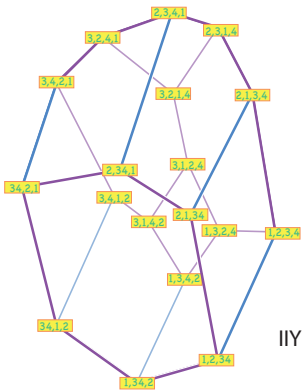




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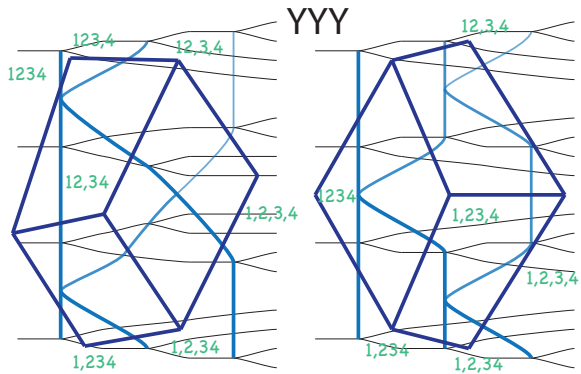
Before being in Asia, I thought that these polytopes were new.

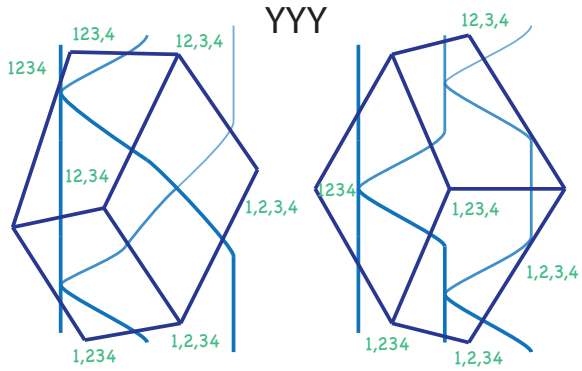
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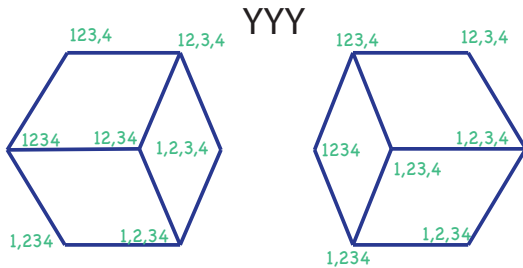
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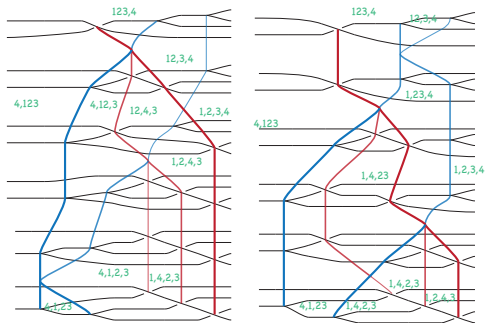


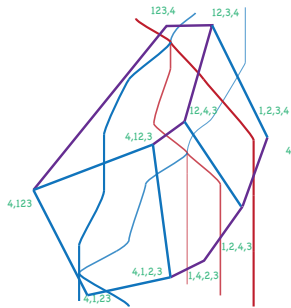




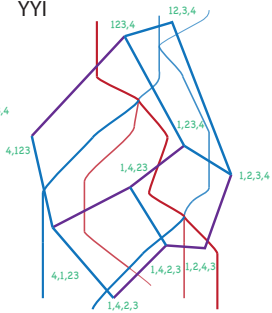
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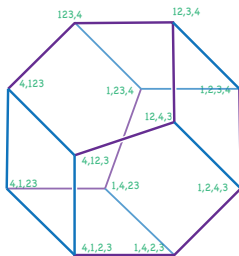
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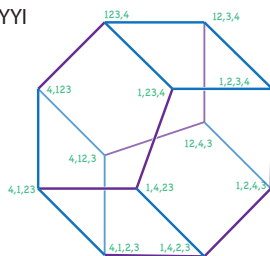


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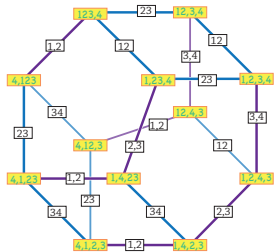




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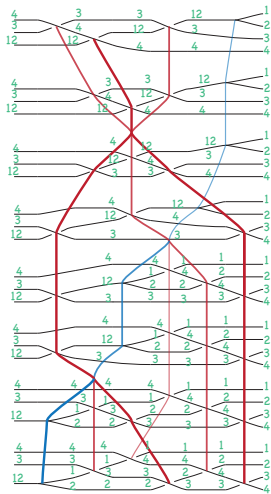
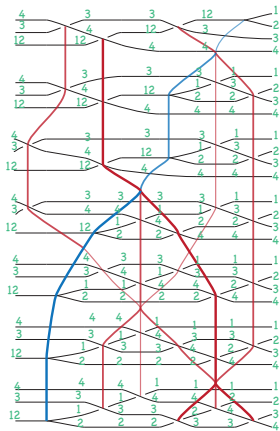


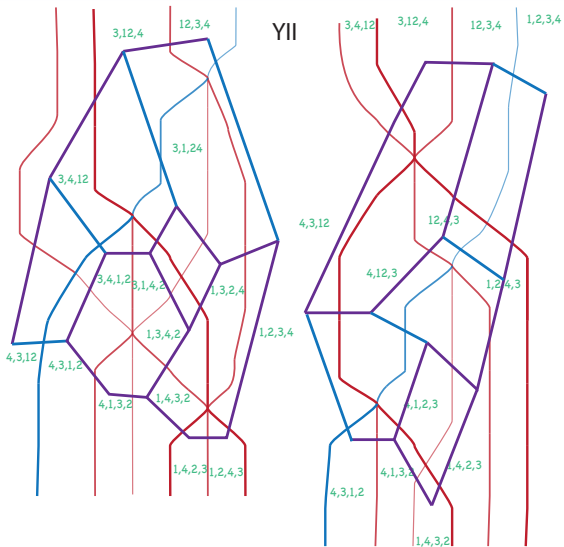
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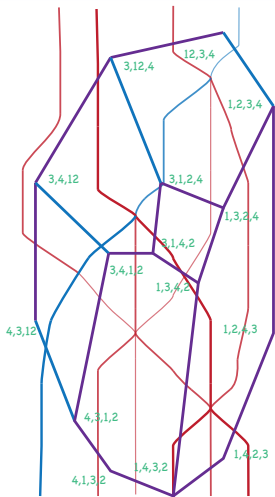


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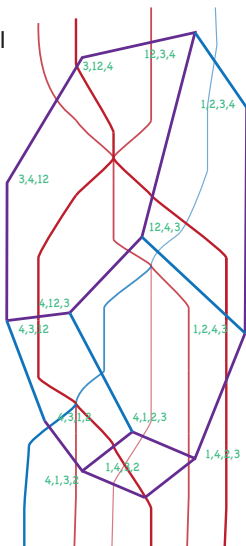
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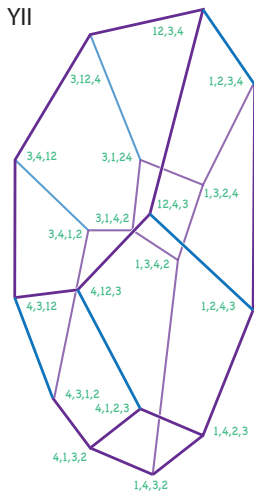
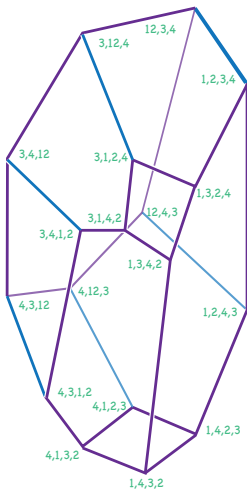




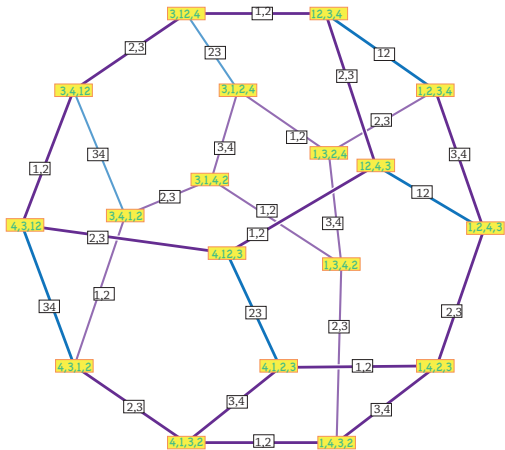


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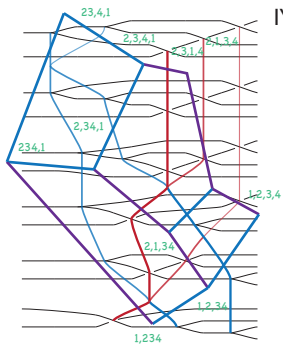




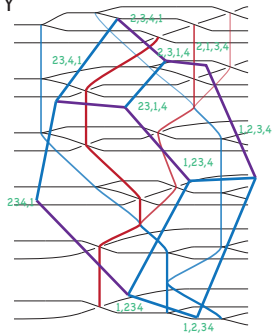
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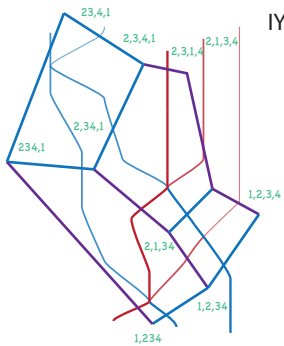


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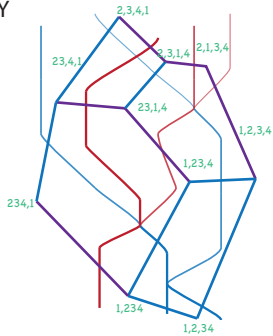


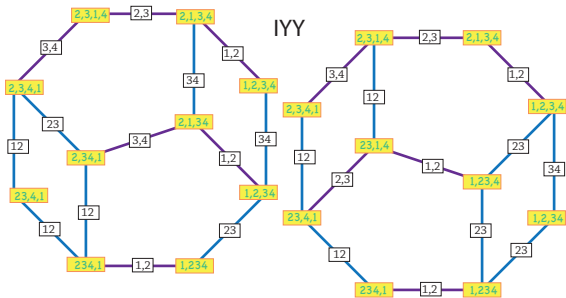
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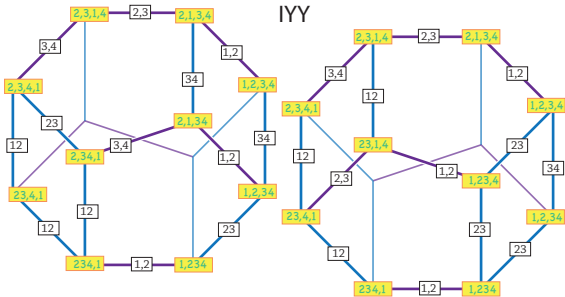


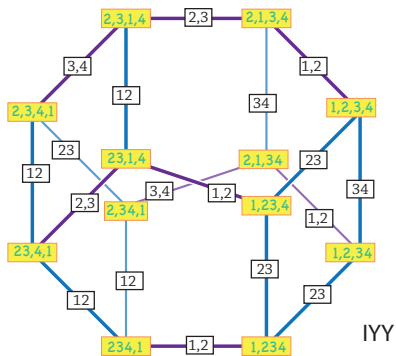


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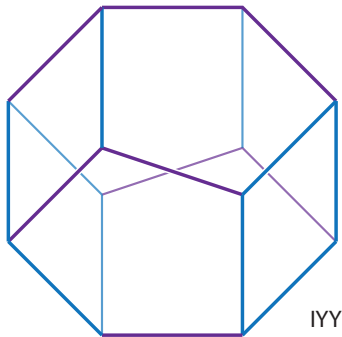




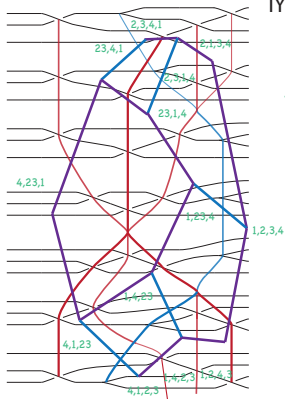




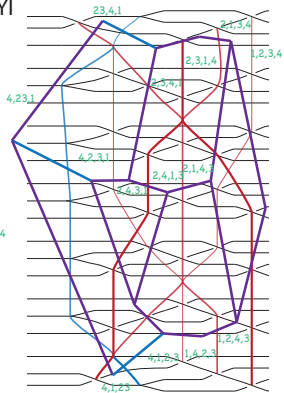
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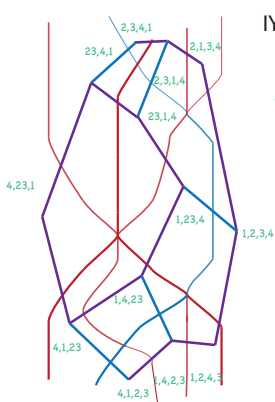


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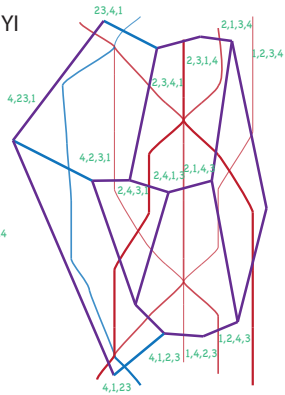


IV

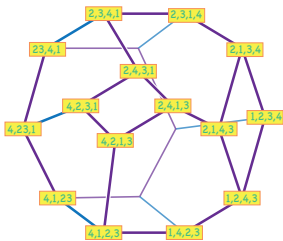
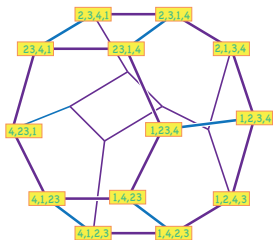


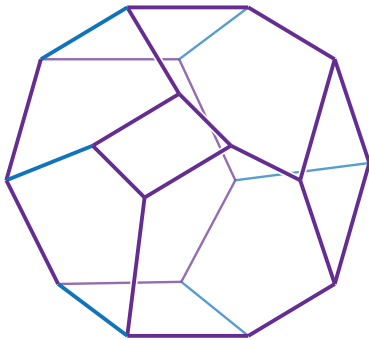


II



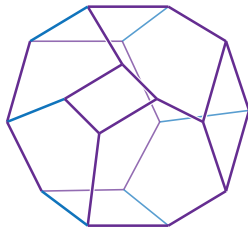
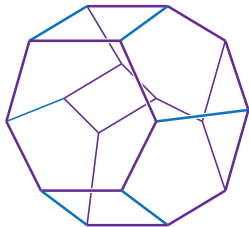
IV



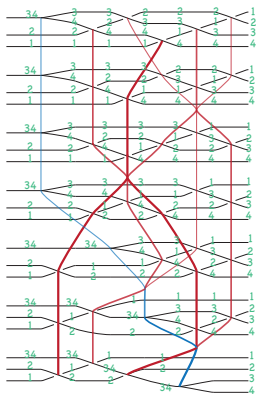


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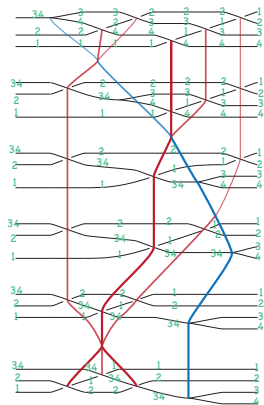
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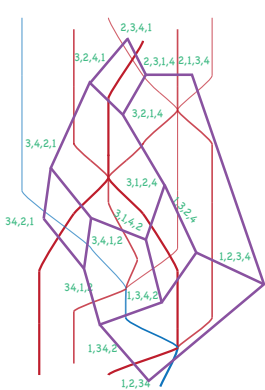


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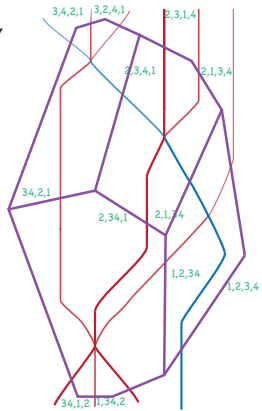


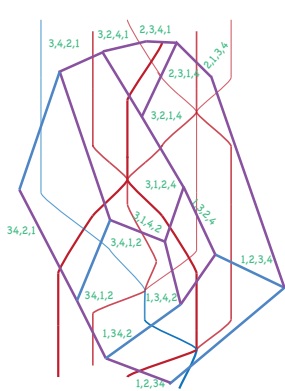
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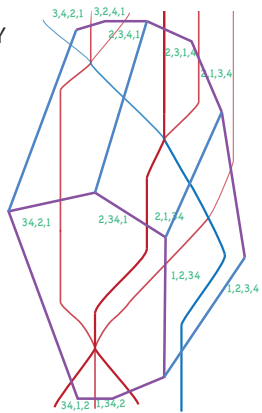


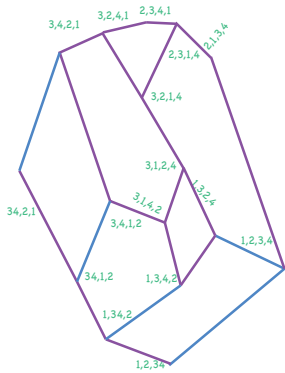
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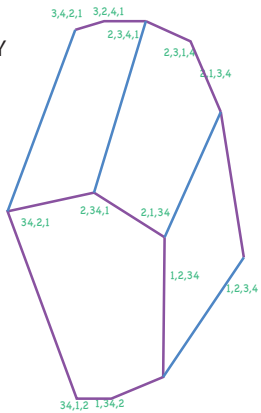


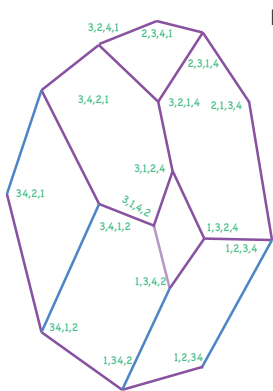
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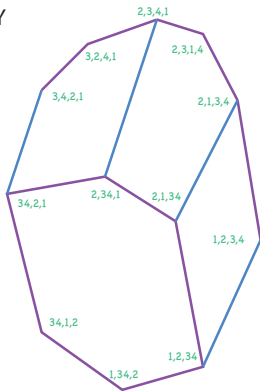


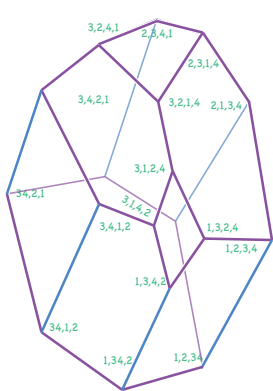
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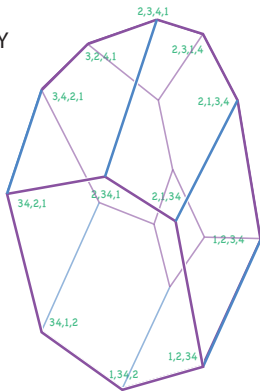


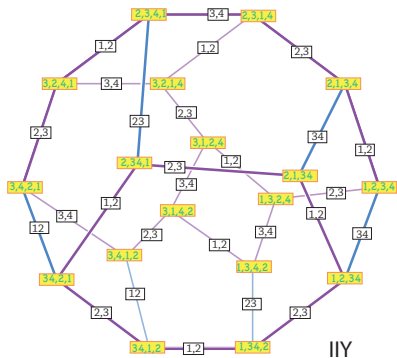
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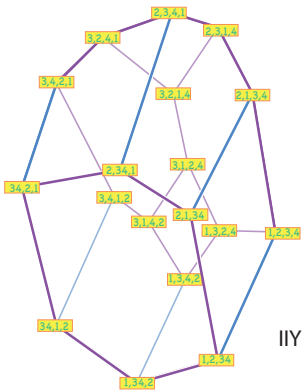


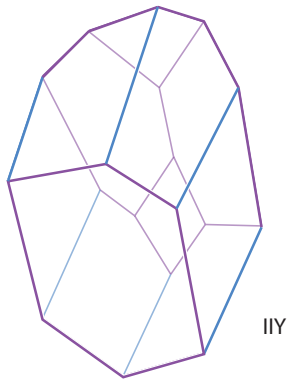


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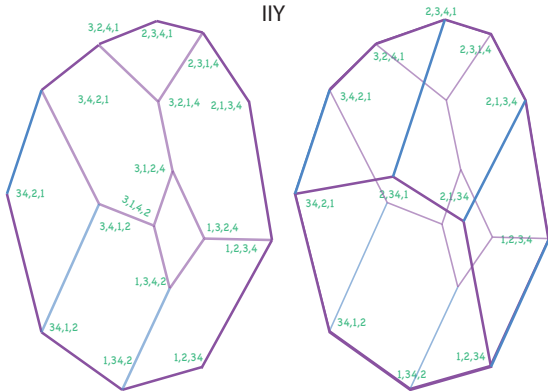




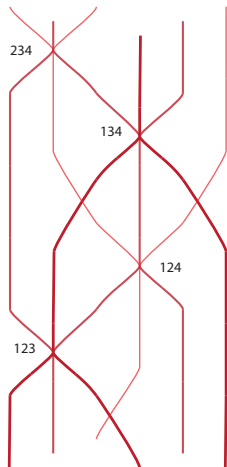
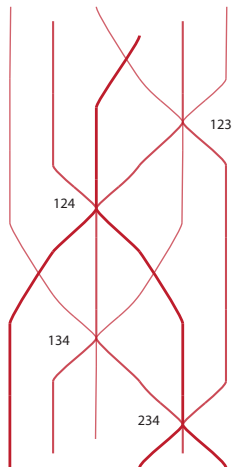




IY



Next we'll use the graphical structure to formulate a series of Abstract tensor equations.

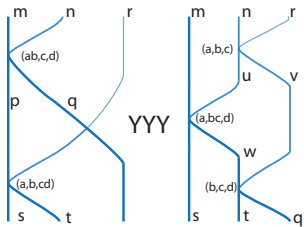


One of these equations is the Zamolodchikov tetrahedral equation.

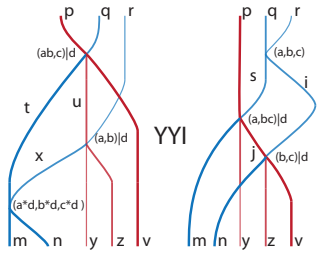
$$S_{123}S_{124}S_{134}S_{234} = S_{234}S_{134}S_{124}S_{123}.$$

One of these equations is the Zamolodchikov tetrahedral equation.

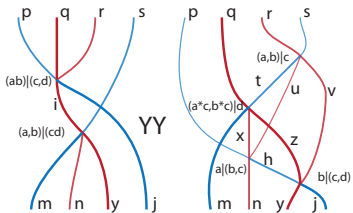
$$S_{123}S_{124}S_{134}S_{234} = S_{234}S_{134}S_{124}S_{123}.$$



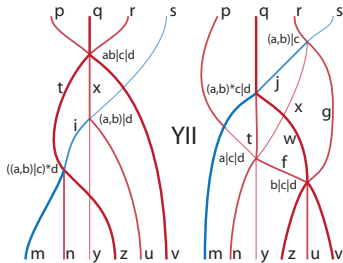
$$K_{p,q}^{m,n} K_{s,t}^{p,r} = K_{u,v}^{n,r} K_{s,w}^{m,u} K_{t,q}^{w,v}$$



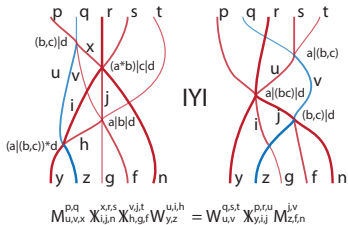
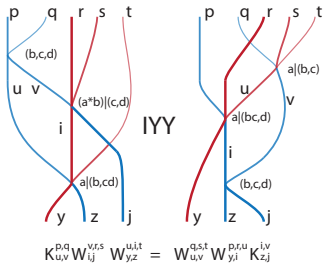
$$M_{t,u,v}^{p,q} M_{x,y,z}^{u,r} K_{m,n}^{t,x} = K_{s,i}^{q,r} M_{m,y,j}^{p,s} M_{n,z,v}^{j,i}$$

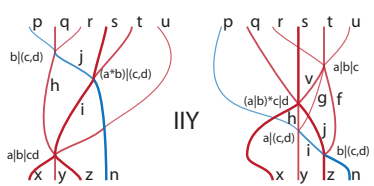


$$W_{ij}^{p,q,r} M_{m,n,y}^{i,s} = M_{t,u,v}^{r,s} M_{m,x,z}^{q,t} W_{n,h}^{p,x,u} W_{y,i}^{h,z,v}$$



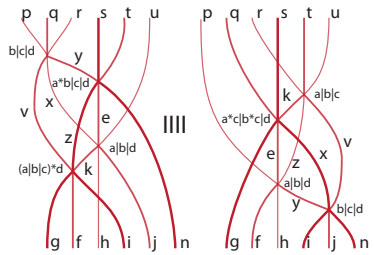
$$X_{t,x,u}^{p,q,r} M_{i,y,u}^{x,s} M_{m,n,z}^{t,i} = M_{j,x,g}^{r,s} M_{m,t,w}^{q,j} X_{n,y,f}^{p,t,x} X_{z,u,v}^{f,w,g}$$





IIY

$$W_{hj}^{p,q,r} W_{in}^{i,s,t} X_{xy,z}^{h,i,u} = X_{v,g,f}^{r,t,u} X_{x,h,j}^{q,s,v} W_{y,i}^{p,h,g} W_{z,n}^{i,j,f}$$



IIII

$$X_{v,x,y}^{p,q,r} X_{z,e,n}^{y,s,t} X_{k,h,j}^{x,e,u} X_{g,f,i}^{v,z,k} = X_{k,z,v}^{r,t,u} X_{g,e,x}^{q,s,k} X_{f,h,y}^{p,e,z} X_{i,j,n}^{y,x,v}$$

In 3-D

A	$Y_j^{\ell,n} Y_i^{j,k} = Y_p^{n,k} Y_i^{\ell,p}$
VI	$Y_d^{a,b} X_{e,f}^{d,c} = X_{c,h}^{b,c} X_{e,g}^{a,c} Y_f^{g,h}.$
IV	$X_{d,g}^{a,b} X_{h,f}^{g,c} Y_e^{d,h} = Y_i^{b,c} X_{e,f}^{a,i}$
III	$X_{d,e}^{a,b} X_{i,h}^{e,c} X_{f,g}^{d,i} = X_{j,k}^{b,c} X_{f,e}^{a,j} X_{g,h}^{e,k}$

$$Y_c^{a,b} = \begin{cases} 1 & \text{if } c = ab, \\ 0 & \text{otherwise;} \end{cases}$$

$$X_{c,d}^{a,b} = \begin{cases} 1 & \text{if } c = b \text{ \& } d = b^{-1}ab, \\ 0 & \text{otherwise.} \end{cases}$$

In 4-D

YYY	$K_{p,q}^{m,n} K_{s,t}^{p,r} = K_{u,v}^{n,r} K_{s,w}^{m,u} K_{t,q}^{w,v}$
YYI	$M_{t,u,v}^{p,q} M_{x,y,z}^{u,r} K_{m,n}^{t,x} = K_{s,i}^{q,r} M_{m,y,j}^{p,s} M_{n,z,v}^{j,i}$
YY	$W_{i,j}^{p,q,r} M_{m,n,y}^{i,s} = M_{t,u,v}^{r,s} M_{m,x,z}^{q,t} W_{n,h}^{p,x,u} W_{y,i}^{h,z,v}$
YII	$\mathbb{X}_{t,x,v}^{p,q,r} M_{y,u,v}^{x,s} M_{m,n,z}^{t,i} = M_{j,x,g}^{r,s} M_{m,t,x}^{q,j} \mathbb{X}_{n,y,f}^{p,t,x} \mathbb{X}_{z,u,v}^{f,w,g}$
IYY	$K_{u,v}^{p,q} W_{i,j}^{v,r,s} W_{y,z}^{u,i,t} = W_{u,v}^{q,s,t} W_{y,i}^{p,r,u} K_{z,j}^{i,v}$
IYI	$M_{u,v,x}^{p,q} \mathbb{X}_{i,j,n}^{x,r,s} \mathbb{X}_{h,g,f}^{v,j,t} W_{y,z}^{u,i,h} = W_{u,v}^{q,s,t} \mathbb{X}_{y,i,j}^{p,r,u} M_{z,f,n}^{j,v}$
IYY	$W_{h,j}^{p,q,r} W_{i,n}^{j,s,t} \mathbb{X}_{x,y,z}^{h,i,u} = \mathbb{X}_{v,g,f}^{r,t,u} \mathbb{X}_{x,h,j}^{q,s,v} W_{y,i}^{p,h,g} W_{z,n}^{i,j,f}$
IIII	$\mathbb{X}_{v,x,y}^{p,q,r} \mathbb{X}_{z,e,n}^{y,s,t} \mathbb{X}_{k,h,j}^{x,e,u} \mathbb{X}_{g,f,i}^{v,z,k} = \mathbb{X}_{k,z,v}^{r,t,u} \mathbb{X}_{g,e,x}^{q,s,k} \mathbb{X}_{f,h,y}^{p,e,z} \mathbb{X}_{i,j,n}^{y,x,v}$

$$K_{r,s}^{p,q} = \begin{cases} 1 & \text{if } \exists a, b, c \text{ s.t.} \\ & \begin{cases} p = ((ab, c), 0|0), & q = ((a, b), 0|0), \\ r = ((a, bc), 0|0), & \& s = ((b, c), 0|0). \end{cases} \\ 0 & \text{otherwise;} \end{cases}$$

$$M_{p,q,r}^{i,j} = \begin{cases} 1 & \text{if } \exists a, b, c \text{ s.t.} \\ & \begin{cases} i = ((0, 0), ab|c), & j = ((a, b), 0|0), \\ p = ((a, b) \triangleleft c, 0|0), & q = ((0, 0), a|c), \\ \& r = ((0, 0), b|c), \end{cases} \\ 0 & \text{otherwise;} \end{cases}$$

$$W_{p,q}^{i,j,\ell} = \begin{cases} 1 & \text{if } \exists a, b, c \text{ s.t.} \\ & \begin{cases} i = ((b, c), 0|0), & j = ((0, 0), (a \triangleleft b)|c), \\ \ell = ((0, 0), a|b), \\ p = ((0, 0), a|(bc)), & \& q = ((b, c), 0|0), \end{cases} \\ 0 & \text{otherwise;} \end{cases}$$

$$\chi_{s,t,u}^{p,q,r} = \begin{cases} 1 & \text{if } \exists a, b, c \text{ s.t.} \\ & \begin{cases} p = ((0, 0), b|c), & q = ((0, 0), (a \triangleleft b)|c), \\ r = ((0, 0), a|b), \\ s = ((0, 0), (a|b) \triangleleft c), & t = ((0, 0), a|c), \\ \& u = ((0, 0), b|c), \end{cases} \\ 0 & \text{otherwise;} \end{cases}$$

where, *e.g.*, $a \triangleleft c = c^{-1}ac$, $(a, b) \triangleleft c = (a \triangleleft c, b \triangleleft c)$,
and $(a|b) \triangleleft c = (a \triangleleft c)|(b \triangleleft c)$

These tensors form a solution set over a module whose basis is determined by the underlying algebraic structure precisely because $\partial \circ \partial = 0$ in the chain complex.

Thanks

Thank you for your attention!